

Tutorial on Fairness Modeling

Part 1: Fairness in Optimization Models

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Two Tutorials

- *This tutorial:* modeling **fairness in optimization models**
 - Social welfare functions that incorporate fairness.
 - Practical LP/MILP/NLP models.
 - A bit of social choice theory.
- *Next tutorial:* modeling **group fairness in AI**
 - Crash course in deontological ethics.
 - Group parity metrics & their assessment.
 - Connections with social welfare functions.

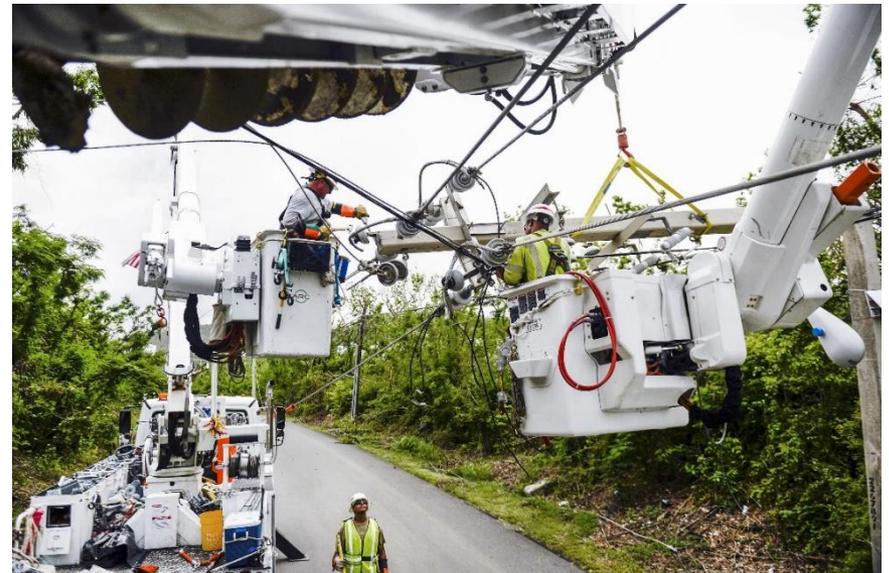
Modeling Fairness

- A growing interest in incorporating **fairness** into **optimization models...**
 - Health care resources.
 - Facility location (e.g., emergency services, infrastructure).
 - Telecommunications.
 - Traffic signal timing
 - Disaster recovery (e.g., power restoration)...



Modeling Fairness

- Example: disaster relief
 - Power restoration can focus on **urban** areas first (**efficiency**).
 - This can leave rural areas without power for weeks/months.
 - This happened in Puerto Rico after Hurricane Maria (2017).
- A more **equitable** solution
 - ...would give some priority to rural areas without overly sacrificing efficiency.



Modeling Fairness

- It is far from obvious how to formulate equity concerns **mathematically**.
 - Less straightforward than maximizing total benefit or minimizing total cost.
 - Still less obvious how to **combine** equity with total benefit.



Modeling Fairness

- There is **no one** concept of equity or fairness.
 - The appropriate concept **depends on the application.**
- We therefore survey a range of formulations.
 - Describe their **mathematical properties.**
 - Indicate their **strengths** and **weaknesses.**
 - State what appears to be the **most practical model.**
 - So that one can select the formulation that **best suits** a given application.
- Also a brief excursion into **social choice theory.**
 - ...and into **structural properties** of fair solutions.

References

- References and more details may be found in

V. Chen & J. N. Hooker, [A guide to formulating equity and fairness in an optimization model](#), *Annals of OR*, 2023.

Inequality measures

Criterion	Linear?	Contin?
Relative range	yes	yes
Relative mean deviation	yes	yes
Coefficient of variation	no	yes
Gini coefficient	yes	yes
Hoover index	yes	yes

Fairness for the disadvantaged

Criterion	Linear?	Contin?
Maximin (Rawlsian)	yes	yes
Leximax (lexicographic)	yes	yes
McLoone index	yes	no

Linear = fairness model introduces only **linear** expressions
Contin. = fairness model introduces only **continuous** variables

Combining efficiency & fairness

Convex combinations

Criterion	Linear?	Contin?
Utility + Gini coefficient	no	yes
Utility * Gini coefficient	yes	yes
Utility + maximin	yes	yes

Combining efficiency & fairness

Classical methods

Criterion	Linear?	Contin?
Alpha fairness	yes	yes
Proportional fairness (Nash bargaining)	yes	yes
Kalai-Smorodinsky bargaining	no	yes

Linear = fairness model introduces only **linear** expressions
Contin. = fairness model introduces only **continuous** variables

Combining efficiency & fairness

Threshold methods

Criterion	Linear?	Contin?
Utility + maximin – Utility threshold	yes	no
Utility + maximin – Equity threshold	yes	yes
Utility + leximax – Predefined priorities	yes	no
Utility + leximax – No predefined priorities	yes	no

Linear = fairness model introduces only **linear** expressions

Contin. = fairness model introduces only **continuous** variables

Generic Model

- We formulate each fairness criterion as a **social welfare function (SWF)**.

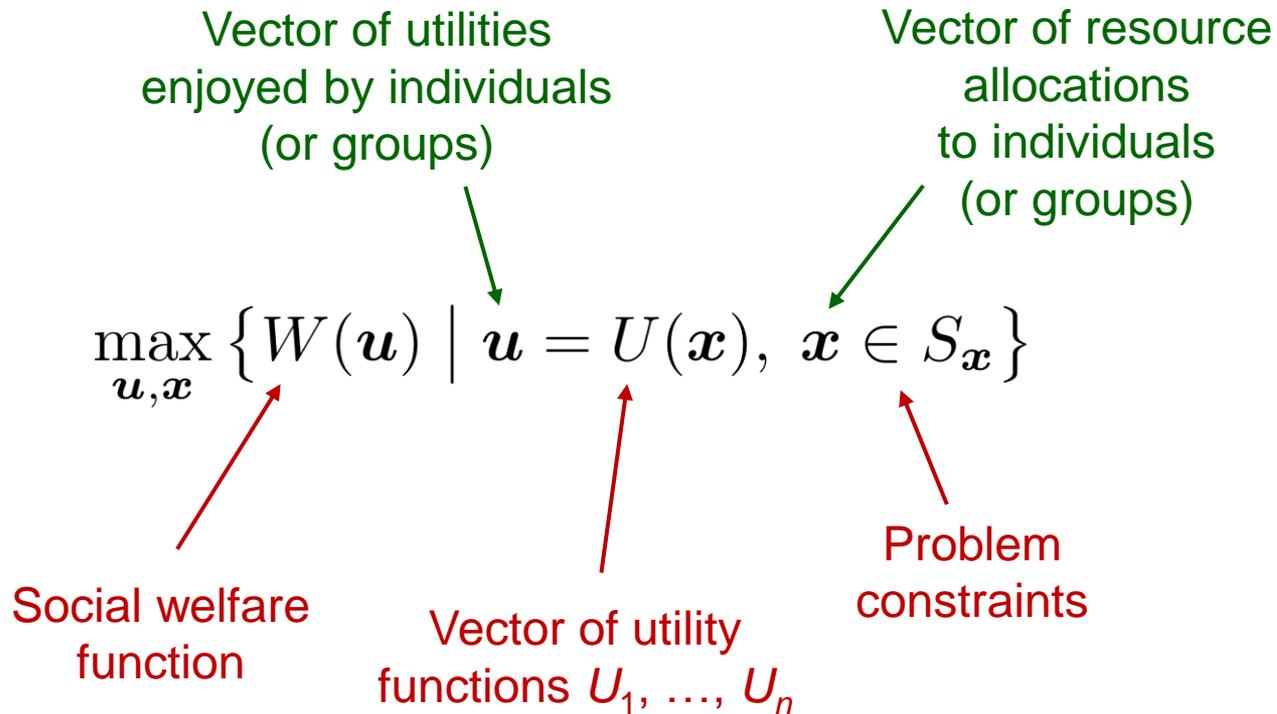
Individual utilities

$$W(\mathbf{u}) = W(u_1, \dots, u_n)$$


- Measures desirability of the **magnitude and distribution of utilities** across individuals.
- **Utility** can be wealth, health, negative cost, etc.
- The **SWF** becomes the **objective function** of the optimization model.

Generic Model

The social welfare optimization problem



Generic Model

Example – *Medical triage*

Utility functions are $U_i(\mathbf{x}) = \alpha_i + \beta_i x_i$.

QALYs without treatment

Additional QALYs due to treatment

resources allocated to patient group i

$$\max_{\mathbf{u}, \mathbf{x}} \left\{ W(\mathbf{u}) \left| \begin{array}{l} u_i = \alpha_i + \beta_i x_i, 0 \leq x_i \leq d'_i, \text{ all } i \\ \sum_i a'_i x_i \leq B' \end{array} \right. \right\}$$

Social welfare function

Budget constraint

Bounds on group i resource consumption

Generic Model

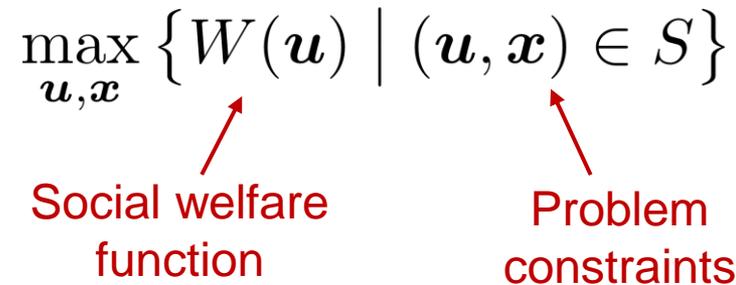
The social welfare optimization problem

Incorporate $\mathbf{u} = U(\mathbf{x})$ into problem constraints.

$$\max_{\mathbf{u}, \mathbf{x}} \{ W(\mathbf{u}) \mid (\mathbf{u}, \mathbf{x}) \in S \}$$

Social welfare function

Problem constraints

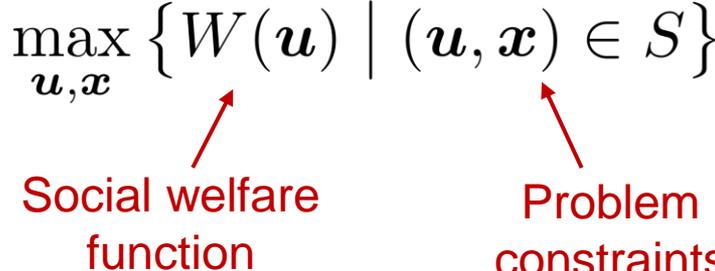


Generic Model

The social welfare optimization problem

Incorporate $\mathbf{u} = U(\mathbf{x})$ into problem constraints.

$$\max_{\mathbf{u}, \mathbf{x}} \{ W(\mathbf{u}) \mid (\mathbf{u}, \mathbf{x}) \in S \}$$


Social welfare function Problem constraints

In the triage problem, we can eliminate x_i because $u_i = \alpha_i + \beta_i x_i$:

$$\max_{\mathbf{u}, \mathbf{x}} \left\{ W(\mathbf{u}) \mid \sum_i a_i u_i \leq B, \quad c_i \leq u_i \leq d_i \right\}$$

where $a_i = \frac{a'_i}{\beta_i}$, $B = B' + \sum_i \frac{a'_i \alpha_i}{\beta_i}$, $(c_i, d_i) = (\alpha_i \beta_i, d'_i)$.

Inequality Measures

<i>Criterion</i>	<i>Linear?</i>	<i>Conti?</i>
Relative range	yes	yes
Relative mean deviation	yes	yes
Coefficient of variation	no	yes
Gini coefficient	yes	yes
Hoover index	yes	yes

Inequality Measures

Equality vs fairness

Two views on ethical importance of equality:

Parfit 1997

- **Irreducible:** Inequality is inherently unfair.
- **Reducible:** Inequality is unfair only insofar as it reduces utility.

Scanlon 2003

Frankfurt 2015

Possible problems with inequality measures:

- No preference for an identical distribution with **higher utility**.
- Even when average utility is fixed, no preference for reducing inequality at the **bottom** rather than the **top** of the distribution.

Inequality Measures

Equality vs fairness

We can perhaps agree on this much:

- Equality is **not the same concept** as fairness, even when it is closely related.
- An inequality metric can be appropriate when a specifically **egalitarian** distribution is the goal, **without regard** to efficiency and other forms of equity.

Inequality Measures

Relative range

$$W(\mathbf{u}) = \frac{u_{\max} - u_{\min}}{\bar{u}}$$

Rationale:

- Perceived inequality is relative to the best off.
- So, move everyone closer to the best off.

Problem:

- Ignores distribution **between** extremes.

Inequality Measures

Relative range

- Problem is **linearized** using same change of variable as in linear-fractional programming.

Let $\mathbf{u} = \mathbf{u}'/t$ and $\mathbf{x} = \mathbf{x}'/t$. The optimization problem is

$$\min_{\substack{\mathbf{x}', \mathbf{u}', t \\ u'_{\min}, u'_{\max}}} \left\{ u'_{\max} - u'_{\min} \mid \begin{array}{l} u'_{\min} \leq u'_i \leq u'_{\max}, \text{ all } i \\ \bar{u}' = 1, t \geq 0, (\mathbf{u}', \mathbf{x}') \in S' \end{array} \right\}$$

where t, u'_{\min}, u'_{\max} are new variables.

Charnes & Cooper 1962

Inequality Measures

Relative range

Model:

$$\min_{\substack{\mathbf{x}', \mathbf{u}', t \\ u'_{\min}, u'_{\max}}} \left\{ u'_{\max} - u'_{\min} \mid \begin{array}{l} u'_{\min} \leq u'_i \leq u'_{\max}, \text{ all } i \\ \bar{u}' = 1, t \geq 0, (\mathbf{u}', \mathbf{x}') \in S' \end{array} \right\}$$

The difficulty of constraints $(\mathbf{u}', \mathbf{x}') \in S'$ depends on nature of S .

If S is linear $A\mathbf{u} + B\mathbf{x} \leq \mathbf{b}$, it remains linear: $A\mathbf{u}' + B\mathbf{x}' \leq t\mathbf{b}$.

If S is $\mathbf{g}(\mathbf{u}, \mathbf{x}) \leq \mathbf{b}$ for homogeneous \mathbf{g} , it retains almost the same form: $\mathbf{g}(\mathbf{u}', \mathbf{x}') \leq t\mathbf{b}$.

Inequality Measures

Relative mean deviation

$$W(\mathbf{u}) = -\frac{1}{\bar{u}} \sum_i |u_i - \bar{u}|$$

Rationale:

- Considers all utilities.

Model:

- Again, linearized by change of variable.

$$\min_{\mathbf{x}', \mathbf{u}', \mathbf{v}, t} \left\{ \sum_i v_i \mid \begin{array}{l} -v_i \leq u'_i - \bar{u}' \leq v_i, \text{ all } i \\ \bar{u}' = 1, t \geq 0, (\mathbf{u}', \mathbf{x}') \in S' \end{array} \right\}$$

where \mathbf{v} is vector of new variables.

Inequality Measures

Coefficient of variation

$$W(\mathbf{u}) = -\frac{1}{\bar{u}} \left[\frac{1}{n} \sum_i (u_i - \bar{u})^2 \right]^{\frac{1}{2}}$$

Rationale:

- Familiar. Outliers receive extra weight.

Problem:

- Nonlinear (but convex)

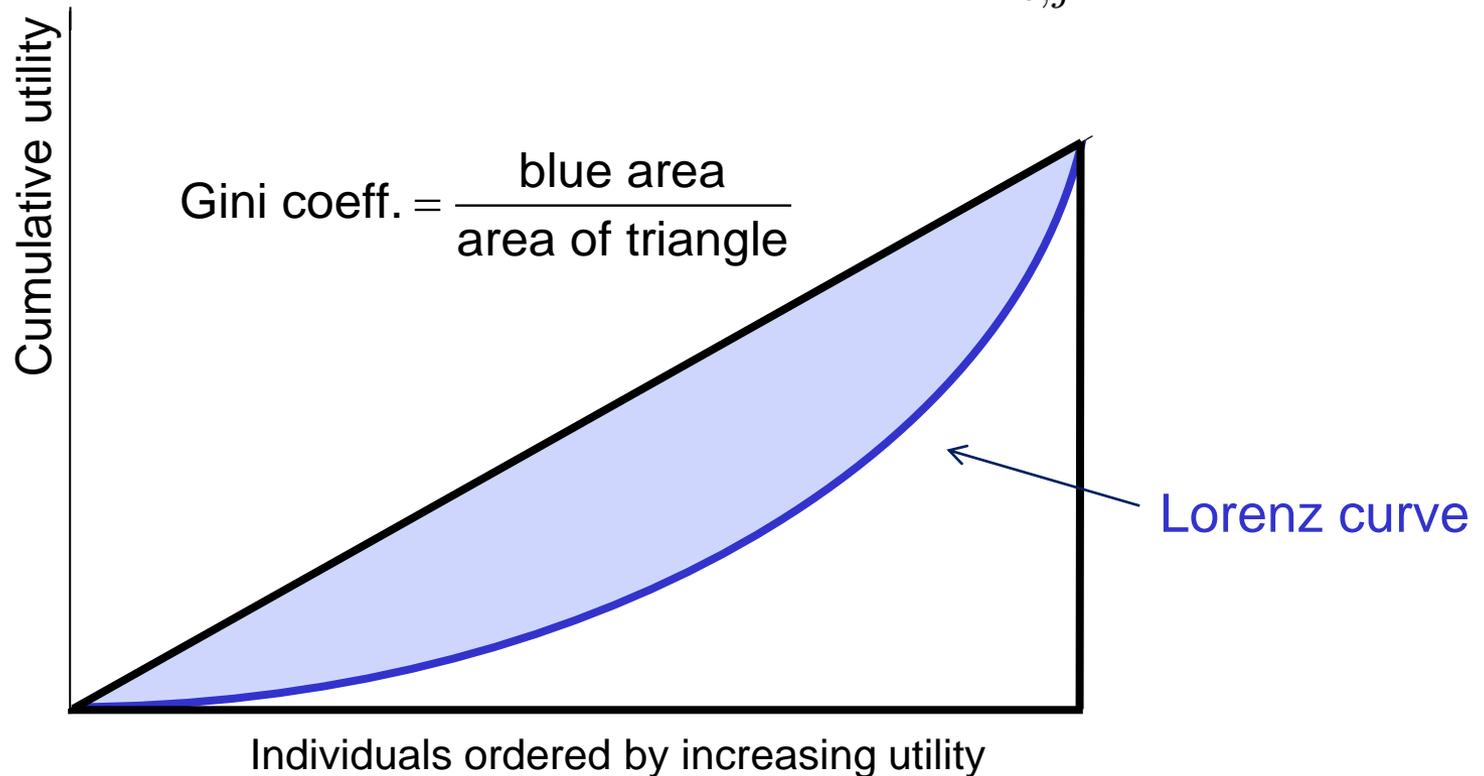
Model:

$$\min_{\mathbf{x}', \mathbf{u}', \mathbf{v}, t} \left\{ \frac{1}{n} \sum_i (u'_i - \bar{u}')^2 \mid \begin{array}{l} \bar{u}' = 1, t \geq 0 \\ (\mathbf{u}', \mathbf{x}') \in S' \end{array} \right\}$$

Inequality Measures

Gini coefficient

$$W(\mathbf{u}) = -G(\mathbf{u}), \quad \text{where } G(\mathbf{u}) = \frac{1}{2\bar{u}n^2} \sum_{i,j} |u_i - u_j|$$



Inequality Measures

Gini coefficient

$$W(\mathbf{u}) = -G(\mathbf{u}), \quad \text{where } G(\mathbf{u}) = \frac{1}{2\bar{u}n^2} \sum_{i,j} |u_i - u_j|$$

Rationale:

- Relationship to Lorenz curve.
- Widely used.

Model:

- Linear:
$$\min_{\mathbf{x}', \mathbf{u}', V, t} \left\{ \frac{1}{2n^2} \sum_{i,j} v_{ij} \mid \begin{array}{l} -v_{ij} \leq u'_i - u'_j \leq v_{ij}, \text{ all } i, j \\ \bar{u}' = 1, t \geq 0, (\mathbf{u}', \mathbf{x}') \in S' \end{array} \right\}$$

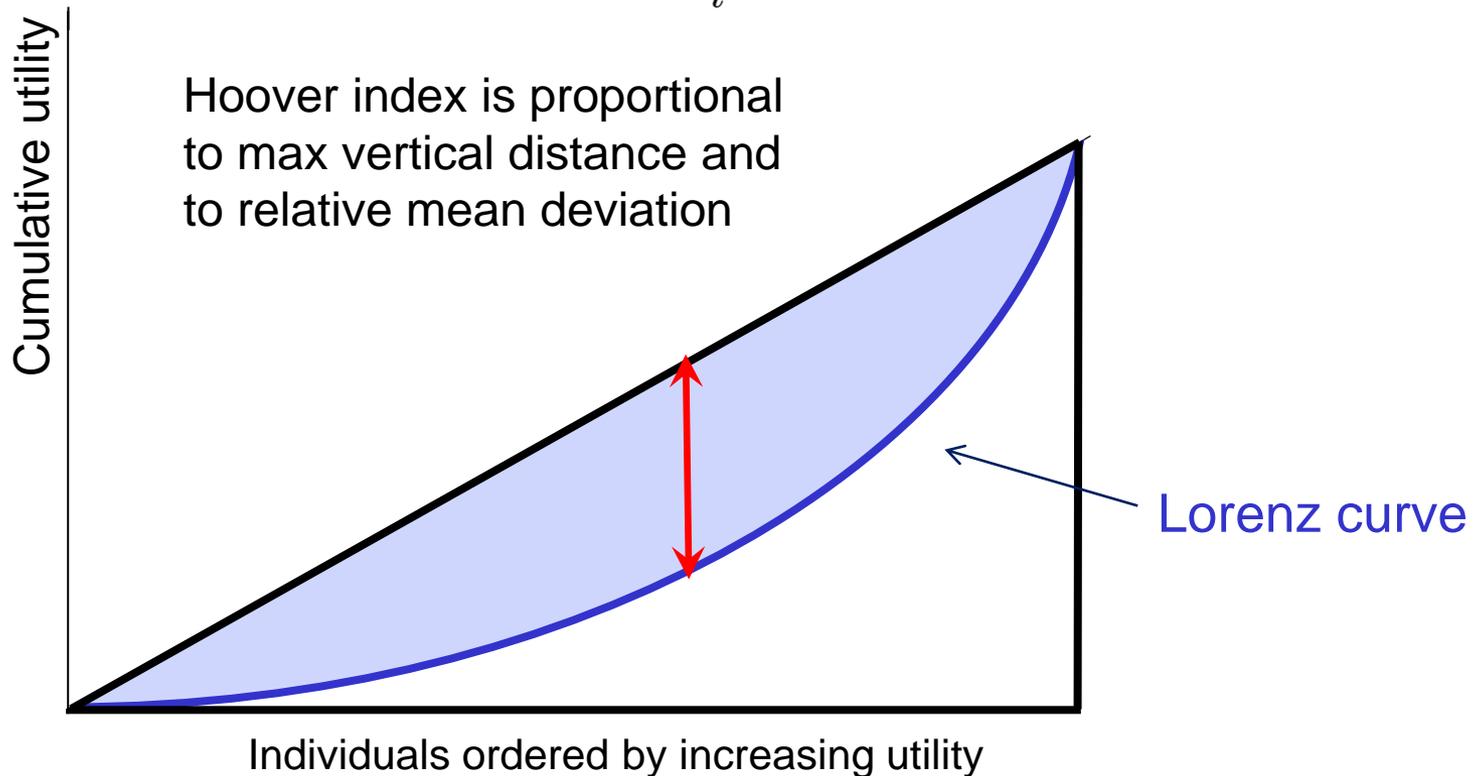
where V is a matrix of new variables.

Inequality Measures

Hoover index

$$W(\mathbf{u}) = -\frac{1}{2n\bar{u}} \sum_i |u_i - \bar{u}|$$

Hoover 1936



Inequality Measures

Hoover index

$$W(\mathbf{u}) = -\frac{1}{2n\bar{u}} \sum_i |u_i - \bar{u}|$$

Rationale:

- Hoover index is fraction of total utility that would have to be redistributed to achieve perfect equality.

Model:

- Same as relative mean deviation.

Fairness for the Disadvantaged

Criterion	<i>Linear?</i>	<i>Contin?</i>
Maximin (Rawlsian)	yes	yes
Leximax (lexicographic)	yes	yes
McLoone index	yes	no

Fairness for the Disadvantaged

Maximin

$$W(\mathbf{u}) = \min_i \{u_i\}$$

Rationale:

- Based on **difference principle** of John Rawls.
- Inequality is justified only to the extent that it increases the utility of the worst-off.
- Originally intended only for the design of **social institutions** and distribution of **primary goods** (goods that any rational person would want).
- Can be adopted as a general principle of equity: maximize the minimum utility.

Rawls 1971, 1999

Fairness for the Disadvantaged

Maximin

$$W(\mathbf{u}) = \min_i \{u_i\}$$

Social contract argument:

- We decide on social policy in an “original position,” behind a “veil of ignorance” as to our position on society.
- All parties must be willing to **endorse** the policy, no matter what position they end up assuming.
- No rational person can endorse a policy that puts him/her on the **bottom** of society – unless that person would be even **worse off** under another social arrangement.
- Therefore, an agreed-upon social policy must maximize the welfare of the worst-off.

Fairness for the Disadvantaged

Maximin

$$W(\mathbf{u}) = \min_i \{u_i\}$$

Model: $\max_{\mathbf{x}, \mathbf{u}, w} \{w \mid w \leq u_i, \text{ all } i; (\mathbf{u}, \mathbf{x}) \in S\}$

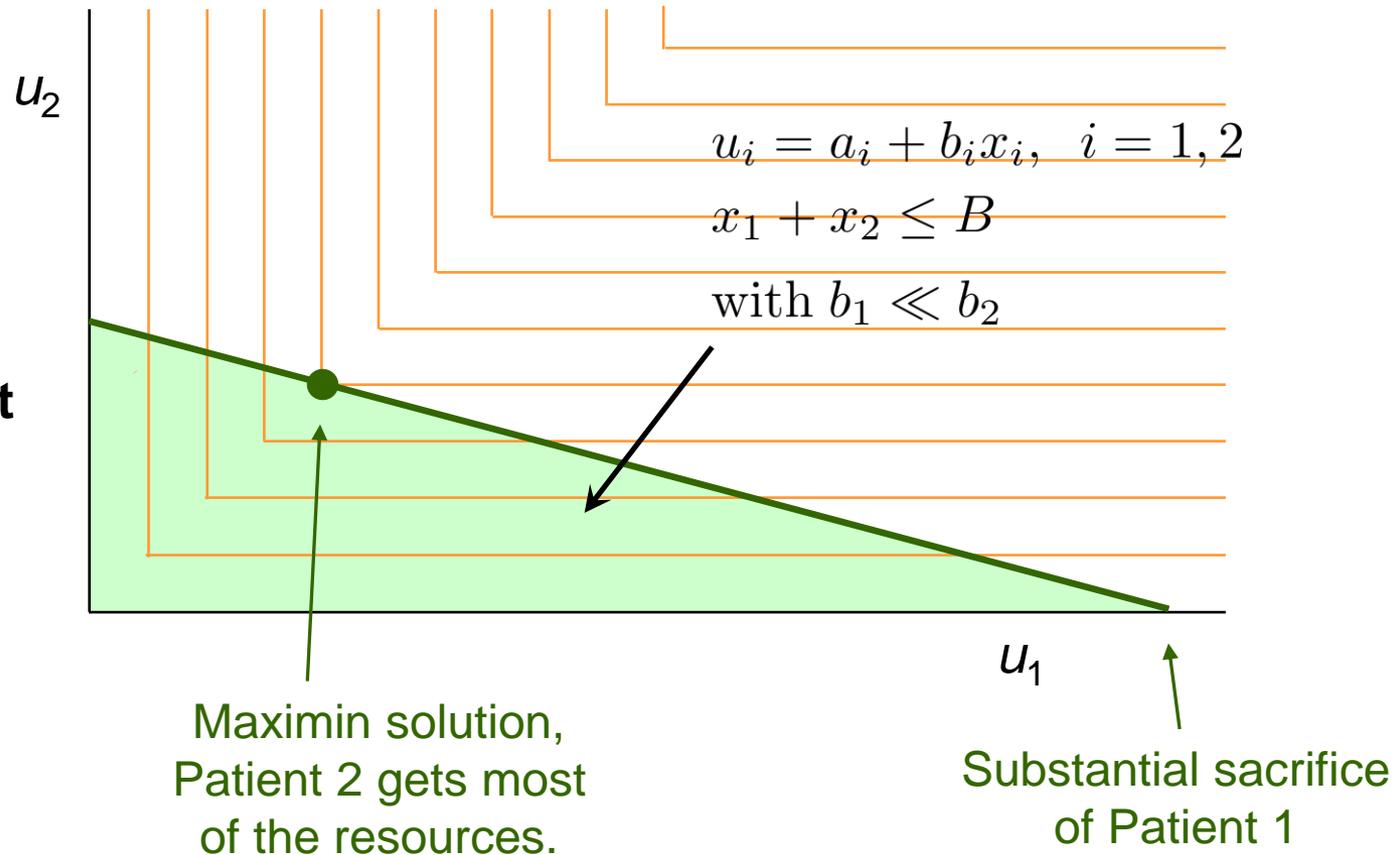
Problems:

- Can force equality even when this is extremely costly in terms of total utility.
- Does not care about 2nd worst off, etc., and so can waste resources.

Fairness for the Disadvantaged

Maximin

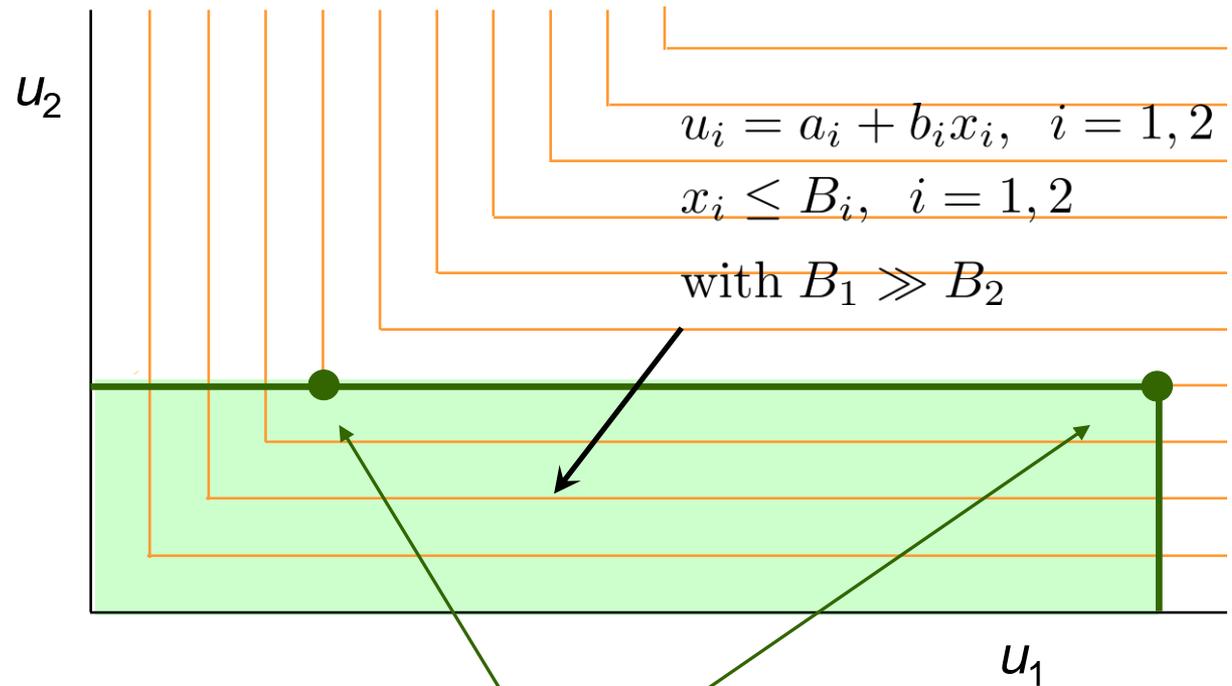
Medical example
with
budget constraint



Fairness for the Disadvantaged

Maximin

Medical example
with
resource bounds



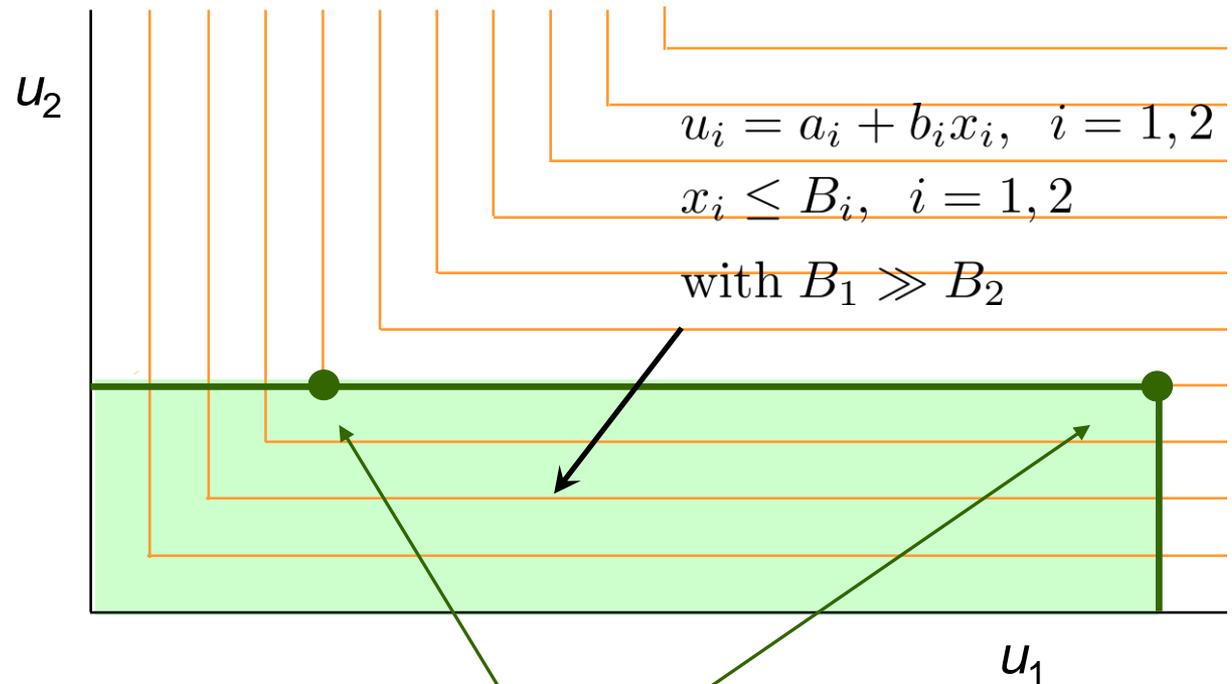
These solutions have same social welfare!

Fairness for the Disadvantaged

Maximin

Medical example
with
resource bounds

Remedy: use
leximax solution



These solutions have same social welfare!

Fairness for the Disadvantaged

Leximax

Rationale:

- Takes in account 2nd worst-off, etc., and avoids wasting utility.
- Can be justified with Rawlsian argument.

Model:

Solve sequence of optimization problems

$$\max_{\mathbf{x}, \mathbf{u}, w} \left\{ w \mid \begin{array}{l} w \leq u_i, \quad u_i \geq \hat{u}_{i_{k-1}}, \quad i \in I_k \\ (\mathbf{u}, \mathbf{x}) \in S \end{array} \right\}$$

for $k = 1, \dots, n$, where i_k is defined so that $\hat{u}_{i_k} = \min_{i \in I_k} \{\hat{u}_i\}$, and where $I_k = \{1, \dots, n\} \setminus \{i_1, \dots, i_{k-1}\}$, $(\bar{\mathbf{x}}, \bar{\mathbf{u}})$ is an optimal solution of problem k , and $\hat{u}_{i_0} = -\infty$.

If $\hat{u}_j = \min_{i \in I_k} \{\hat{u}_i\}$ for multiple j , **must enumerate all solutions** that result from breaking the tie.

Fairness for the Disadvantaged

McLoone index

$$W(\mathbf{u}) = \frac{1}{|I(\mathbf{u})|\tilde{u}} \sum_{i \in I(\mathbf{u})} u_i$$

where \tilde{u} is the median of utilities in \mathbf{u} and $I(\mathbf{u})$ is the set of indices of utilities at or below the median

Rationale:

- Compares total utility of those at or below the median to the utility that would result from bringing them up to the median.
- Index = 1 if no one is below median, \rightarrow 0 for long lower tail.
- Focus on **all** the **disadvantaged**.
- Often used for public goods (e.g., educational benefits).

Fairness for the Disadvantaged

McLoone index

Model: Nonlinear, requires 0-1 variables.

$$\max_{\substack{\mathbf{x}, \mathbf{u}, m \\ \mathbf{y}, \mathbf{z}, \delta}} \left\{ \begin{array}{l} \frac{\sum_i y_i}{\sum_i z_i} \quad \left| \quad \begin{array}{l} m - M\delta_i \leq u_i \leq m + M(1 - \delta_i), \text{ all } i \\ y_i \leq u_i, y_i \leq M\delta_i, \delta_i \in \{0, 1\}, \text{ all } i \\ z_i \geq 0, z_i \geq m - M(1 - \delta_i), \text{ all } i \\ \sum_i \delta_i \leq n/2, (\mathbf{u}, \mathbf{x}) \in S \end{array} \right. \end{array} \right.$$

Linearize with change of variable, obtain MILP.

$$\max_{\substack{\mathbf{x}', \mathbf{u}', m' \\ \mathbf{y}', \mathbf{z}', t, \delta}} \left\{ \begin{array}{l} \sum_i y'_i \quad \left| \quad \begin{array}{l} u'_i \geq m' - M\delta_i, \text{ all } i \\ u'_i \leq m' + M(1 - \delta_i), \text{ all } i \\ y'_i \leq u'_i, y'_i \leq M\delta_i, \delta_i \in \{0, 1\}, \text{ all } i \\ z'_i \geq 0, z'_i \geq m' - M(1 - \delta_i), \text{ all } i \\ \sum_i z'_i = 1, t \geq 0 \\ \sum_i \delta_i \leq n/2, (\mathbf{u}', \mathbf{x}') \in S' \end{array} \right. \end{array} \right.$$

Social Choice Theory

- The economics literature derives social welfare functions from **axioms of rational choice**.
- The social welfare function depends on degree of **interpersonal comparability** of utilities.
- Arrow's impossibility theorem was the first result, but there are many others.

Social Choice Theory

Axioms

Anonymity (symmetry)

Social preferences are the same if indices of u_i s are permuted.

Strict pareto

If $\mathbf{u} > \mathbf{u}'$, then \mathbf{u} is preferred to \mathbf{u}' .

Independence

The preference of \mathbf{u} over \mathbf{u}' depends only on \mathbf{u} and \mathbf{u}' and not on what other utility vectors are possible.

Separability

Individuals i for which $u_i = u'_i$ do not affect the relative ranking of \mathbf{u} and \mathbf{u}' .

Social Choice Theory

Interpersonal comparability

- The properties of social welfare functions that satisfy the axioms depend on the degree to which utilities can be **compared** across individuals.

Invariance transformations

- These are transformations of utility vectors that indicate the degree of interpersonal comparability.
- Applying an invariance transformation to utility vectors does not change the **ranking** of distributions.

An invariance transformation has the form $\phi = (\phi_1, \dots, \phi_n)$, where ϕ_i is a transformation of individual utility i .

Social Choice Theory

Unit comparability.

- Invariance transformation has the form $\phi_i(u_i) = \beta u_i + \gamma_i$
- So, it is possible to compare utility **differences** across individuals:
 $u'_i - u_i > u'_j - u_j$ if and only if $\phi_i(u'_i) - \phi_i(u_i) > \phi_j(u'_j) - \phi_j(u_j)$

Theorem. Given anonymity, strict pareto, and independence axioms, the social welfare criterion must be **utilitarian**.

$$W(\mathbf{u}) = \sum_i u_i$$

Social Choice Theory

Level comparability.

- Invariance transformation has the form

$$\phi(\mathbf{u}) = (\phi_0(u_1), \dots, \phi_0(u_n))$$

where ϕ_0 is strictly increasing.

- So, it is possible to compare utility **levels** across individuals.

$$u_i > u_j \text{ if and only if } \phi_i(u_i) > \phi_j(u_j)$$

Theorem. Given anonymity, strict pareto, independence, and separability axioms, the social welfare criterion must be **maximin** or **minimax**.

Social Choice Theory

Problem with the utilitarian proof.

- The proof assumes that utilities have **no more** than unit comparability.
- This immediately rules out a maximin criterion, since identifying the minimum utility presupposes that utility **levels** can be compared.

Problem with the maximin proof.

- The proof assumes that utilities have **no more** than level comparability.
- This immediately rules out criteria that consider the spread of utilities.
- So, it rules out all the criteria we consider after maximin.

Utility & Fairness – Convex Combinations

Criterion	Linear?	Contin?
Utility + Gini coefficient	no	yes
Utility * Gini coefficient	yes	yes
Utility + maximin	yes	yes

Utility & Fairness – Convex Combinations

Utility + Gini coefficient

$$W(\mathbf{u}) = (1 - \lambda) \sum_i u_i + \lambda(1 - G(\mathbf{u}))$$

Rationale.

- Takes into account both efficiency and equity.
- Allows one to adjust their relative importance.

Problem.

- Combines utility with a dimensionless quantity.
- How to interpret λ , or choose a λ for a given application?
- Choice of λ is an issue with convex combinations in general.

Utility & Fairness – Convex Combinations

Utility * Gini coefficient

$$W(\mathbf{u}) = (1 - G(\mathbf{u})) \sum_i u_i$$

Rationale.

Eisenhandler & Tzur 2019

- Gets rid of λ .
- Equivalent to SWF that is easily linearized:

$$W(\mathbf{u}) = \sum_i u_i - \frac{1}{n} \sum_{i < j} |u_j - u_i|$$

Problem.

- It is still a convex combination of utility and an equality metric (mean absolute difference).
- Implicit multiplier $\lambda = 1/2$. Why this multiplier?

Utility & Fairness – Convex Combinations

Utility + Gini-weighted utility

$$W(\mathbf{u}) = \sum_i u_i + \mu(1 - G(\mathbf{u})) \sum_i u_i$$

Rationale.

- Combines quantities measured in same units.

Mostajabdaveh, Gutjahr, Salman 2019

Problem.

- Equivalent to utility*(1-Gini) with multiplier $\lambda = \mu(1 + 2\mu)^{-1}$.
- How to interpret μ ?

Utility & Fairness – Convex Combinations

Utility + Maximin

$$W(\mathbf{u}) = (1 - \lambda) \sum_i u_i + \lambda \min_i \{u_i\}$$

Rationale.

- Explicitly considers individuals other than worst off.

Problem.

- If u_k is smallest utility, this is simply the linear combination

$$W(\mathbf{u}) = u_k + (1 - \lambda) \sum_{i \neq k} u_i$$

- How to interpret λ ?

Utility & Fairness – Classical Methods

<i>Criterion</i>	<i>Linear?</i>	<i>Contin?</i>
Alpha fairness	yes	yes
Proportional fairness (Nash bargaining)	yes	yes
Kalai-Smorodinsky bargaining	no	yes

Alpha Fairness

$$W_{\alpha}(\mathbf{u}) = \begin{cases} \frac{1}{1-\alpha} \sum_i u_i^{1-\alpha} & \text{for } \alpha \geq 0, \alpha \neq 1 \\ \sum_i \log(u_i) & \text{for } \alpha = 1 \end{cases}$$

Mo & Walrand 2000; Verloop, Ayesta & Borst 2010

Rationale.

- Continuous and well-defined adjustment of equity/efficiency tradeoff.

Utility u_j must be reduced by $(u_j/u_i)^{\alpha}$ units to compensate for a unit increase in u_i ($< u_j$) while maintaining constant social welfare.

- Integral of power law $\sum_i u_i^{-\alpha}$
- Utilitarian when $\alpha = 0$, maximin when $\alpha \rightarrow \infty$
- Can be derived from certain axioms.

Lan & Chiang 2011

Alpha Fairness

$$W_\alpha(\mathbf{u}) = \begin{cases} \frac{1}{1-\alpha} \sum_i u_i^{1-\alpha} & \text{for } \alpha \geq 0, \alpha \neq 1 \\ \sum_i \log(u_i) & \text{for } \alpha = 1 \end{cases}$$

Model

- Nonlinear but concave.

$$\max_{\mathbf{x}, \mathbf{u}} \{W_\alpha(\mathbf{u}) \mid (\mathbf{u}, \mathbf{x}) \in S\}$$

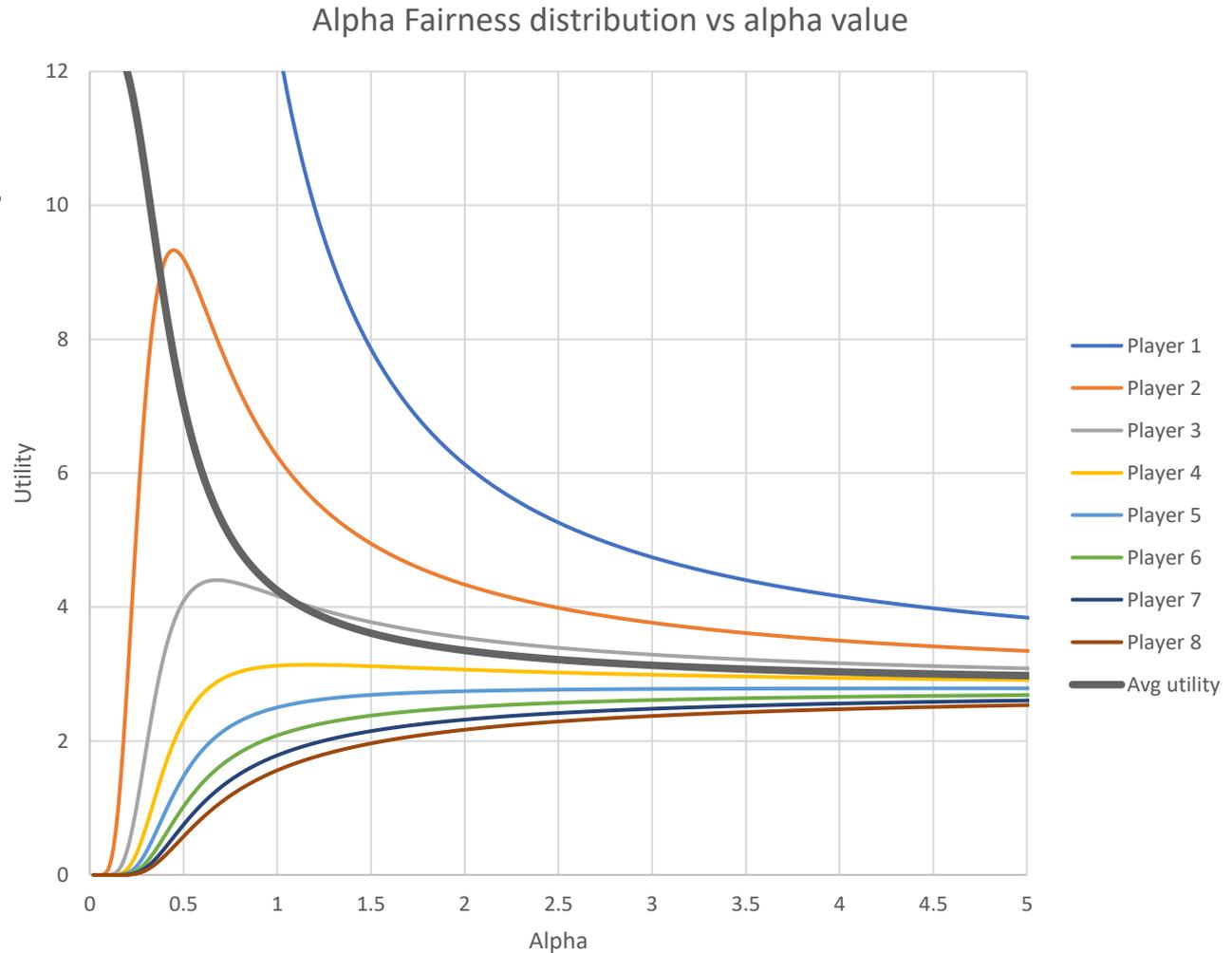
- Can be solved by efficient algorithms if constraints are linear (or perhaps if S is convex).

Alpha Fairness

Example:

Maximum alpha fairness
subject to budget constraint

$$u_1 + 2u_2 + \dots + 8u_8 \leq 100$$



Alpha Fairness

$$W_\alpha(\mathbf{u}) = \begin{cases} \frac{1}{1-\alpha} \sum_i u_i^{1-\alpha} & \text{for } \alpha \geq 0, \alpha \neq 1 \\ \sum_i \log(u_i) & \text{for } \alpha = 1 \end{cases}$$

Possible problems

- Parameter α has unobvious interpretation.
- Unclear how to choose α in practice.
- An egalitarian distribution can have same social welfare as arbitrarily extreme inequality.

In a 2-person problem, the distribution $(u_1, u_2) = (1, 1)$ has the same social welfare as $(2^{1/(1-\alpha)}, \infty)$ when $\alpha > 1$.

Proportional Fairness

$$W(\mathbf{u}) = \sum_i \log(u_i)$$

Nash 1950

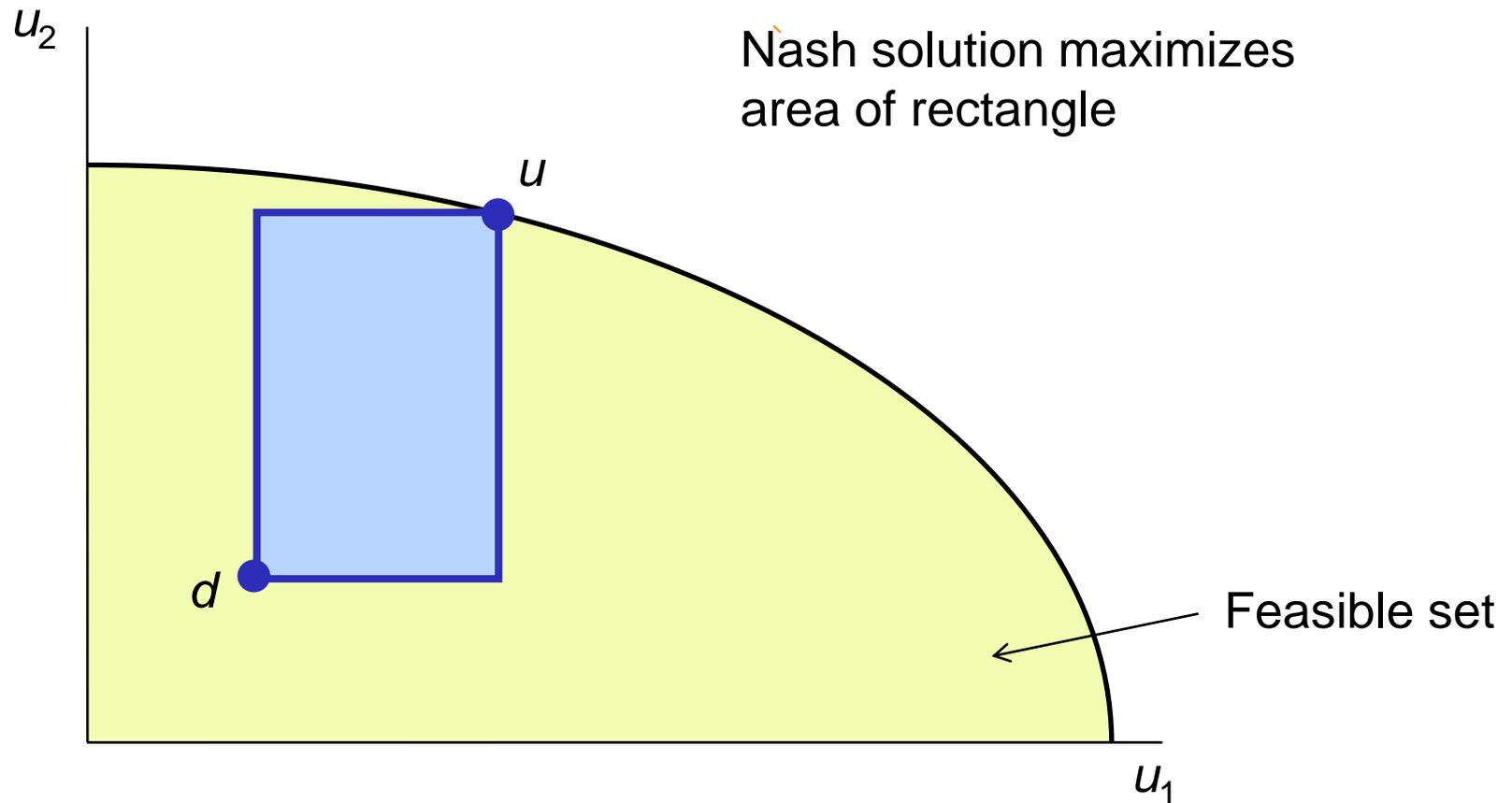
- Special case of alpha fairness ($\alpha = 1$).
- Also known as **Nash bargaining solution**, in which case bargaining starts with a default distribution \mathbf{d} .

$$W(\mathbf{u}) = \sum_i \log(u_i - d_i) \quad \text{or} \quad W(\mathbf{u}) = \prod_i (u_i - d_i)$$

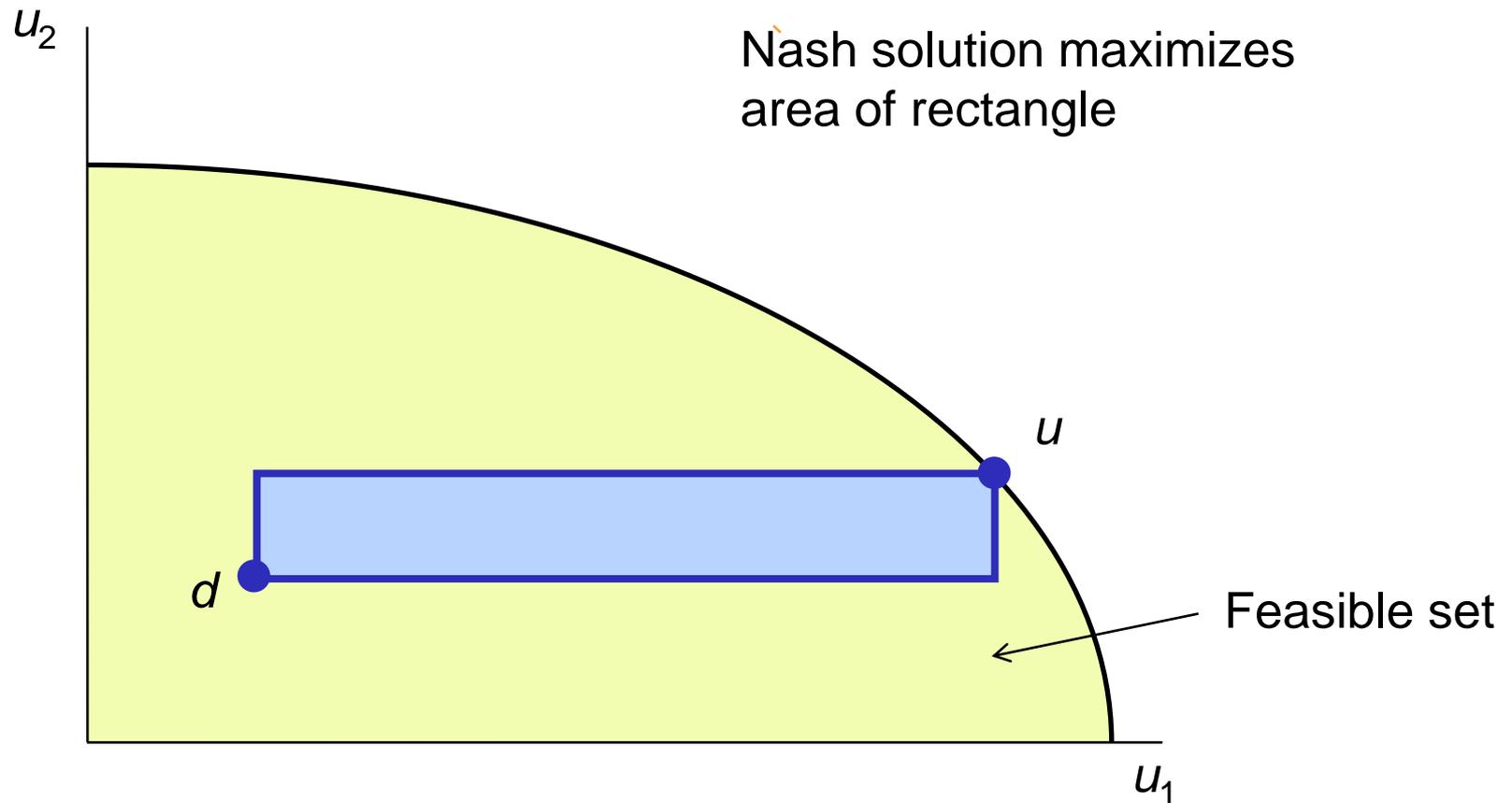
Rationale

- Has nice geometric interpretation.
- Can be derived from axiomatic and bargaining arguments.
- Used in engineering applications (telecom, traffic signaling).

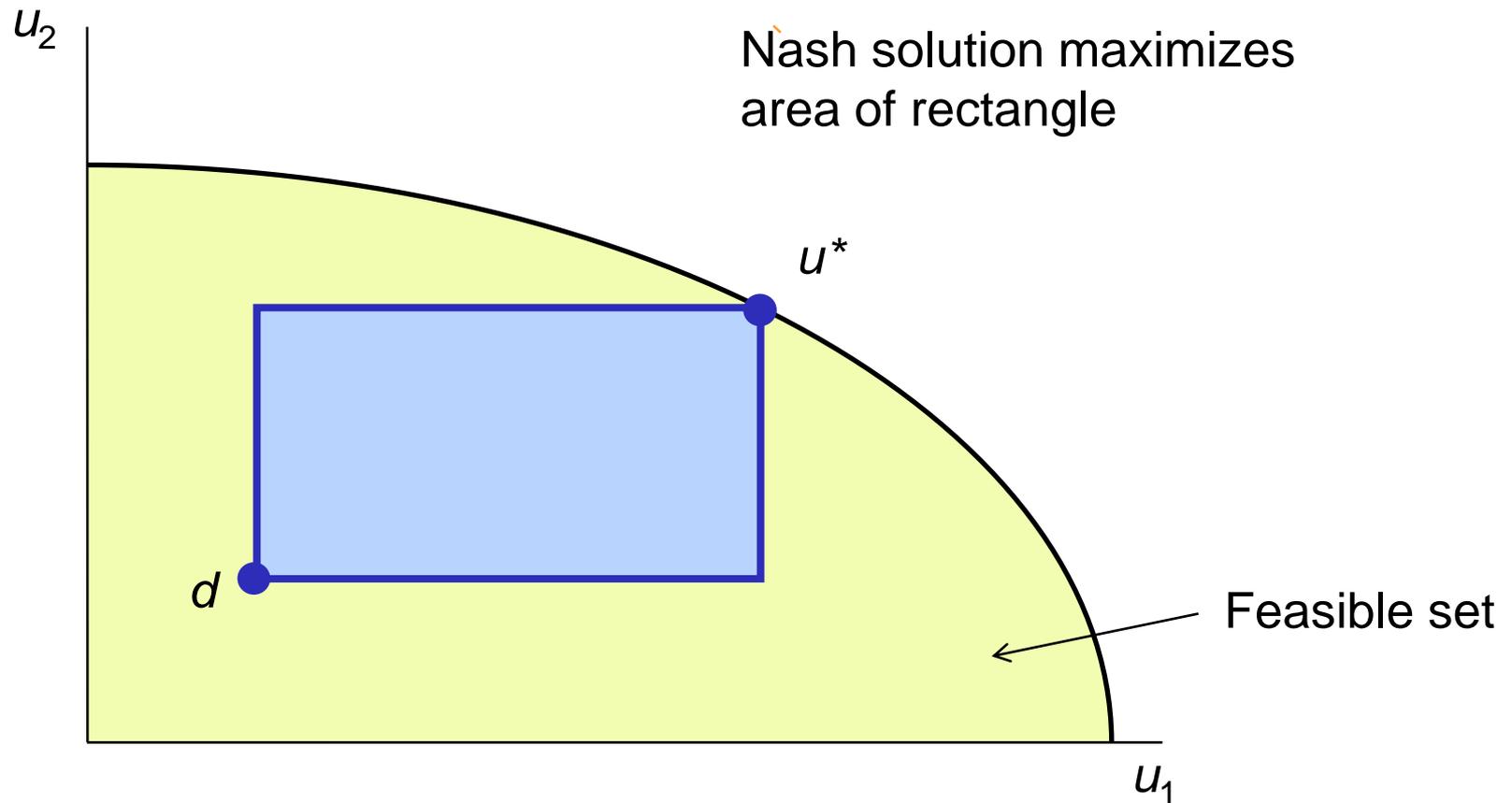
Proportional Fairness



Proportional Fairness



Proportional Fairness



Back to Social Choice Theory

Axiomatic derivation of proportional fairness

From Nash's article, based on:

- **Anonymity, Pareto** and **independence** axioms
- **Scale invariance:** invariance transformation $\phi_i(u_i) = \beta_i u_i$

Nash 1950

Back to Social Choice Theory

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- **Scale invariance:** invariance transformation $\phi_i(u_i) = \beta_i u_i$

Nash 1950

Possible problem

Invariance under individual rescaling is better suited to negotiation procedures than assessing just distributions.

Back to Social Choice Theory

Bargaining justifications

“Rational” negotiation converges to the Nash bargaining solution. Assumes an initial utility distribution to which parties return if negotiation fails.

- Finite convergence (assuming a minimum distance between offers), based on a bargaining procedure of Zeuthen.

Harsanyi 1977

Zeuthen 1930

- Asymptotic convergence based on equilibrium modeling.

Rubinstein 1982

Binmore, Rubinstein, Wolinsky 1986

Back to Social Choice Theory

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Harsanyi 1977

Zeuthen 1930

- Asymptotic convergence based on equilibrium modeling.

Rubinstein 1982

Binmore, Rubinstein, Wolinsky 1986

Possible problem

Not clear that **negotiation** leads to **justice**.

Back to Social Choice Theory

Axiomatic derivation of alpha fairness

- Certain axioms lead to a **family** of SWFs containing **alpha fairness**, along with logarithmic functions (including Theil & Atkinson indices).
- Key to the proof is an **axiom of partition**:

Lan and Chiang 2011

There exists a mean function h such that for any partition $(\mathbf{u}_1, \mathbf{u}_2)$ of \mathbf{u} and any two distributions \mathbf{u} and \mathbf{u}' ,

$$\frac{W(t\mathbf{u})}{W(t\mathbf{u}')} = h\left(\frac{W(\mathbf{u}_1)}{W(\mathbf{u}'_1)}, \frac{W(\mathbf{u}_2)}{W(\mathbf{u}'_2)}\right)$$

where $t > 0$ is an arbitrary scalar. This implies that h must be a geometric or power mean.

Back to Social Choice Theory

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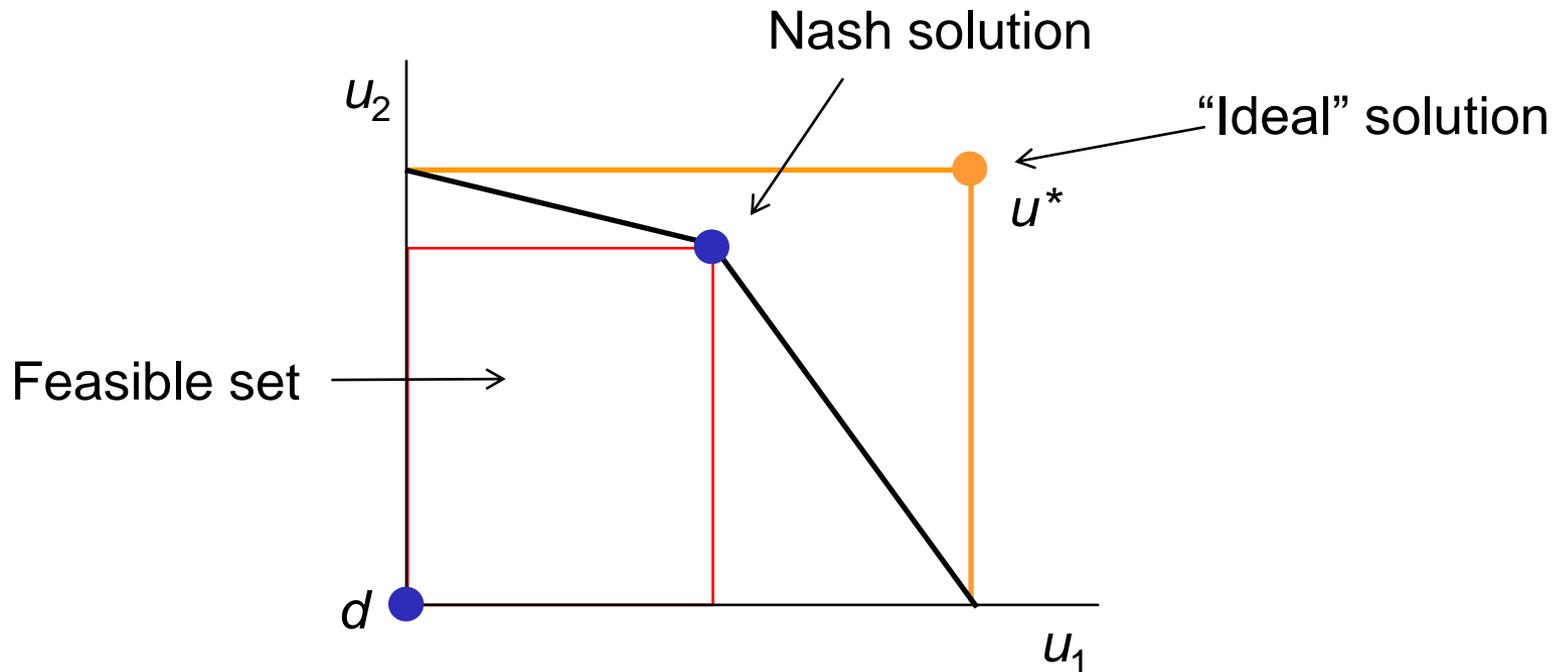
where $t > 0$ is an arbitrary scalar. This implies that h must be a geometric or power mean.

Possible problem

It is hard to interpret the axiom of partition.

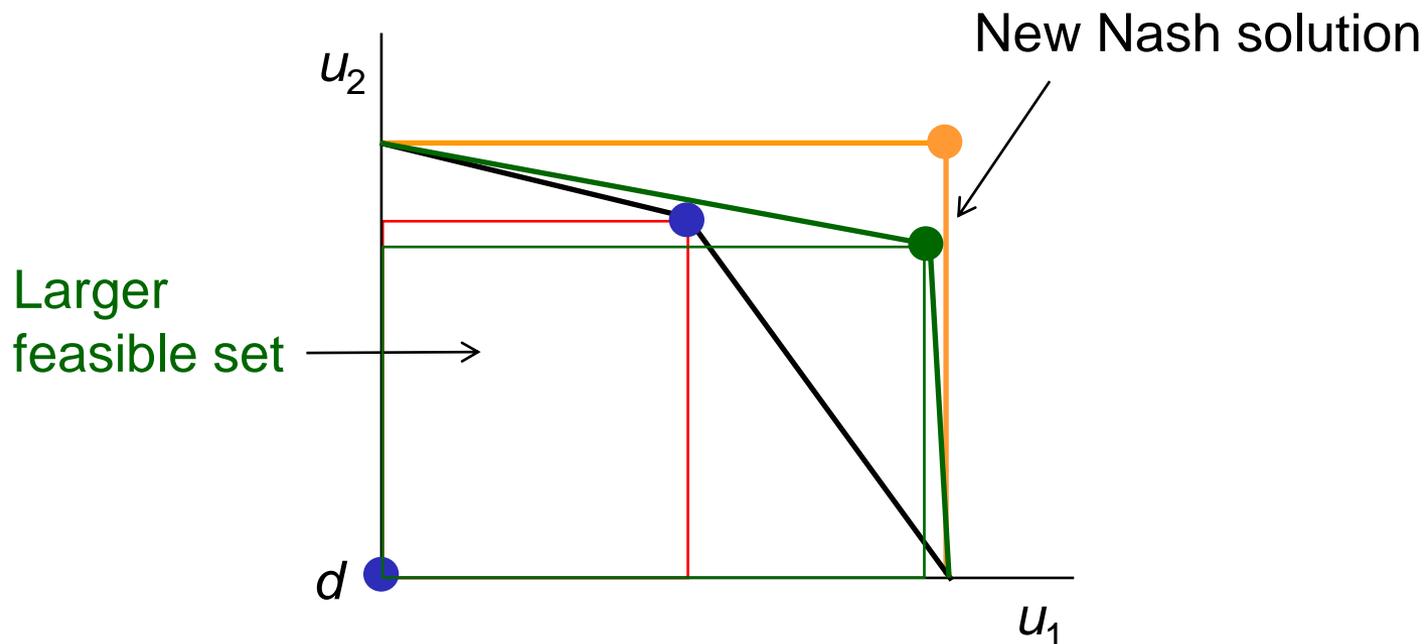
Kalai-Smorodinsky Bargaining

- Begins with a critique of the Nash bargaining solution.



Kalai-Smorodinsky Bargaining

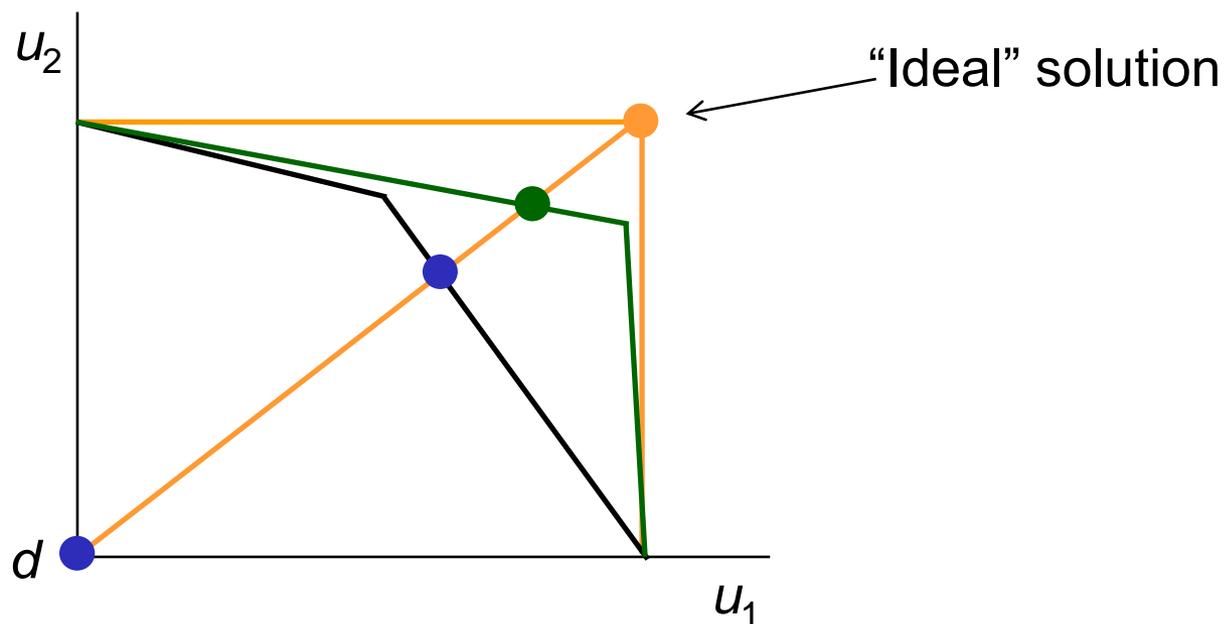
- Begins with a critique of the Nash bargaining solution.
- The new Nash solution is **worse** for player 2 even though the feasible set is **larger**.



Kalai-Smorodinsky Bargaining

- **Proposal:** Bargaining solution is pareto optimal point on line from d to ideal solution.

Kalai & Smorodinsky 1975



Kalai-Smorodinsky Bargaining

Social welfare function

$$W(\mathbf{u}) = \begin{cases} \sum_i u_i, & \text{if } \mathbf{u} = (1 - \beta)\mathbf{d} + \beta\mathbf{u}^{\max} \text{ for some } \beta \text{ with } 0 \leq \beta \leq 1 \\ 0, & \text{otherwise} \end{cases}$$

where $u_i^{\max} = \max_{\mathbf{x}, \mathbf{u}} \{u_i \mid (\mathbf{u}, \mathbf{x}) \in S\}$.

Model

$$\max_{\beta, \mathbf{x}, \mathbf{u}} \{ \beta \mid \mathbf{u} = (1 - \beta)\mathbf{d} + \beta\mathbf{u}^{\max}, (\mathbf{u}, \mathbf{x}) \in S, \beta \leq 1 \}$$

Kalai-Smorodinsky Bargaining

Rationale

- Follows from Nash's axiomatic derivation if **monotonicity replaces independence** axiom.
- Seems reasonable for **price or wage negotiation**.
- Adapts Rawlsian maximin to **relative** utility (wrt the ideal).
- Defended by some social contract theorists (e.g., "contractarians")

Gauthier 1983, Thompson 1994

Kalai-Smorodinsky Bargaining

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Gauthier 1983, Thompson 1994

Possible problem

- In some contexts, it may not be ethical to allocate utility in proportion to one's potential.
- For example, when allocating resources to those with minor ailments vs chronic diseases.

Utility & Fairness – Threshold Methods

<i>Criterion</i>	<i>Linear?</i>	<i>Contin?</i>
Utility + maximin – Utility threshold	yes	no
Utility + maximin – Equity threshold	yes	yes
Utility + leximax – Predefined priorities	yes	no
Utility + leximax – No predefined priorities	yes	no

Threshold Methods

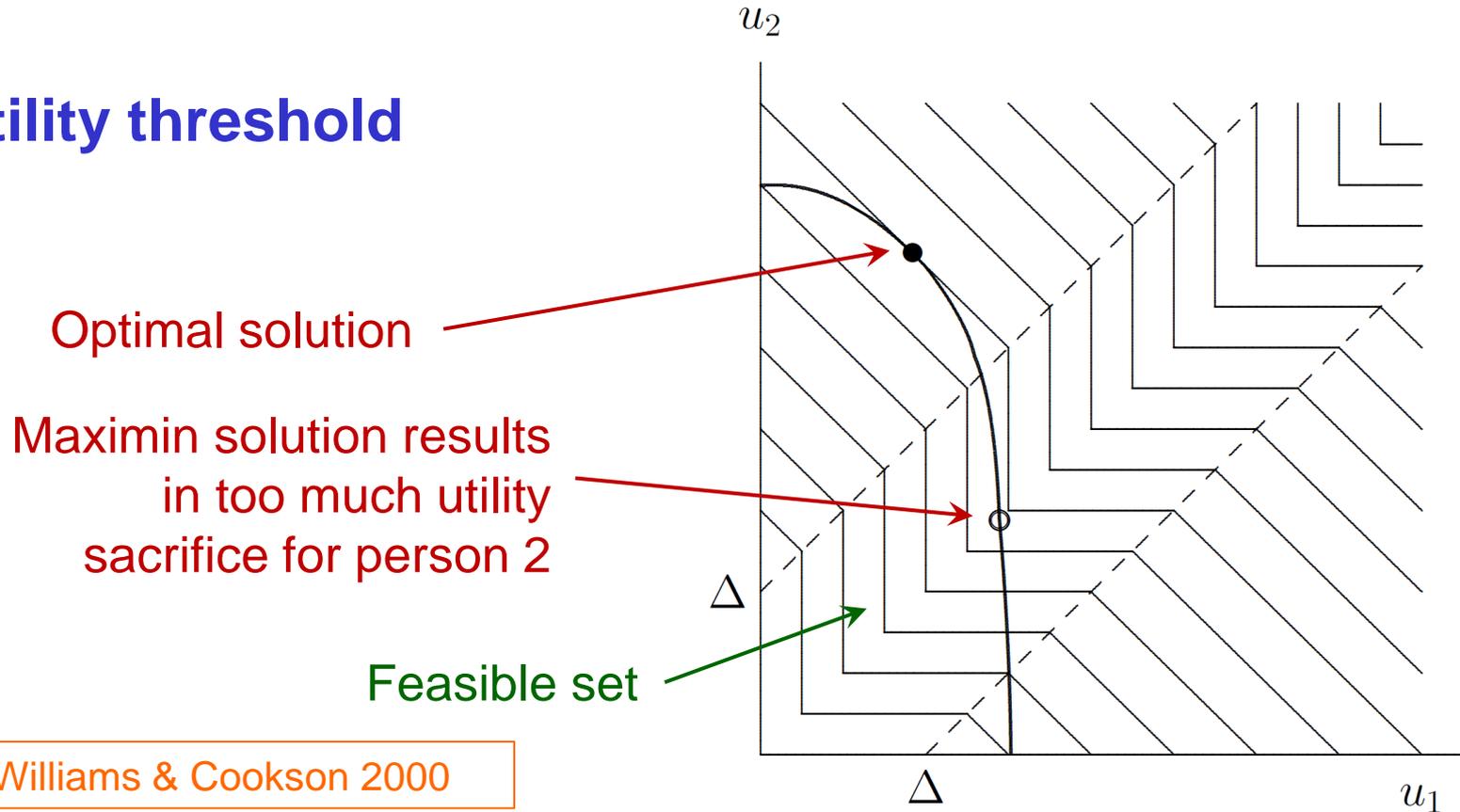
Combining utility and maximin

- **Utility threshold:** Use a maximin criterion until the utility cost becomes too great, then switch to a utilitarian criterion.
- **Equity threshold:** Use a utilitarian criterion until the inequity becomes too great, then switch to a maximin criterion.

Williams & Cookson 2000

Threshold Methods

Utility threshold



Williams & Cookson 2000

$$W(u_1, u_2) = \begin{cases} u_1 + u_2, & \text{if } |u_1 - u_2| \geq \Delta \\ 2 \min\{u_1, u_2\} + \Delta, & \text{otherwise} \end{cases}$$

Threshold Methods

Utility threshold

Generalization to n persons

$$W(\mathbf{u}) = (n - 1)\Delta + \sum_{i=1}^n \max \{u_i - \Delta, u_{\min}\}$$

where $u_{\min} = \min_i \{u_i\}$

JH & Williams 2012

Rationale

- Utilities within Δ of the lowest are in the **fair region**.
- Trade-off parameter Δ has a **practical interpretation**.
- Δ is chosen so that individuals in fair region are sufficiently deprived to **deserve priority**.
- Suitable when **equity** is the initial concern, but without paying **too high a cost** for fairness (healthcare, politically sensitive contexts).
- $\Delta = 0$ corresponds to utilitarian criterion, $\Delta = \infty$ to maximin.

Threshold Methods

Utility threshold

Model

$$\max_{\mathbf{x}, \mathbf{u}, \boldsymbol{\delta}, \mathbf{v}, w, z} \left\{ n\Delta + \sum_i v_i \right. \left. \begin{array}{l} u_i - \Delta \leq v_i \leq u_i - \Delta\delta_i, \text{ all } i \\ w \leq v_i \leq w + (M - \Delta)\delta_i, \text{ all } i \\ u_i - u_j \leq M, \text{ all } i, j \\ u_i \geq 0, \delta_i \in \{0, 1\}, \text{ all } i \\ (\mathbf{u}, \mathbf{x}) \in S \end{array} \right\}$$

- Tractable MILP model.
- Model is **sharp** without $(\mathbf{u}, \mathbf{x}) \in S$.
- Easily generalized to differently-sized **groups** of individuals.

JH & Williams 2012

Possible problem

- Due to maximin component, many solutions with different equity properties have same social welfare value.

Threshold Methods

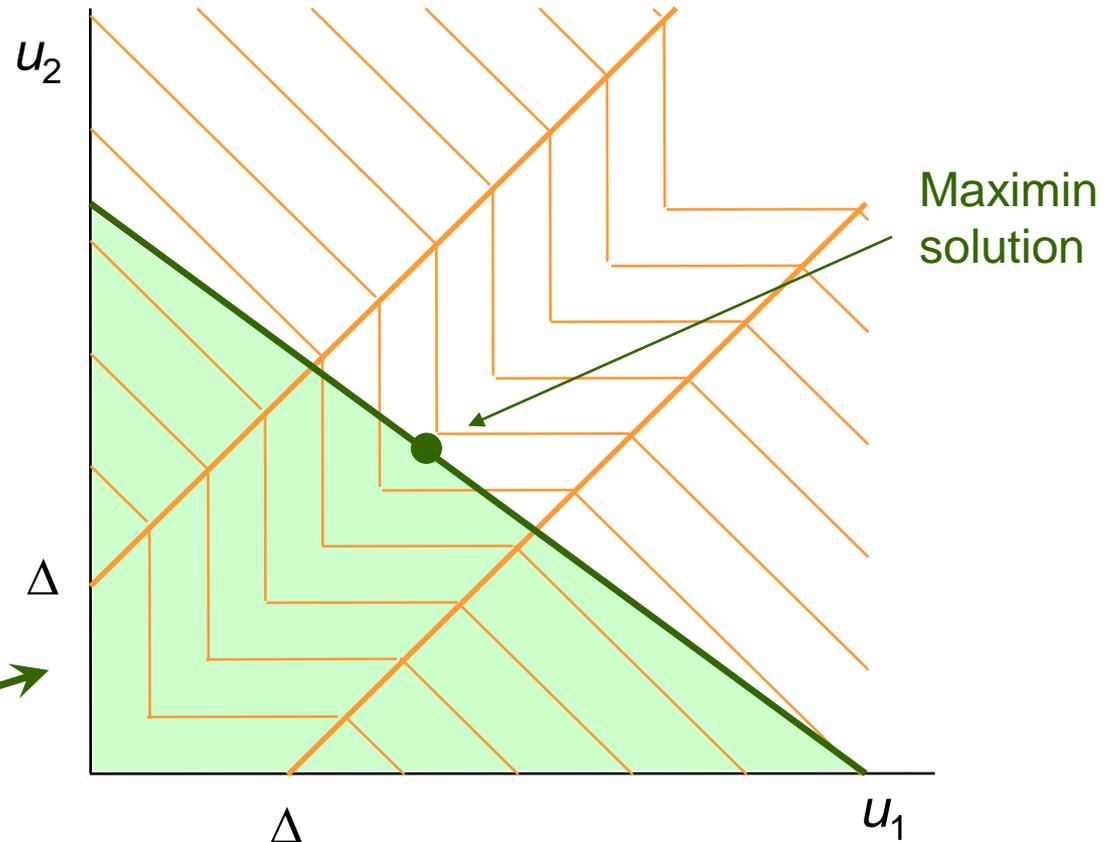
Utility threshold

Theorem. When maximizing the SWF subject to a **budget constraint**, the optimal solution is purely **maximin** or **purely utilitarian**.

Purely maximin if

$$\Delta \geq B \left(\frac{1}{a_{\langle 1 \rangle}} - \frac{n}{\sum_i a_i} \right) \Delta$$

Here, patients have **similar** treatment costs, or Δ is **large**.



Elçi, JH, and Zhang 2022

Threshold Methods

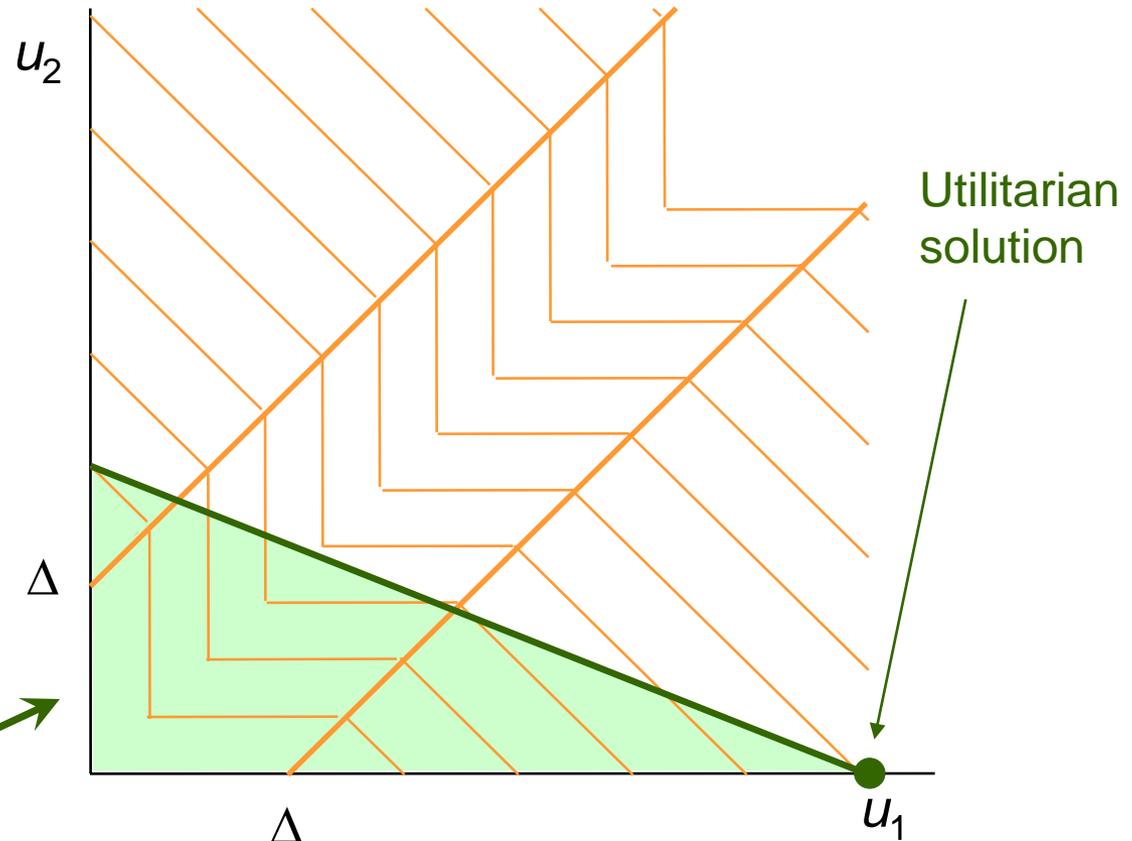
Utility threshold

Theorem. When maximizing the SWF subject to a **budget constraint**, the optimal solution is purely **maximin** or **purely utilitarian**.

Purely utilitarian if

$$\Delta \leq B \left(\frac{1}{a_{\langle 1 \rangle}} - \frac{n}{\sum_i a_i} \right) \Delta$$

Here, patients have **very different** treatment costs, or Δ is **small**.



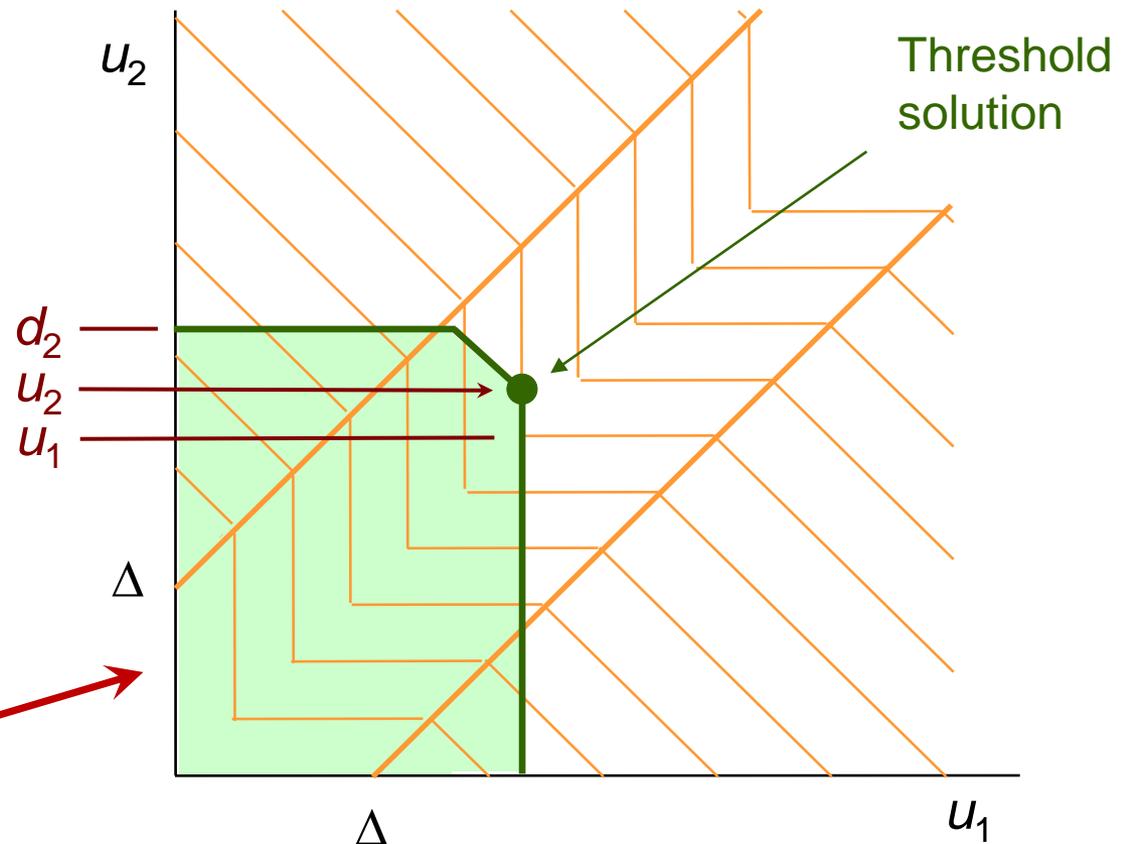
Elçi, JH, and Zhang 2022

Threshold Methods

Utility threshold

Theorem. When maximizing the SWF subject to a **budget constraint** and **upper bounds d_i** at most one utility is **strictly between** its upper bound and the smallest utility.

Here, **one** utility u_2 is **strictly between** upper bound d_2 and the smallest utility u_1 .



Elçi, JH, and Zhang 2022

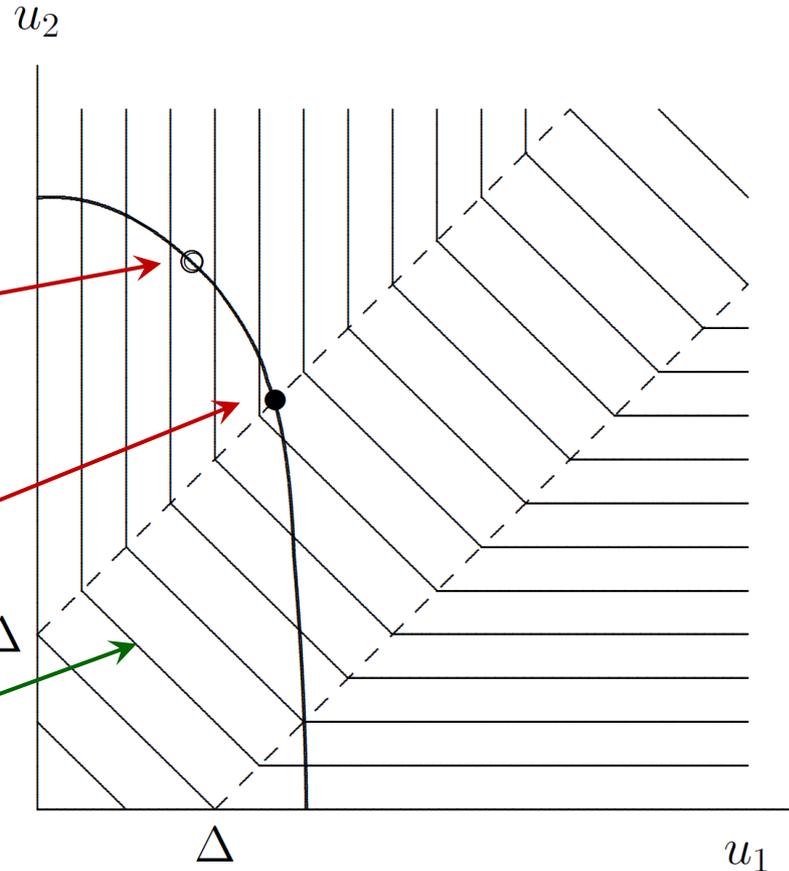
Threshold Methods

Equity threshold

Utilitarian solution
leaves person 1
overly deprived

Optimal solution

Feasible set



Williams & Cookson 2000

$$W(u_1, u_2) = \begin{cases} 2 \min\{u_1, u_2\} + \Delta, & \text{if } |u_1 - u_2| \geq \Delta \\ u_1 + u_2, & \text{otherwise} \end{cases}$$

Threshold Methods

Equity threshold

Generalization to n persons

$$W(\mathbf{u}) = n\Delta + \sum_{i=1}^n \min\{u_i - \Delta, u_{min}\}$$

Elçi, JH, and Zhang 2023

Rationale

- Utilities more than Δ above the lowest are in the **fair region**.
- Trade-off parameter Δ has a **practical interpretation**.
- Δ is chosen so that well-off individuals (those in fair region) **do not deserve more utility** unless smaller utilities are also increased.
- Suitable when **efficiency** is the initial concern, but one does not want to create **excessive inequality** (traffic management, telecom, disaster recovery).

Threshold Methods

Equity threshold

Model

$$\max_{\mathbf{x}, \mathbf{u}, \mathbf{v}, w, z} \left\{ n\Delta + \sum_i v_i \mid \begin{array}{l} v_i \leq w \leq u_i, \text{ all } i \\ v_i \leq u_i - \Delta, \text{ all } i \\ w \geq 0, v_i \geq 0, \text{ all } i \\ (\mathbf{u}, \mathbf{x}) \in S \end{array} \right\}$$

Elçi, JH, and Zhang 2023

- Linear model.
- Easily generalized to differently-sized **groups** of individuals.

Possible problem

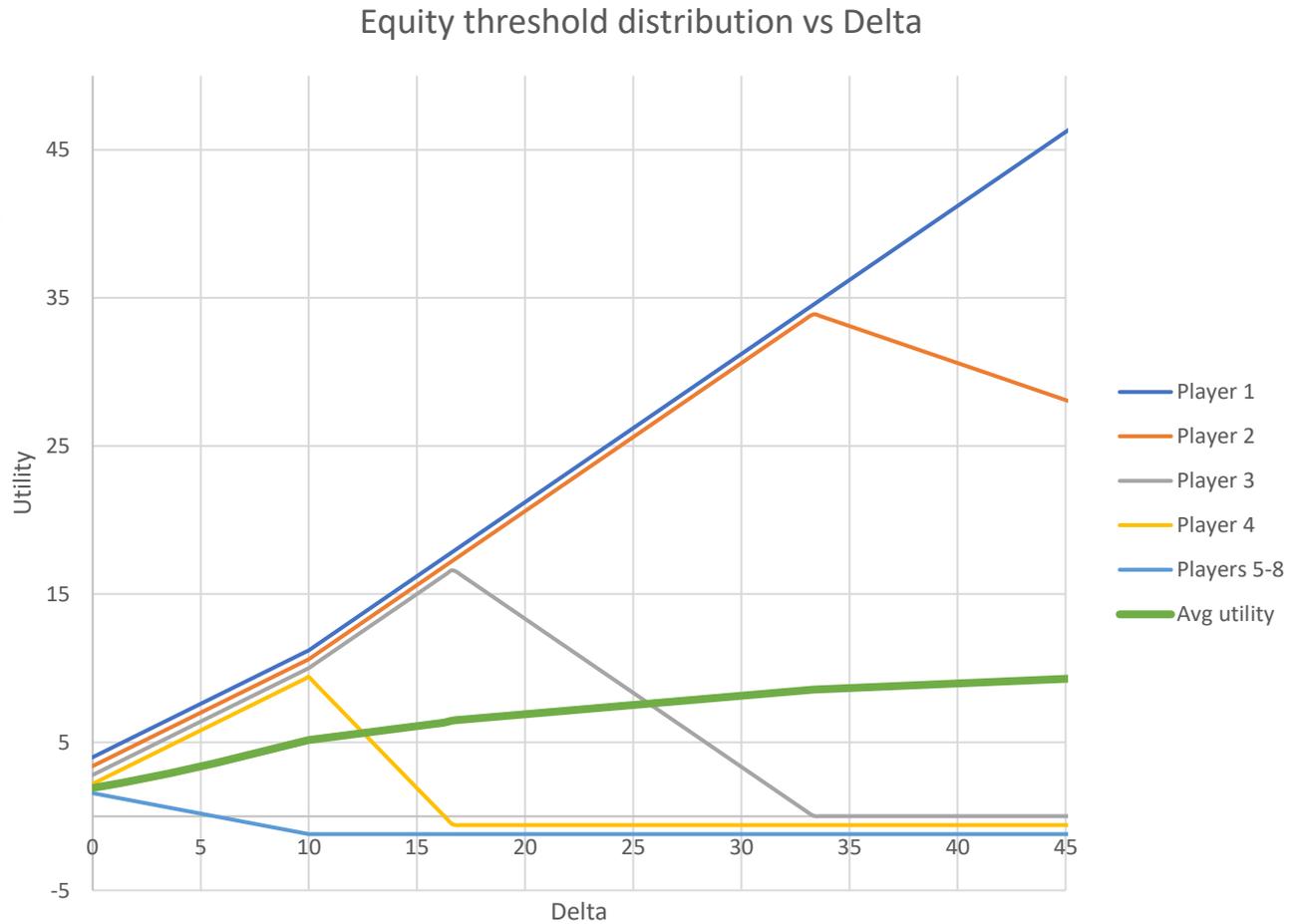
- As with threshold model, many solutions with different equity properties have same social welfare value.

Threshold Methods

Example:

Maximum equity
threshold SWF
subject to budget constraint

$$u_1 + 2u_2 + \dots + 8u_8 \leq 100$$



Threshold Methods

Utility + leximax, predetermined preferences

$$W(\mathbf{u}) = \begin{cases} nu_1, & \text{if } |u_i - u_j| \leq \Delta \text{ for all } i, j \\ \sum_i u_i + \text{sgn}(u_1 - u_i)\Delta, & \text{otherwise} \end{cases}$$

where preference order is u_1, \dots, u_n .

McElfresh & Dickerson 2018

Rationale

- Takes into account utility levels of individuals in the fair region.
- Successfully applied to kidney exchange.

Threshold Methods

Utility + leximax, predetermined preferences

Model (MILP)

$$\max_{\substack{\mathbf{u}, \mathbf{x} \\ w_1, w_2 \\ \mathbf{y}, \phi, \delta}} \left\{ w_1 + w_2 \left| \begin{array}{l} w_1 \leq nu_1, w_1 \leq M\phi \\ w_2 \leq \sum_i (u_i + y_i), w_2 \leq M(1 - \phi) \\ u_i - u_j - \Delta \leq M(1 - \phi), \text{ all } i, j \\ y_i \leq \Delta, y_i \leq -\Delta + M\delta_i, u_i - u_1 \leq M(1 - \delta_i), \text{ all } i \\ (\mathbf{u}, \mathbf{x}) \in S; \phi, \delta_i \in \{0, 1\}, \text{ all } i \end{array} \right. \right\}$$

where preference order is u_1, \dots, u_n .

Threshold Methods

Utility + leximax, predetermined preferences

Possible problems

- SWF is discontinuous.
- Preferences cannot be pre-ordered in many applications.
- Leximax is not incorporated in the SWF, but is applied only to SWF-maximizing solutions.

Threshold Methods

Utility + leximax, sequence of SWFs

SWFs W_1, \dots, W_n are maximized sequentially, where W_1 is the utility threshold SWF defined earlier, and W_k for $k \geq 2$ is

$$W_k(\mathbf{u}) = \sum_{i=1}^{k-1} (n - i + 1)u_{\langle i \rangle} + (n - k + 1) \min \{u_{\langle 1 \rangle} + \Delta, u_{\langle k \rangle}\} + \sum_{i=k}^n \max \{0, u_{\langle i \rangle} - u_{\langle 1 \rangle} - \Delta\}$$

where $u_{\langle 1 \rangle}, \dots, u_{\langle n \rangle}$ are u_1, \dots, u_n in nondecreasing order.

Rationale

Chen & JH 2021

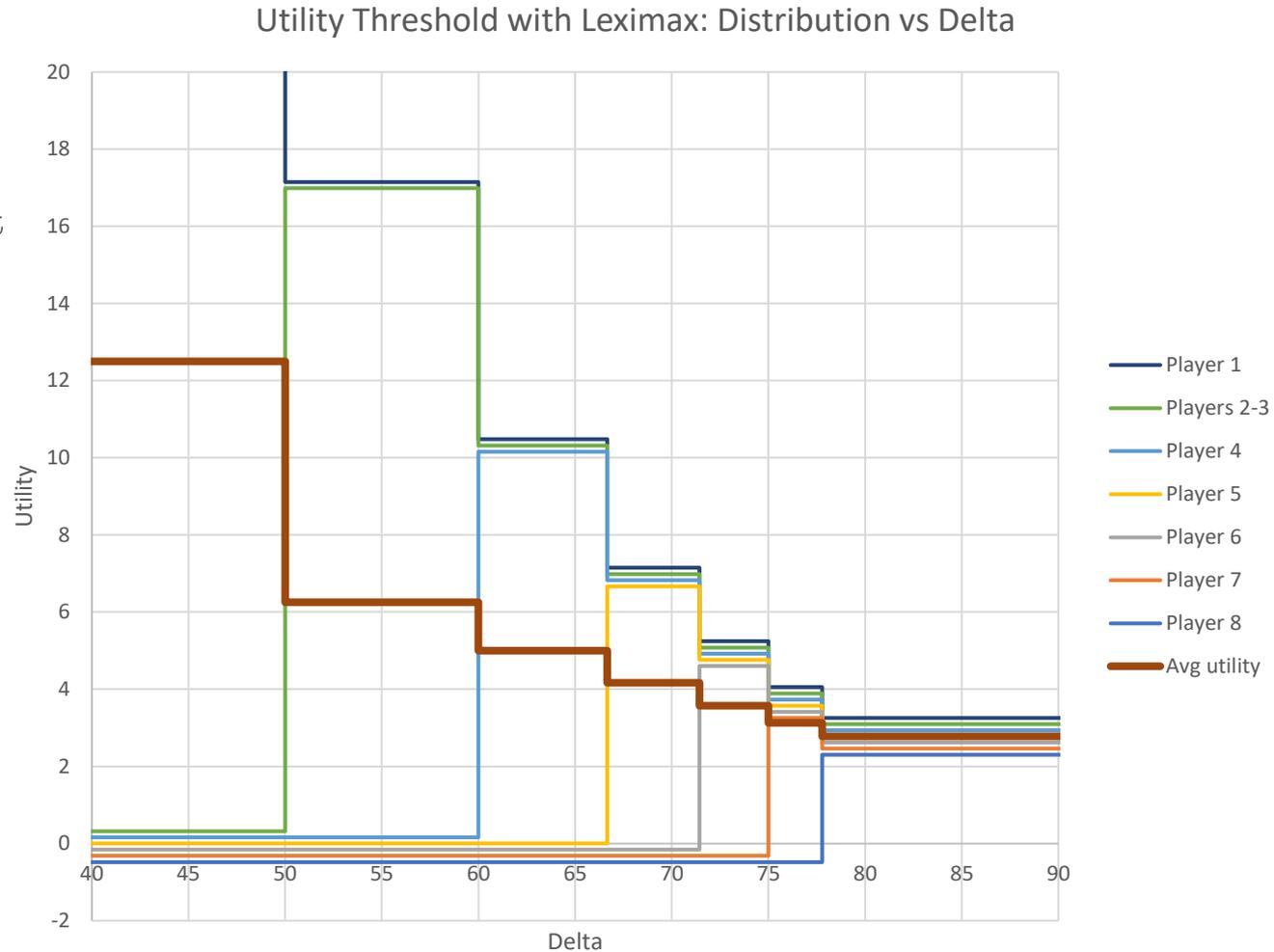
- Does not require pre-ordered preferences.
- Takes into account utility levels of all individuals in the fair region.
- Tractable MILP models in practice, valid inequalities known.

Threshold Methods

Example:

Maximum utility threshold
SWF with leximax
subject to budget constraint

$$u_1 + 2u_2 + \dots + 8u_8 \leq 100$$



Threshold Methods

Possible problems

- Requires solving a sequence of MILPs.
- Hard to explain and justify on first principles.

Threshold Methods

Utility + leximax, sequence of SWFs

Model (MILP for W_k)

$$\max_{\substack{\mathbf{x}, \mathbf{u}, \delta, \epsilon \\ \mathbf{v}, w, \sigma, z}} \left\{ z \left[\begin{array}{l} z \leq (n - k + 1)\sigma + \sum_{i \in I_k} v_i \\ 0 \leq v_i \leq M\delta_i, \quad i \in I_k \\ v_i \leq u_i - \bar{u}_{i_1} - \Delta + M(1 - \delta_i), \quad i \in I_k \\ \sigma \leq \bar{u}_{i_1} + \Delta \\ \sigma \leq w \\ w \leq u_i, \quad i \in I_k \\ u_i \leq w + M(1 - \epsilon_i), \quad i \in I_k \\ \sum_{i \in I_k} \epsilon_i = 1 \\ w \geq \bar{u}_{i_{k-1}} \\ u_i - \bar{u}_{i_1} \leq M, \quad i \in I_k \\ \delta_i, \epsilon_i \in \{0, 1\}, \quad i \in I_k \end{array} \right. \right.$$

where \bar{u}_{i_k} is the value of the smallest utility in the optimal solution of the k th MILP model, and $I = \{1, \dots, n\} \setminus \{i_1, \dots, i_{k-1}\}$. The socially optimal solution is $(\bar{u}_1, \dots, \bar{u}_n)$.

Threshold Methods – Healthcare Example

- Based on budget decisions in UK National Health Service
- Allocate limited treatment resources to disease/prognosis categories of patients.
- Based on cost, number of patients, and QALY estimates with and without treatment.*
- We will compare **2 utility-threshold SWFs**: utility + maximin and sequential utility + leximax.
- Solution time = fraction of second for each value of Δ .

Problem due to JH & Williams 2012

*QALY = quality adjusted life-year. Data reflect a particular situation and are not valid in general. Solutions presented here should not be taken as a general recommendation for healthcare resource allocation, but only as an illustration of social welfare functions.

QALY
& cost
data

Part 1

Intervention	Cost per person c_i (£)	QALYs gained q_i	Cost per QALY (£)	QALYs without intervention α_i	Subgroup size n_i
<i>Pacemaker for atrioventricular heart block</i>					
Subgroup A	3500	3	1167	13	35
Subgroup B	3500	5	700	10	45
Subgroup C	3500	10	350	5	35
<i>Hip replacement</i>					
Subgroup A	3000	2	1500	3	45
Subgroup B	3000	4	750	4	45
Subgroup C	3000	8	375	5	45
<i>Valve replacement for aortic stenosis</i>					
Subgroup A	4500	3	1500	2.5	20
Subgroup B	4500	5	900	3	20
Subgroup C	4500	10	450	3.5	20
<i>CABG¹ for left main disease</i>					
Mild angina	3000	1.25	2400	4.75	50
Moderate angina	3000	2.25	1333	3.75	55
Severe angina	3000	2.75	1091	3.25	60
<i>CABG for triple vessel disease</i>					
Mild angina	3000	0.5	6000	5.5	50
Moderate angina	3000	1.25	2400	4.75	55
Severe angina	3000	2.25	1333	3.75	60
<i>CABG for double vessel disease</i>					
Mild angina	3000	0.25	12,000	5.75	60
Moderate angina	3000	0.75	4000	5.25	65
Severe angina	3000	1.25	2400	4.75	70

QALY
& cost
data

Part 2

Intervention	Cost per person c_i (£)	QALYs gained q_i	Cost per QALY (£)	QALYs without intervention α_i	Subgroup size n_i
<i>Heart transplant</i>	22,500	4.5	5000	1.1	2
<i>Kidney transplant</i>					
Subgroup A	15,000	4	3750	1	8
Subgroup B	15,000	6	2500	1	8
<i>Kidney dialysis</i>					
<i>Less than 1 year survival</i>					
Subgroup A	5000	0.1	50,000	0.3	8
<i>1-2 years survival</i>					
Subgroup B	12,000	0.4	30,000	0.6	6
<i>2-5 years survival</i>					
Subgroup C	20,000	1.2	16,667	0.5	4
Subgroup D	28,000	1.7	16,471	0.7	4
Subgroup E	36,000	2.3	15,652	0.8	4
<i>5-10 years survival</i>					
Subgroup F	46,000	3.3	13,939	0.6	3
Subgroup G	56,000	3.9	14,359	0.8	2
Subgroup H	66,000	4.7	14,043	0.9	2
Subgroup I	77,000	5.4	14,259	1.1	2
<i>At least 10 years survival</i>					
Subgroup J	88,000	6.5	13,538	0.9	2
Subgroup K	100,000	7.4	13,514	1.0	1
Subgroup L	111,000	8.2	13,537	1.2	1

Threshold Methods – Healthcare Example

Budget constraint

$$\sum_j n_j c_j y_j \leq B$$

Size of treatment group j (points to n_j)
 Unit cost of treatment j (points to c_j)
 Fraction of group treated (points to y_j)

Utility function

$$u_i = q_i y_i + \alpha_i$$

Treatment benefit (QALYs) (points to $q_i y_i$)
 QALYs without treatment (points to α_i)

which implies $y_i = (u_i - \alpha_i) / q_i$

So the optimization problem becomes

$$\max_{\mathbf{u}} \left\{ W(\mathbf{u}) \mid \sum_j \frac{n_j c_j}{q_j} u_j \leq B + \sum_j \frac{n_j c_j \alpha_j}{q_j}; \quad \alpha \leq \mathbf{u} \leq \mathbf{q} + \alpha \right\}$$

Utility + maximin

Δ (QALYs)

Budget = £3 million

0 3.4 4.5 5.5 13.2 15.5

Pacemaker

Hip replace

Aortic valve

2 vessel

3 vessel

Left main

Heart transplant

Kidney transpl.

>10 yr life exp.

5-10 yr

2-5 yr

1-2 yr

<1 yr

Increasing severity →

Avg. utility (QALYs)

7.54

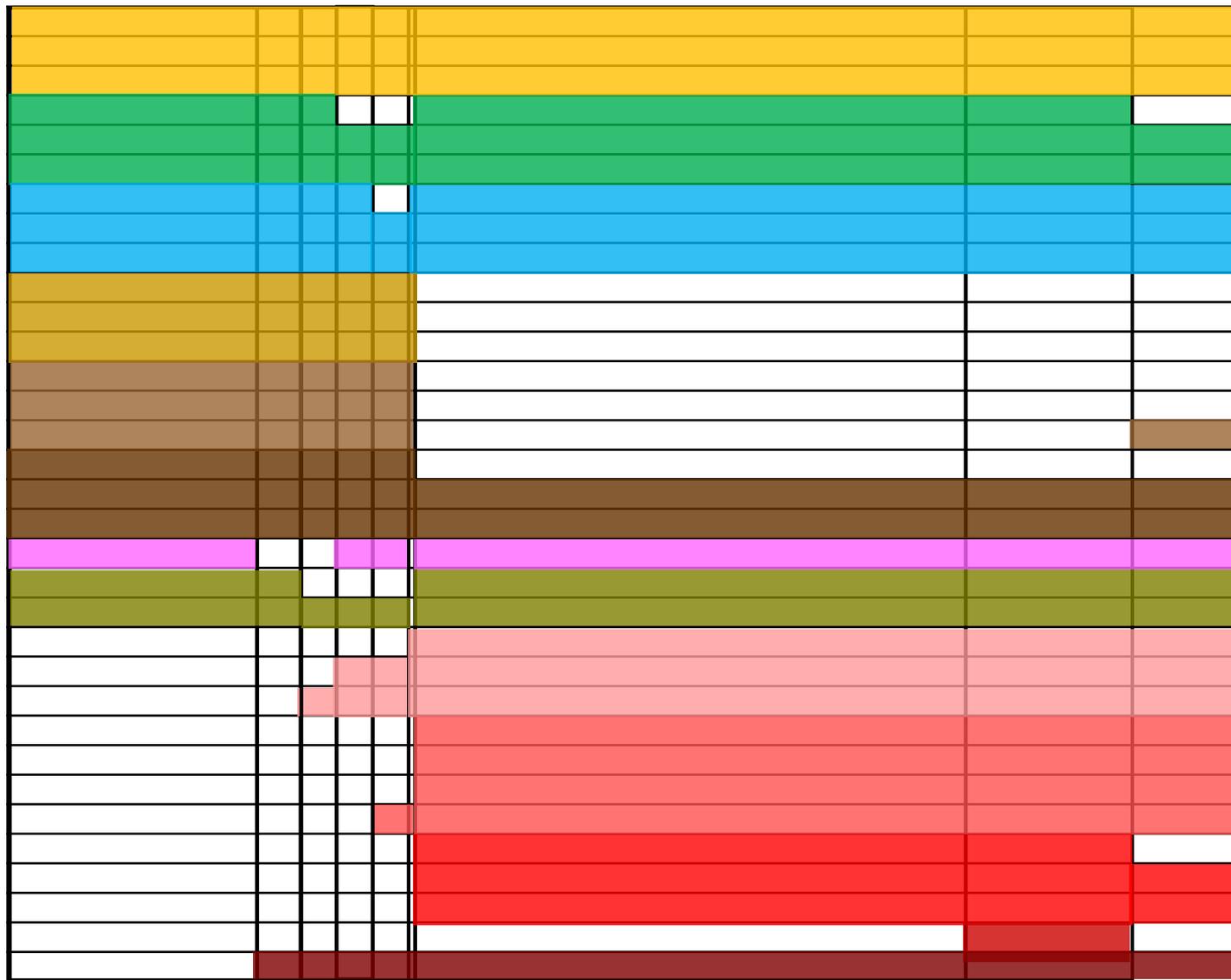
7.43

7.36

7.03

7.19

93



Utility + leximax

Δ (QALYs)

Budget = £3 million

0 1 2 3.4 5.4 6.6 8.4 11.6 13.1

Pacemaker

Hip replace

Aortic valve

2 vessel

3 vessel

Left main

Heart transplant

Kidney transpl.

>10 yr life exp.

5-10 yr

2-5 yr

1-2 yr

<1 yr

Increasing severity →

Avg. utility

7.54

7.21

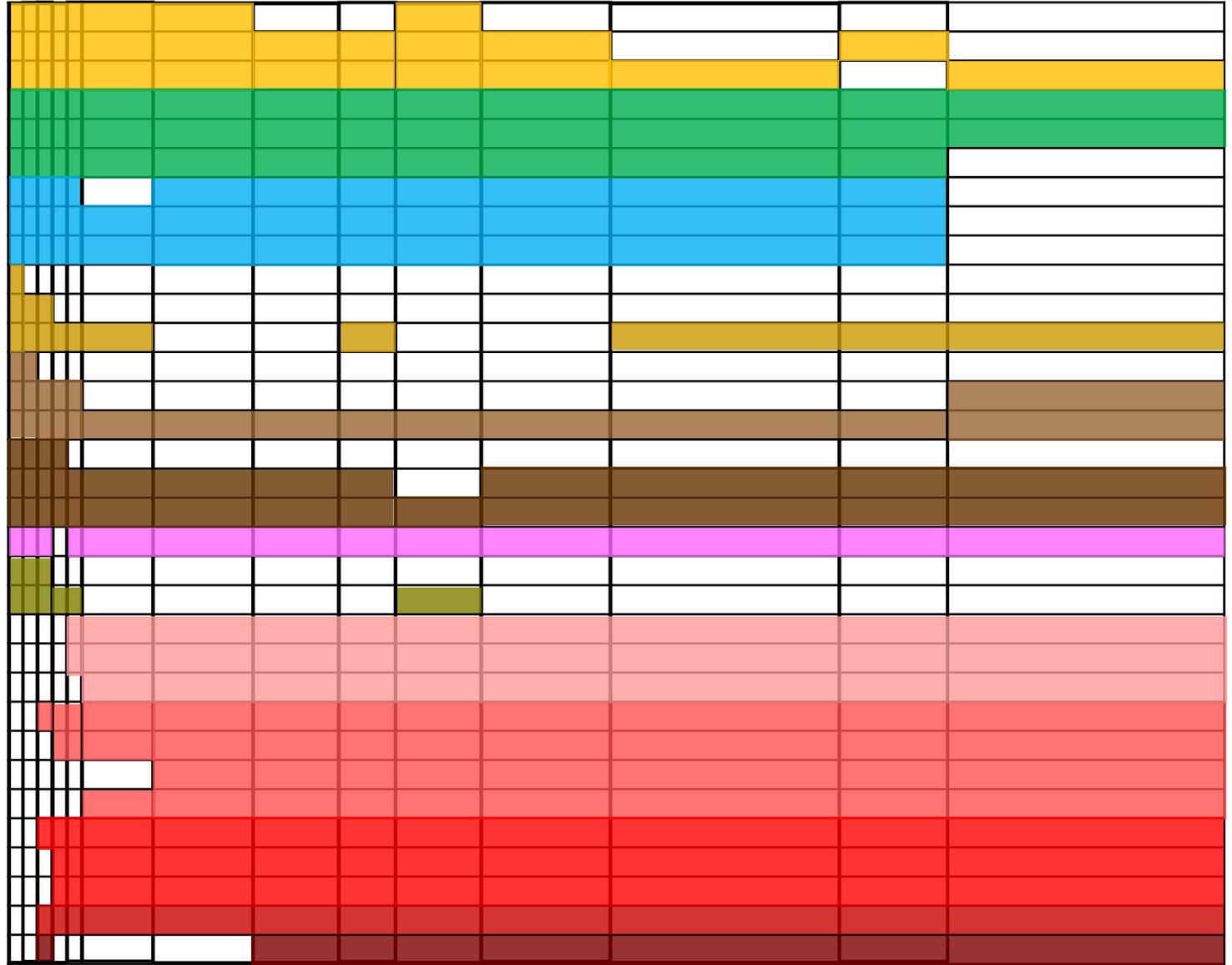
7.12

6.94

6.8

6.41

94



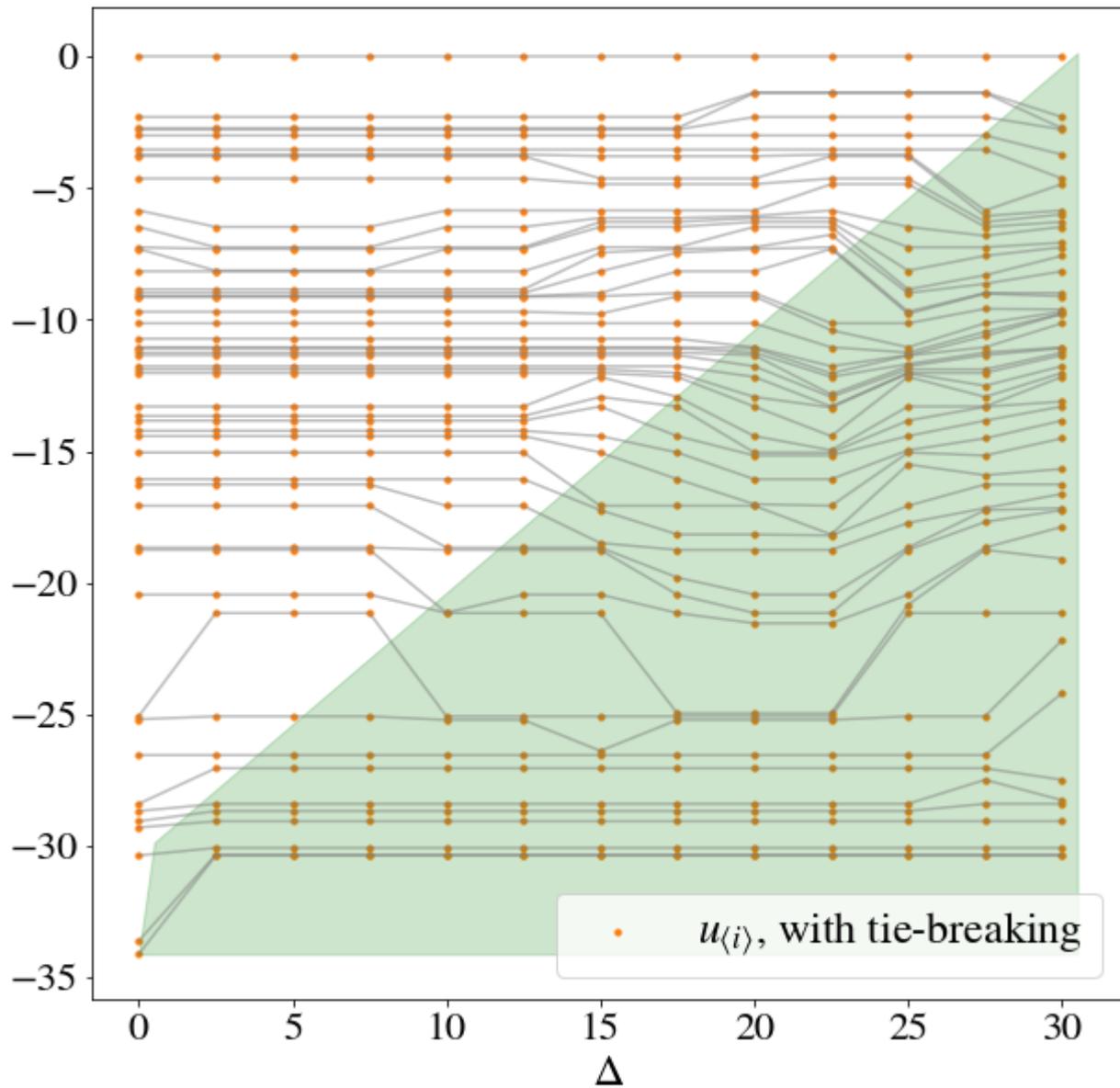
Threshold Methods – Disaster Preparedness Example

- Select earthquake shelter locations.
- Utility = negative distance of each neighborhood to nearest shelter, subject to limited budget.
- We will compare **2 utility-threshold SWFs**: utility + maximin and sequential utility + leximax.
- 50 neighborhoods, 50 potential shelter locations.
- Solution time = 1 to 18 seconds for each value of Δ .

Problem due to Mostajabdaveh, Gutjahr & Salman 2019

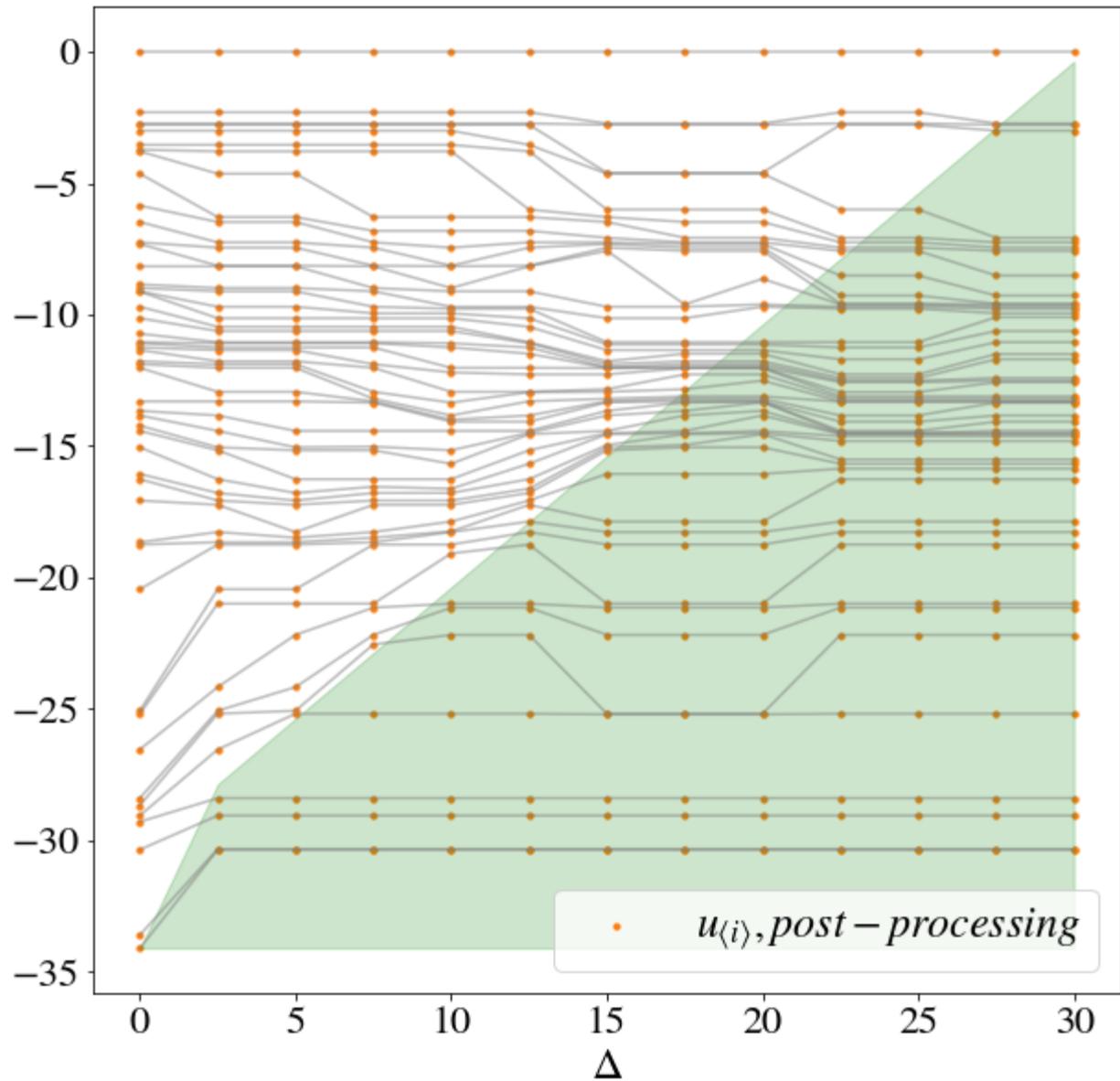
Threshold
SWF

Utility +
maximin



Threshold
SWF

Utility +
leximax



Questions? Comments?

