Optimization Models for Equity

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Modeling Equity

- There is a growing interest in incorporating equity considerations in mathematical programming models.
 - Not enough to minimize cost or maximize revenue.
 - Also concerned about distribution of resources/benefits.
 - Not obvious how to capture equity in the objective function.
 - Still less obvious how to combine it with an efficiency objective.

Modeling Equity

- Some applications...
 - Health care allocation.
 - Facility location (e.g., emergency services).
 - Taxation (revenue vs. progressivity).
 - Relief operations and disaster planning/response.
 - Telecommunications (lexmax, Nash bargaining solution)







Availability

- These slides are available on my website.
 - Google "John Hooker"
 - Originally presented at a workshop at London School of Economics, December 2010.
- Will give only a brief overview today.

Summary Outline

- Fairness and equality:
 - Utilitarianism
 - Piecewise Linear Modeling
 - Rawlsian Difference Principle
 - Axiomatics
 - Measures of Inequality
 - An Allocation Problem
- Bargaining and equality/efficiency combinations:
 - Nash Bargaining Solution
 - Raiffa-Kalai-Smorodinsky Bargaining
 - Disjunctive Modeling
 - Combining Equity and Efficiency
 - Health Care Example

Fairness and Equality Outline

- Utilitarianism
 - Utility and production functions
 - The optimization problem
 - Arguments for utilitarianism
- Piecewise Linear Modeling
 - LP model of concave maximization
 - MILP model of nonconcave maximization

Fairness and Equality Outline

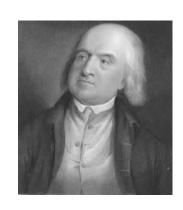
- Rawlsian Difference Principle
 - The social contract argument
 - The lexmax principle
 - The optimization problem
- Axiomatics
 - Interpersonal comparability
 - Axioms of rational choice
 - Social welfare functions

Fairness and Equality Outine

- Measures of Inequality
 - An example
 - Utrilitarian, maximin, and lexmax solution
 - Relative range, max, min
 - Relative mean deviation
 - Variance, coefficient of variation
 - McLoone index
 - Gini coefficient
 - Atkinson index
 - Hoover index
 - Theil index
- An Allocation Problem

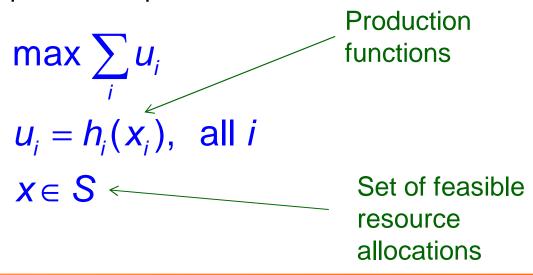
Efficiency vs. Equity

- Two classical criteria for distributive justice:
 - Utilitarianism (efficiency)
 - Difference principle of John Rawls (equity)
- These have the must studied philosophical underpinnings.

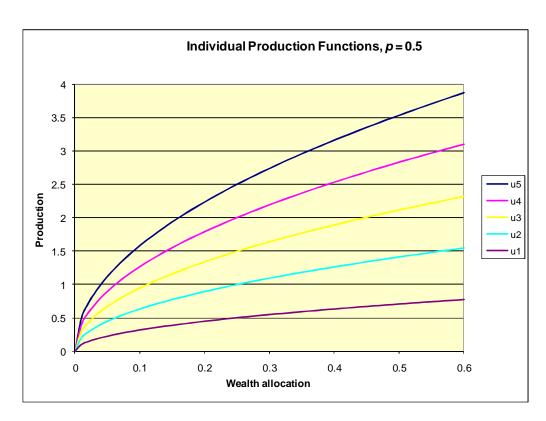




- Utilitarianism seeks allocation of resources that maximizes total utility.
 - Let x_i = resources allocated to person i.
 - Let u_i = utility enjoyed by person i.
 - We have an optimization problem



For example, $h_i(x_i) = a_i x_i^p$ with different a_i s for 5 individuals



- The individual production function h_i has two components.
 - The **value** $v_i(x_i)$ created by the individual, as a result of receiving resources x_i .
 - The **utility** $u_i(v_i(x_i)) = h_i(x_i)$ of the value created $(u_i$ is normally concave).
 - So a_i reflects the value function v_i (productivity), and p reflects the combined shape of both functions v_i and u_i .

Assume resource distribution is constrained only by a fixed budget. We have the optimization problem

$$\max \sum_{i} u_{i}$$

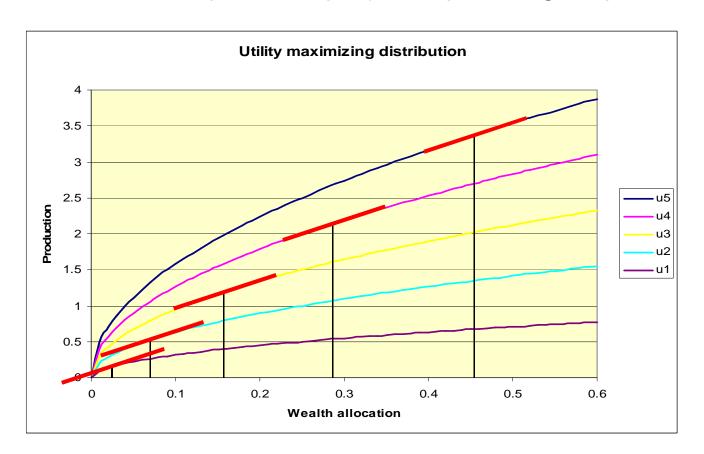
$$u_{i} = a_{i} x_{i}^{p}, \text{ all } i$$

$$\sum_{i} x_{i} = 1, x_{i} \geq 0, \text{ all } i$$

This has a closed-form solution

$$X_i = a_i^{\frac{1}{1-p}} \left(\sum_{j=1}^n a_j^{\frac{1}{1-p}} \right)^{-1}$$

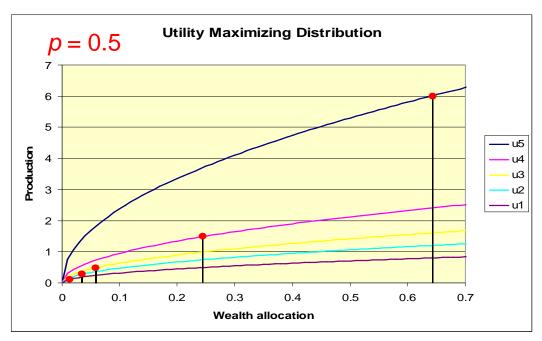
Optimal allocations equalize slope (i.e., equal marginal productivity).



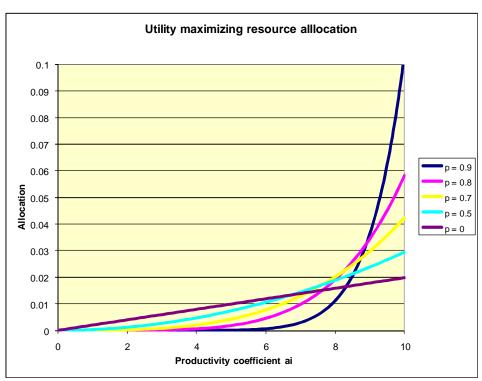
Arguments for utilitarianism

- Can define utility to suit context.
- Utilitarian distributions incorporate some **egalitarian** factors:
- With **concave** production functions, egalitarian distributions create more utility, *ceteris paribus*.
- Inegalitarian distributions create disutility, due to social disharmony.

- Egalitarian distributions create more utility?
 - This effect is **limited**.
 - Utilitarian distributions can be very unequal. Productivity differences are magnified in the allocations.



- Egalitarian distributions create more utility?
 - In the example, the **most egalitarian** distribution ($p \rightarrow 0$) assigns resources in proportion to productivity.



- Unequal distributions create disutility?
 - Perhaps, but modeling this requires **nonseparable** utility functions $u_i = h_i(x_1, \dots, x_n)$

that may result in a problem that is hard to model and solve.

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- Unequal distributions create disutility?
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that may result in a problem that is hard to model and solve.

- More fundamentally, this defense of utilitarianism is based on contingency, not principle.
- If we evaluate the fairness of utilitarian distribution, then there must be another standard of equitable distribution.
- How do we model the standard we really have in mind?

Modeling Utility

- Ideally, production functions are concave, and feasible set is convex.
 - For example, $h_i(x_i) = a_i x_i^p$ for 0 and linear constraints on <math>x.
 - Then we solve the problem

$$\max \sum_{i} h_{i}(x_{i})$$

$$Ax \leq b, x \geq 0$$

by nonlinear programming.

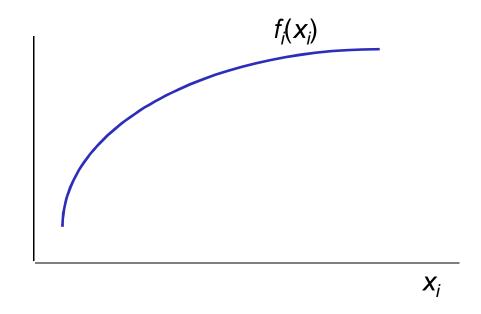
Any local optimum is a global optimum.

- Piecewise linear modeling converts nonlinear programming to LP (linear programming) or MILP (mixed integer/linear programming).
 - A key technique.
 - Applies when functions are separable.
- Suppose we want to solve

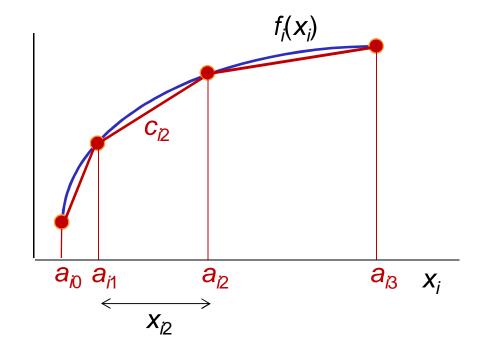
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$$\max \sum_{i} v_{i}$$

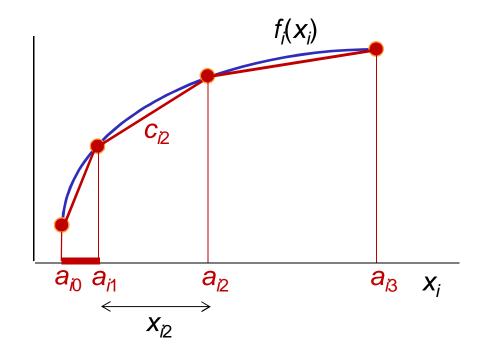
$$v_{i} = f_{i}(a_{0}) + \sum_{j} \frac{\Delta f_{ij}}{\Delta a_{ij}} x_{ij}$$

$$x_{i} = \sum_{j} x_{ij}$$

$$Ax \leq b, \quad x \geq 0$$
where
$$\Delta f_{ij} = f_{i}(a_{ij}) - f_{i}(a_{i,j-1})$$

$$\Delta a_{ii} = a_{ii} - a_{i,i-1}$$

• If each f_i is **concave**, this reduces (approx.) to an **LP**.



The lower intervals "fill up" first.

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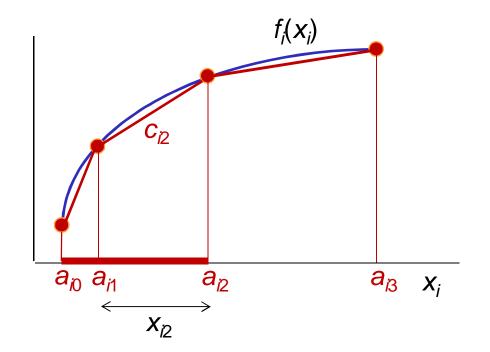
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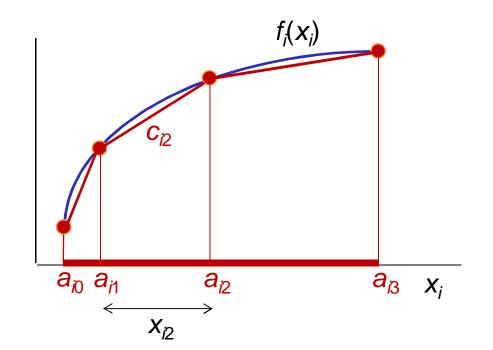
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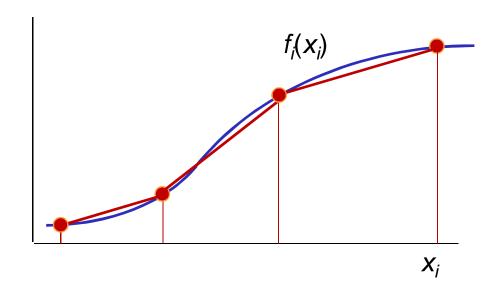
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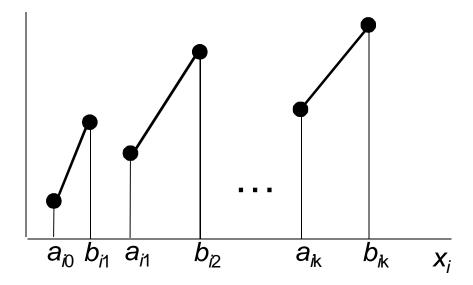
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• If f_i is **nonconcave**, we can use an **MILP** model of the piecewise linear approximation.

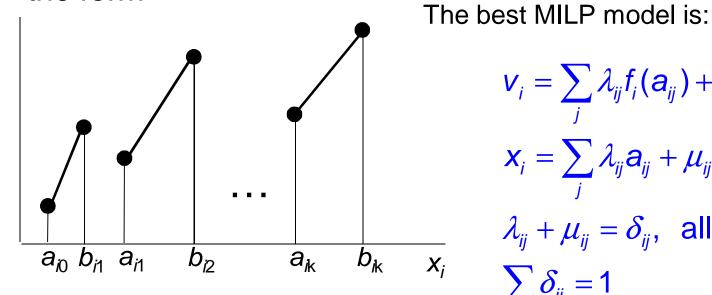


• In general, a piecewise linear approximation v_i of f_i has the form



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The function is continuous when $b_{ij} = a_{i,j+1}$

$$v_{i} = \sum_{j} \lambda_{ij} f_{i}(a_{ij}) + \mu_{ij} f_{i}(b_{ij})$$

$$x_{i} = \sum_{j} \lambda_{ij} a_{ij} + \mu_{ij} b_{ij}$$

$$\lambda_{ij} + \mu_{ij} = \delta_{ij}, \text{ all } j$$

$$\sum_{j} \delta_{ij} = 1$$

$$\lambda_{ij}, \mu_{ij} \geq 0, \quad \delta_{ij} \in \{0,1\}, \text{ all } j$$

When the piecewise linear function is continuous, don't use the "textbook" model

$$v_{i} = \sum_{j=1}^{k+1} \lambda_{ij} f_{i}(a_{ij})$$

$$x_{i} = \sum_{j=1}^{k+1} \lambda_{ij} a_{ij}, \quad \sum_{j=1}^{k} \lambda_{ij} = 1$$

$$\lambda_{ij} \leq \delta_{i,j-1} + \delta_{ij}, \quad j = 2, ..., k$$

$$\lambda_{i1} \leq \delta_{i1}, \quad \lambda_{i,k+1} \leq \delta_{ik}, \quad \sum_{j=1}^{k} \delta_{ij} = 1$$

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The "textbook" may tell you to use only the continuous part of the model

$$v_i = \sum_{j=1}^{k+1} \lambda_{ij} f_i(a_{ij})$$

$$\mathbf{x}_{i} = \sum_{j=1}^{k+1} \lambda_{ij} \mathbf{a}_{ij}$$

and declare the λ_{ii} SOS2.

$$\lambda_{ij}, \mu_{ij} \geq 0, \ \delta_{ij} \in \{0,1\}, \ j = 1,...,k+1$$

where $a_{i,k+1} = b_{ik}$

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$$\mathbf{v}_i = \sum_{j=1}^{k+1} \lambda_{ij} f_i(\mathbf{a}_{ij})$$

$$\mathbf{x}_{i} = \sum_{j=1}^{k+1} \lambda_{ij} \mathbf{a}_{ij}$$

and declare the λ_{ii} SOS2.

This sacrifices the tight relaxation of the next model...

 The best model of a continuous piecewise v_i is the "incremental" formulation:

$$v_{i} = f_{i}(a_{i1}) + \sum_{j=2}^{k+1} \frac{\Delta f_{ij}}{\Delta a_{ij}} x_{ij}$$

$$x_{i} = a_{i1} + \sum_{j=1}^{k} x_{ij}$$

$$\Delta a_{ij} \delta_{ij} \leq x_{ij} \leq \Delta a_{ij} \delta_{i,j-1}, \quad j = 3, ..., k$$

$$\Delta a_{i2} \delta_{ij} \leq x_{i2} \leq \Delta a_{i2}, \quad 0 \leq x_{i,k+1} \leq \Delta a_{i,k+1} \delta_{ik}$$

$$\delta_{ij} \in \{0,1\}, \quad j = 2, ..., k$$

Problems with Utilitarianism

- A utility maximizing distribution may be unjust.
 - Disabled or nonproductive people may be neglected.
 - Less talented people who work hard may receive meager wage.
 - Not all jobs can be equally productive. Those with less productive jobs may receive fewer resources.

Rawlsian Difference Principle

- Rawls' Difference Principle seeks to maximize the welfare of the worst off.
 - Also known as maximin principle.
 - Another formulation: inequality is permissible only to the extent that it is necessary to improve the welfare of those worst off.

$$\max_{i} \min_{i} \{u_{i}\}$$

$$u_{i} = h_{i}(x_{i}), \text{ all } i$$

$$x \in S$$

- The root idea is that when I make a decision for myself, I make a decision for anyone in similar circumstances.
 - It doesn't matter who I am.
- Social contract argument
 - I make decisions (formulate a social contract) in an original position, behind a veil of ignorance as to who I am.
 - I must find the decision acceptable after I learn who I am.
 - I cannot rationally assent to a policy that puts me on the bottom, unless I would have been even worse off under alternative policies.
 - So the policy must maximize the welfare of the worst off.

- Applies only to basic goods.
 - Tings that people want, no matter what else they want.
 - Salaries, tax burden, medical benefits, etc.
 - For example, salary differentials may satisfy the principle if necessary to make the poorest better off.
- Applies to smallest groups for which outcome is predictable.
 - A lottery passes the test even though it doesn't maximize welfare of worst off – the loser is unpredictable.
 - ...unless the lottery participants as a whole are worst off.

- The difference rule implies a lexmax principle.
 - If applied recursively.

• Lexmax (lexicographic maximum) principle:

- Maximize welfare of least advantaged class...
- then next-to-least advantaged class...
- and so forth.

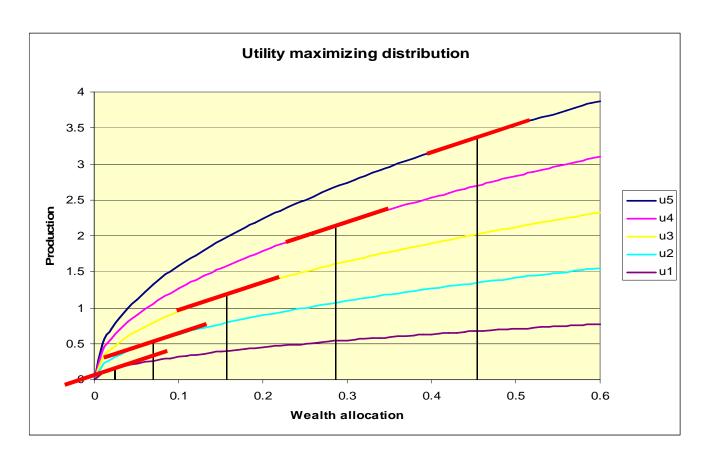
• There is apparently no practical math programming model for lexmax. $\{u_1, ..., u_n\}$

$$u_i = h_i(x_i)$$
, all i
 $x \in S$

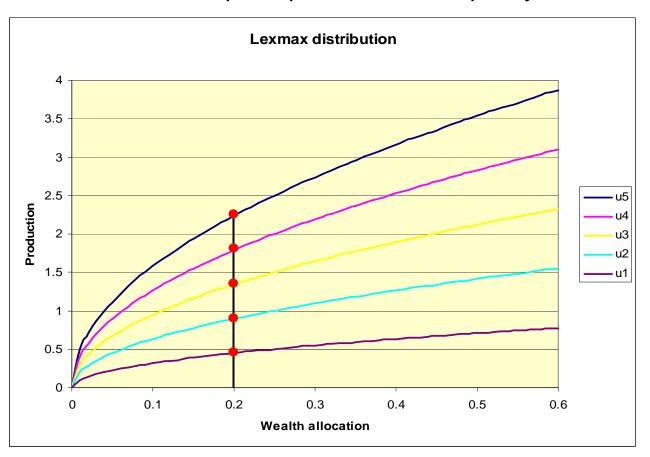
- We can solve the problem sequentially (pre-emptive goal programming).
 - Solve the maximin problem.
 - Fix the smallest u_i to its maximum value.
 - Solve the maximin problem over remaining us.
 - Continue to u_n .

- The Difference and Lexmax Principles need not result in equality.
 - Consider the example presented earlier...

Utilitarian distribution

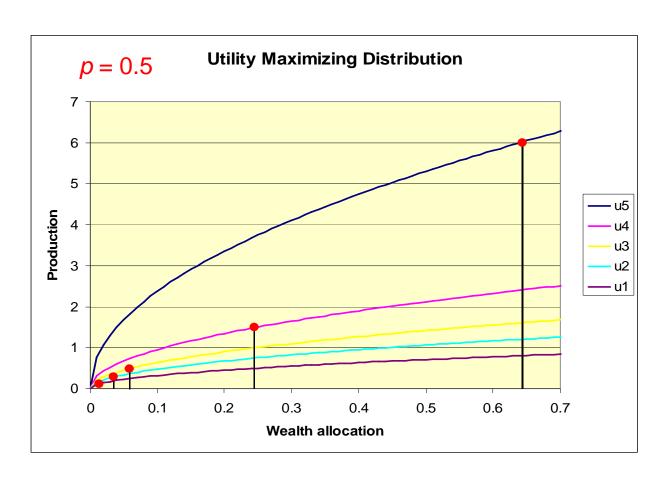


Here, lexmax principle results in equality



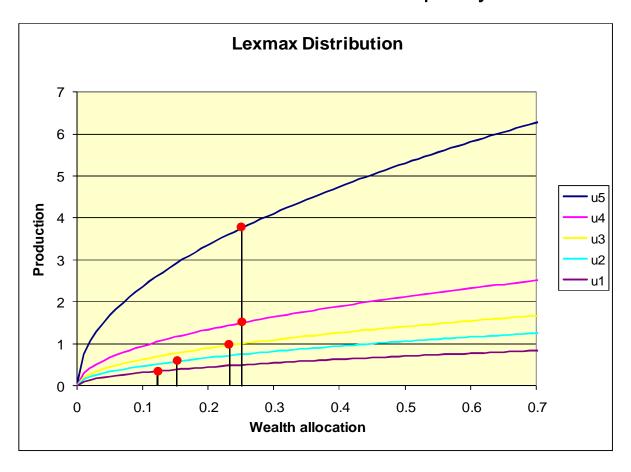
Utilitarianism

But consider this distribution...



Utilitarianism

Lexmax doesn't result in equality



Axiomatics

- The economics literature derives social welfare functions from axioms of rational choice.
 - Some axioms are strong and hard to justify.
 - The social welfare function depends on degree of interpersonal comparability of utilities.
 - Arrow's impossibility theorem was the first result, but there are many others.

Social welfare function

- A function $f(u_1,...,u_n)$ of individual utilities.
- An optimization model can find a distribution of utility that maximizes social welfare.

Interpersonal Comparability

Social Preferences

- Let $u = (u_1, ..., u_n)$ be the vector of utilities allocated to individuals.
- A social welfare function ranks distributions: u is preferable to u' if f(u) > f(u').
- Invariance transformations.
 - These are transformations φ of utility vectors under which the ranking of distributions does not change.
 - Each $\phi = (\phi_1, ..., \phi_n)$, where ϕ_i is a transformation of individual utility u_i .

Interpersonal Comparability

- Ordinal noncomparability.
 - Any $\phi = (\phi_1, ..., \phi_n)$ with strictly increasing ϕ_i s is an invariance transformation.
- Ordinal level comparability.
 - Any $\phi = (\phi_1, ..., \phi_n)$ with strictly increasing and identical ϕ_i s is an invariance transformation.

Interpersonal Comparability

- Cardinal nonncomparability.
 - Any $\phi = (\phi_1, ..., \phi_n)$ with $\phi_i(u_i) = \alpha_i + \beta_i u_i$ and $\beta_i > 0$ is an invariance transformation.
- Cardinal unit comparability.
 - Any $\phi = (\phi_1, ..., \phi_n)$ with $\phi_i(u_i) = \alpha_i + \beta u_i$ and $\beta > 0$ is an invariance transformation.
- Cardinal ratio scale comparability
 - Any $\phi = (\phi_1, ..., \phi_n)$ with $\phi_i(u_i) = \beta u_i$ and $\beta > 0$ is an invariance transformation.

Axioms

- Anonymity
 - Social preferences are the same if indices of u_is are permuted.
- Strict pareto
 - If u > u', then u is preferred to u'.
- Independence of irrelevant alternatives
 - The preference of *u* over *u'* depends only on *u* and *u'* and not on what other utility vectors are possible.
- Separability of unconcerned individuals
 - Individuals *i* for which $u_i = u_i'$ don't affect the ranking of u and u'.

Axiomatics

Theorem

Given **ordinal level comparability**, any social welfare function *f* that satisfies the axioms is lexicographically increasing or lexicographically decreasing. So we get a **lexmax** or **lexmin** objective.

Theorem

Given **cardinal unit comparability**, any social welfare function f that satisfies the axioms has the form $f(u) = \sum_i a_i u_i$ for $a_i \ge 0$. Se we get a **utilitarian** objective.

Axiomatics

Theorem

Given **cardinal noncomparability**, any social welfare function f that satisfies the axioms (except anonimity and separability) has the form $f(u) = u_i$ for some fixed i. So individual i is a **dictator**.

Theorem

Given **cardinal ratio scale comparability**, any social welfare function f that satisfies the axioms has the form $f(u) = \sum_i u_i^p/p$. Se we get the production function used in the example.

Measures of Inequality

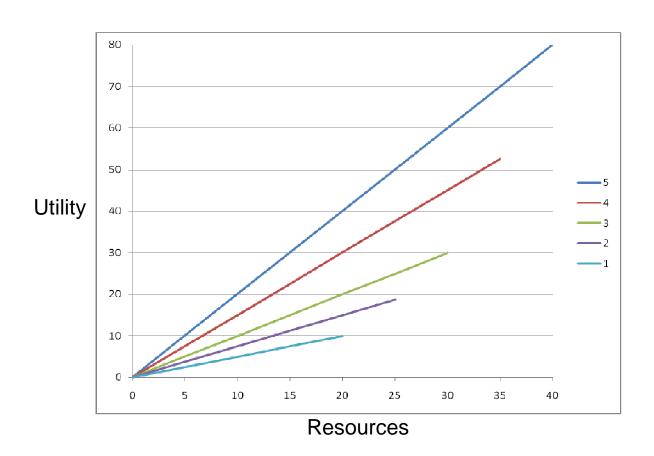
- Assume we wish to minimize inequality.
 - We will survey several measures of inequality.
 - They have different strengths and weaknesses.
 - Minimizing inequality may result in less total utility.
- Pigou-Dalton condition.
 - One criterion for evaluating an inequality measure.
 - If utility is transferred from one who is worse off to one who is better off, inequality should increase.

Measures of Inequality

- Measures of Inequality
 - An example
 - Utrilitarian, maximin, and lexmax solution
 - Relative range, max, min
 - Relative mean deviation
 - Variance, coefficient of variation
 - McLoone index
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- An Allocation Problem

Example

Production functions for 5 individuals



Utilitarian

$$\max \sum_{i} u_{i}$$

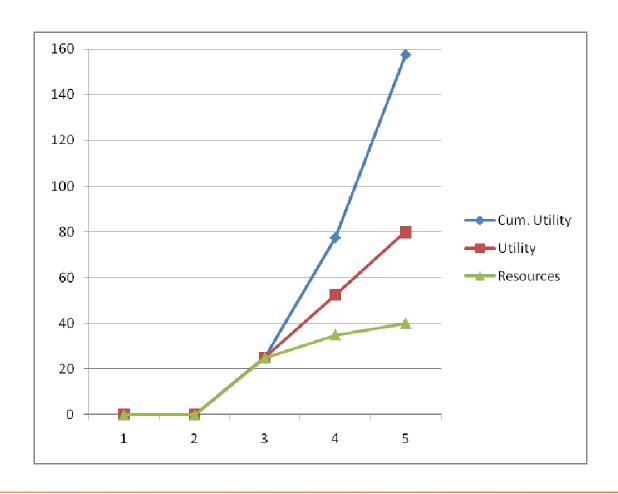
LP model:
$$\max \sum_{i=1}^{5} u_i$$

 $u_i = a_i x_i, \ 0 \le x_i \le b_i, \ \text{all } i, \ \sum_i x_i = B$

where
$$(a_1, ..., a_5) = (0.5, 0.75, 1, 1.5, 2)$$

 $(b_1, ..., b_5) = (20, 25, 30, 35, 40)$
 $B = 100$

Utilitarian



Rawlsian

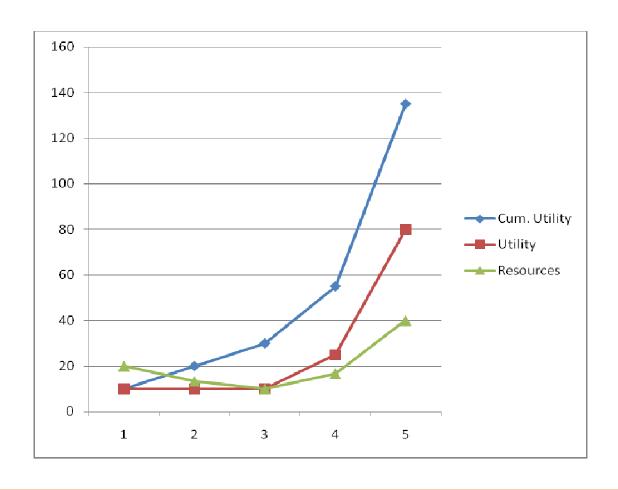
$$\max \left\{ \min_{i} \left\{ u_{i} \right\} \right\}$$

LP model:
$$\max u_{\min} + \in \sum_{i} u_{i} \leftarrow \max_{\substack{i \text{ pareto optimal} \\ u_{\min} \leq u_{i}, \text{ all } i}} \text{Ensures that solution is Pareto optimal}$$

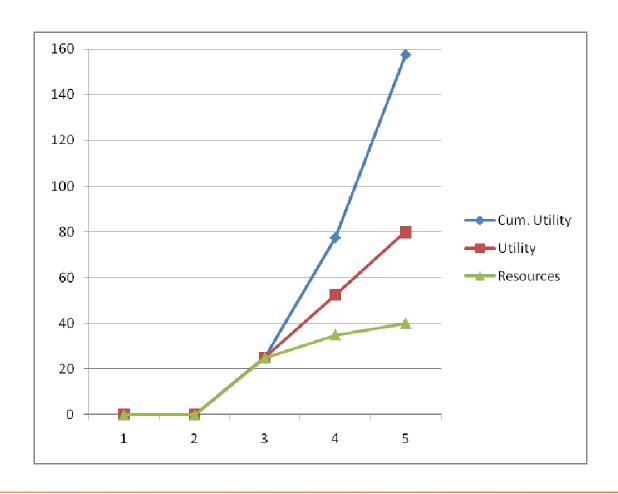
$$u_{\min} \leq u_{i}, \text{ all } i$$

$$u_{i} = a_{i}x_{i}, \quad 0 \leq x_{i} \leq b_{i}, \text{ all } i, \quad \sum_{i} x_{i} = B$$

Rawlsian



Utilitarian



Lexmax

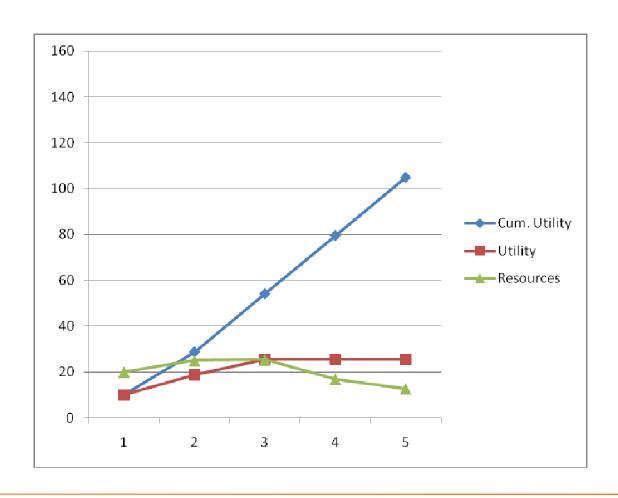
$$\operatorname{lexmax}\left\{u_{1},\ldots,u_{n}\right\}$$

Sequence of
$$\max u_{\min}$$

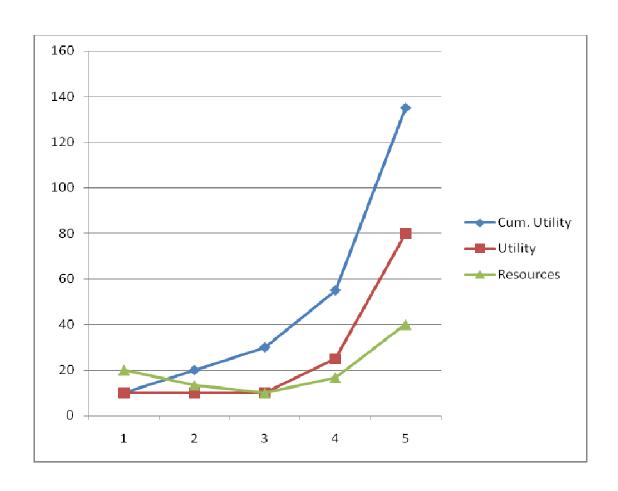
LP models, $u_i = u_i^*$, all $i < k$
 $u_{\min} \le u_i$, all $i \ge k$
 $u_i = a_i x_i$, $0 \le x_i \le b_i$, all i , $\sum_i x_i = B$

Re-index for each k so that u_i for i < k were fixed in previous iterations.

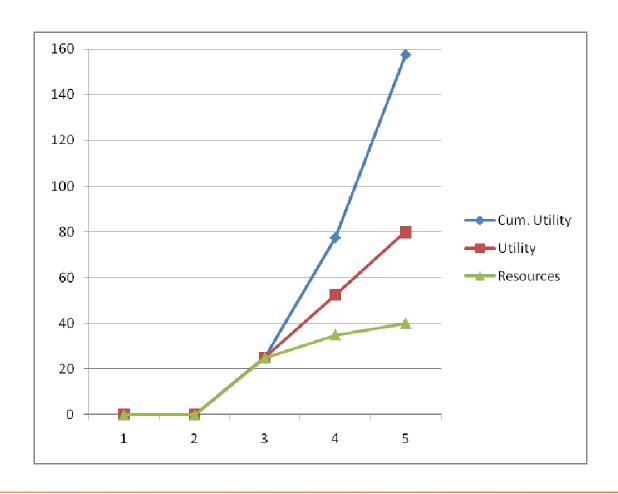
Lexmax



Rawlsian



Utilitarian



$$\frac{U_{\max} - U_{\min}}{\overline{U}}$$

where
$$u_{\text{max}} = \max_{i} \{u_i\}$$
 $u_{\text{min}} = \min_{i} \{u_i\}$ $\overline{u} = (1/n)\sum_{i} u_i$

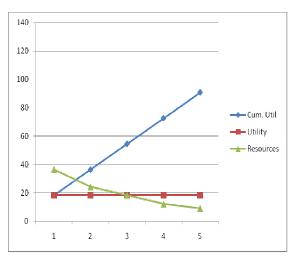
Rationale:

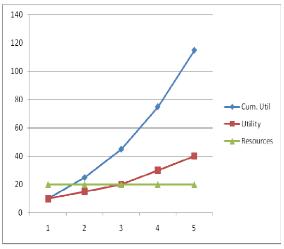
- Perceived inequality is relative to the best off.
- A distribution should be judged by the position of the worst-off.
- Therefore, minimize gap between top and bottom.

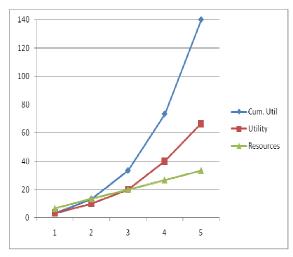
Problems:

- Ignores distribution between extremes.
- Violates Pigou-Dalton condition

Equality Measures: Comparison







Relative range: 0 1.30 2.26

$$\frac{U_{\max} - U_{\min}}{\overline{U}}$$

This is a **fractional linear programming** problem.

Use Charnes-Cooper transformation to an LP. In general,

$$\min \frac{cx + c_0}{dx + d_0}$$

$$Ax \ge b$$

$$x \ge 0$$

$$\min cx' + c_0z$$

$$Ax' \ge bz$$

$$dx' + d_0z = 1$$

$$x', z \ge 0$$

after change of variable x = x'/z and fixing denominator to 1.

$$\frac{U_{\max} - U_{\min}}{\overline{U}}$$

Fractional LP model:
$$\min \frac{u_{\max} - u_{\min}}{(1/n)\sum_{i} u_{i}}$$

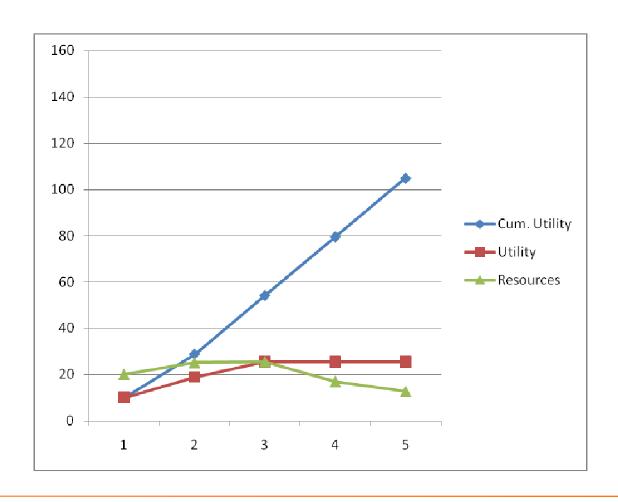
 $u_{\max} \ge u_{i}, \ u_{\min} \le u_{i}, \ \text{all } i$
 $u_{i} = a_{i}x_{i}, \ 0 \le x_{i} \le b_{i}, \ \text{all } i, \ \sum_{i} x_{i} = B$

LP model:
$$\min u_{\max} - u_{\min}$$

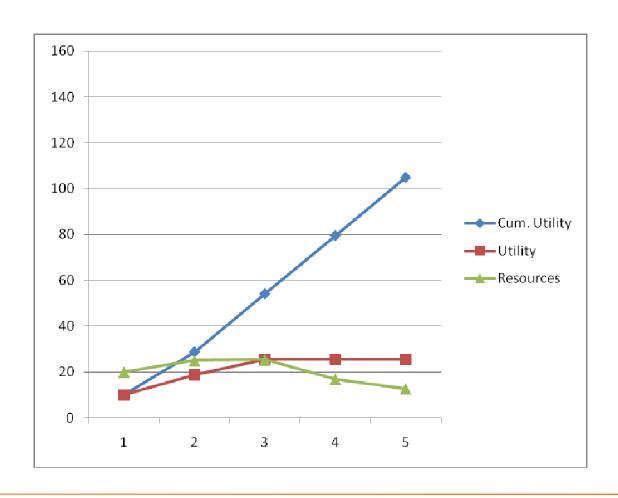
$$u_{\max} \ge u'_i, \ u_{\min} \le u'_i, \ \text{all } i$$

$$u'_i = a_i x'_i, \ 0 \le x'_i \le b_i z, \ \text{all } i, \quad \sum_i x'_i = Bz$$

$$(1/n) \sum_i u'_i = 1$$



Lexmax



Relative Max

$$\frac{U_{\max}}{\overline{U}}$$

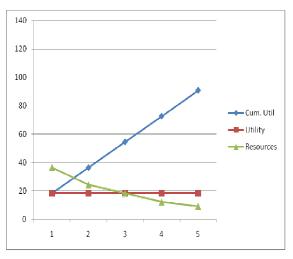
Rationale:

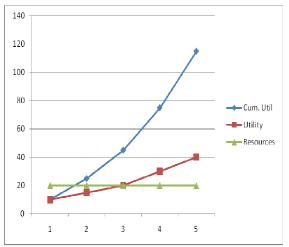
- Perceived inequality is relative to the best off.
- Possible application to salary levels (typical vs. CEO)

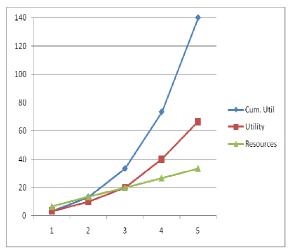
Problems:

- Ignores distribution below the top.
- Violates Pigou-Dalton condition

Equality Measures: Comparison







Relative range:

0

1.30

2.26

Relative max:

- 1

1.73

2.38

Relative Max

$$\frac{u_{\max}}{\overline{u}}$$

$$\min \frac{u_{\text{max}}}{(1/n)\sum_{i} u_{i}}$$

$$u_{\text{max}} \ge u_{i}, \text{ all } i$$

$$u_{i} = a_{i}x_{i}, 0 \le x_{i} \le b_{i}, \text{ all } i, \sum_{i} x_{i} = B$$

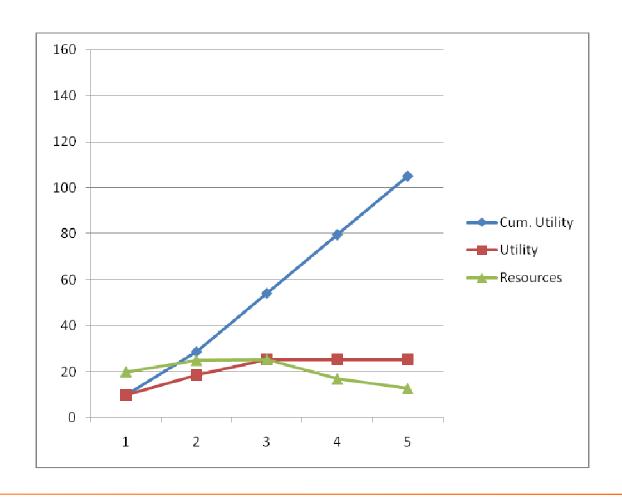
min
$$u_{\text{max}}$$

$$u_{\text{max}} \ge u'_{i} \text{ all } i$$

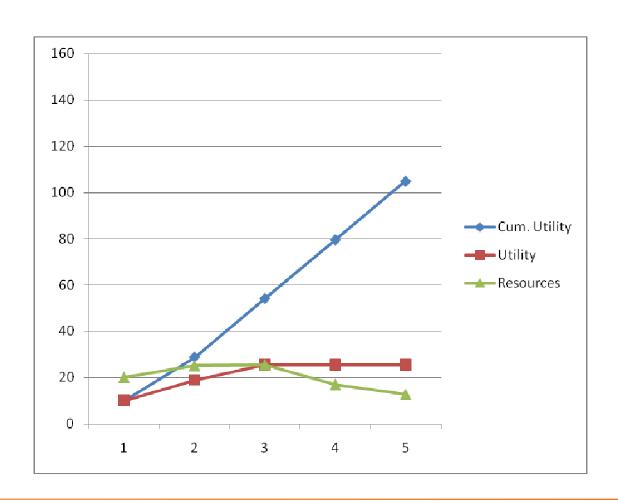
$$u'_{i} = a_{i}x'_{i}, \quad 0 \le x'_{i} \le b_{i}z, \quad \text{all } i, \quad \sum_{i} x'_{i} = Bz$$

$$(1/n)\sum_{i} u'_{i} = 1$$

Relative Max



Relative Range



Relative Min

$$\frac{U_{\min}}{\overline{U}}$$

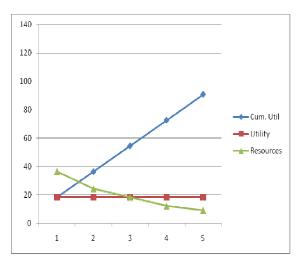
Rationale:

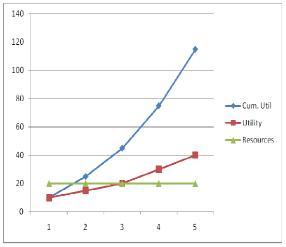
- Measures adherence to Rawlsian Difference Principle.
- ...relativized to mean

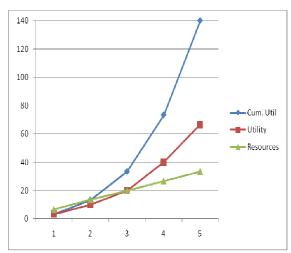
Problems:

- Ignores distribution above the bottom.
- Violates Pigou-Dalton condition

Equality Measures: Comparison







Relative range: Relative max:

Relative min:

1.30

1.73

0.43

2.26

2.38

0.12

Relative Min

$$\frac{U_{\min}}{\overline{U}}$$

$$\max \frac{u_{\min}}{(1/n)\sum_{i} u_{i}}$$

$$u_{\min} \leq u_{i}, \text{ all } i$$

$$u_i = a_i x_i$$
, $0 \le x_i \le b_i$, all i , $\sum_i x_i = B$

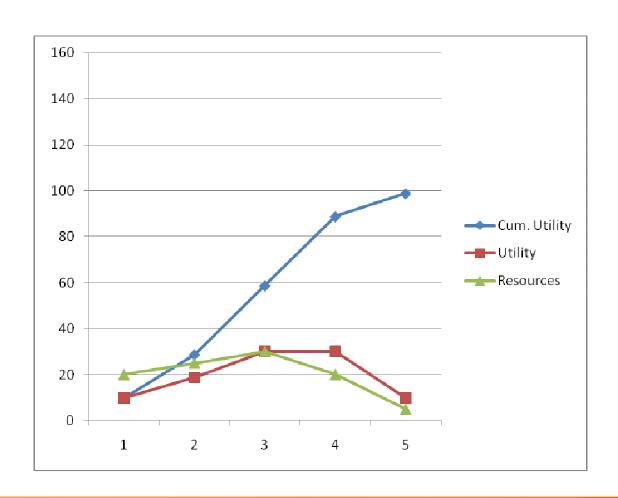
$$\max u_{\min}$$

$$u_{\min} \ge u'_{i} \text{ all } i$$

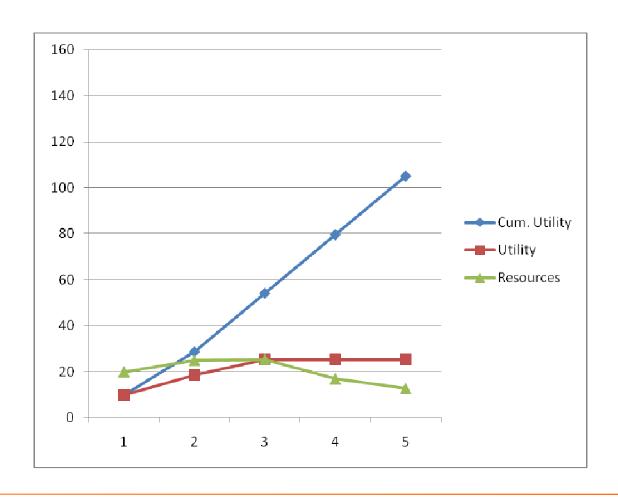
$$u'_{i} = a_{i}x'_{i}, \quad 0 \le x'_{i} \le b_{i}z, \quad \text{all } i, \quad \sum_{i} x'_{i} = Bz$$

$$(1/n)\sum_{i} u'_{i} = 1$$

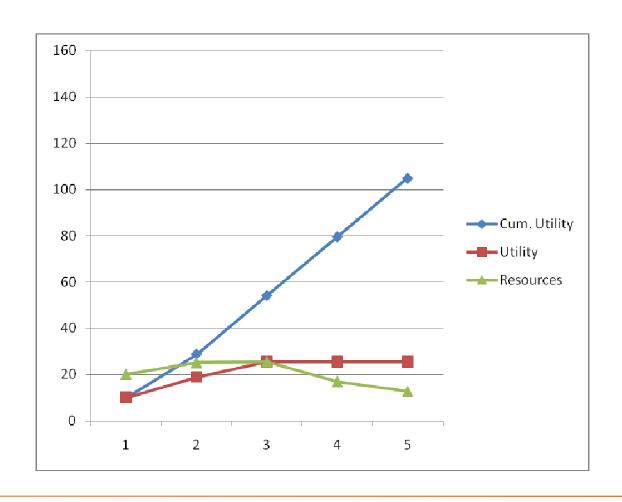
Relative Min



Relative Max



Relative Range



Relative Mean Deviation

$$\frac{\sum_{i} |u_{i} - \overline{u}|}{\overline{u}}$$

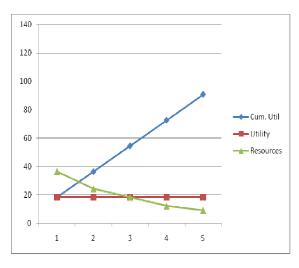
Rationale:

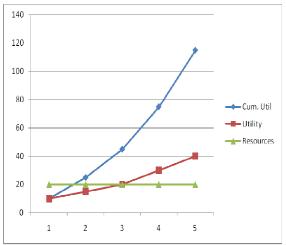
- Perceived inequality is relative to average.
- Entire distribution should be measured.

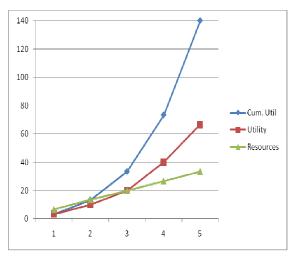
Problems:

- Violates Pigou-Dalton condition
- Insensitive to transfers on the same side of the mean.
- Insensitive to placement of transfers from one side of the mean to the other.

Equality Measures: Comparison







Relative range:

0

1.30

2.26

Rel. mean dev.:

 \mathbf{O}

0.42

0.72

Relative Mean Deviation

Fractional LP model:
$$\frac{\sum_{i} \left| u_{i} - \overline{u} \right| }{\overline{u}}$$

$$u_{i}^{+} \geq u_{i} - \overline{u}, \ u_{i}^{-} \geq \overline{u} - u_{i}, \ \text{all } i$$

$$u_{i} = a_{i} x_{i}, \ 0 \leq x_{i} \leq b_{i}, \ \text{all } i, \quad \sum_{i} x_{i} = B$$

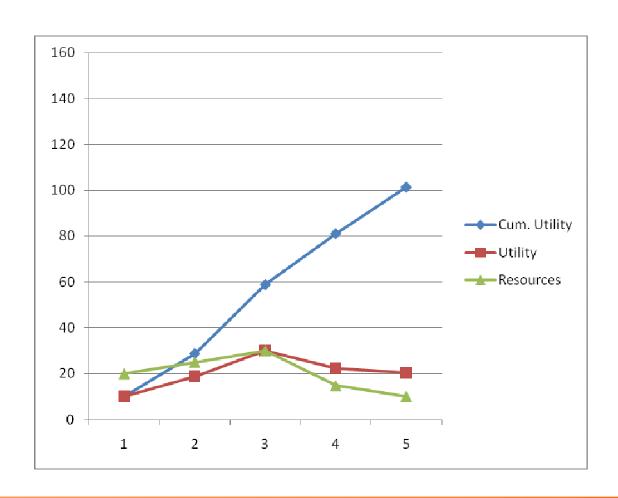
$$\text{LP model:} \quad \max \sum_{i} (u_{i}^{+} + u_{i}^{-})$$

$$u_{i}^{+} \geq u_{i}^{'} - 1, \ u_{i}^{-} \leq u_{i}^{'} - 1, \ \text{all } i$$

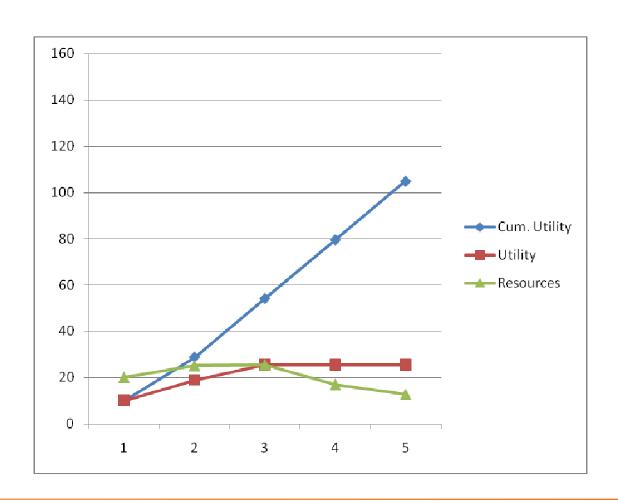
$$(1/n) \sum_{i} u_{i}^{'} = 1$$

$$u_{i}^{'} = a_{i} x_{i}^{'}, \ 0 \leq x_{i}^{'} \leq b_{i} z, \ \text{all } i, \quad \sum x_{i}^{'} = Bz$$

Relative Mean Deviation



Relative Range



$$(1/n)\sum_{i}(u_{i}-\overline{u})^{2}$$

Rationale:

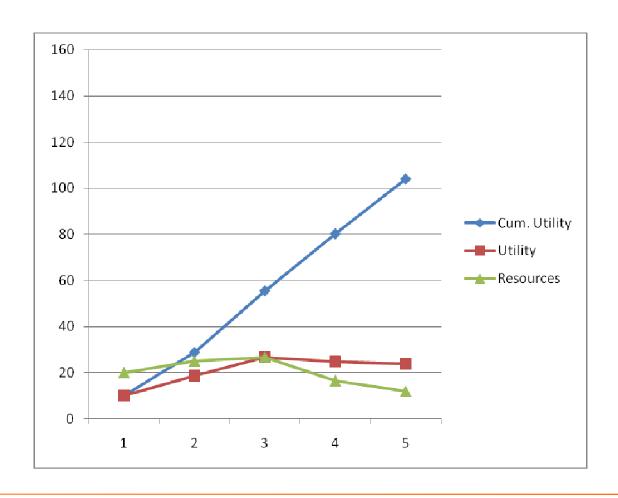
- Weight each utility by its distance from the mean.
- Satisfies Pigou-Dalton condition.
- Sensitive to transfers on one side of the mean.
- Sensitive to placement of transfers from one side of the mean to the other.

Problems:

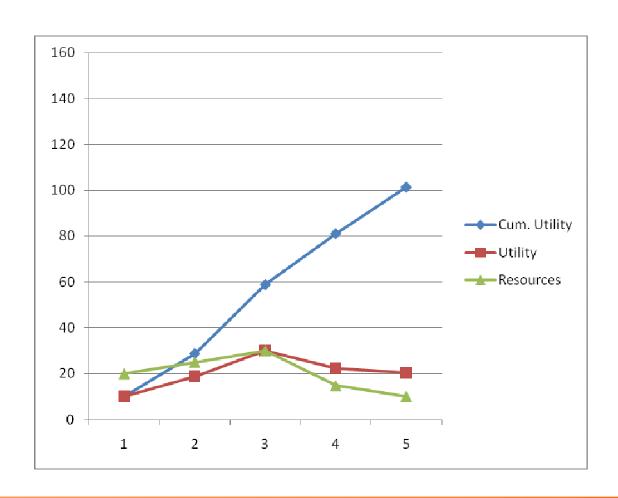
- Weighting is arbitrary?
- Variance depends on scaling of utility.

$$(1/n)\sum_{i}(u_{i}-\overline{u})^{2}$$

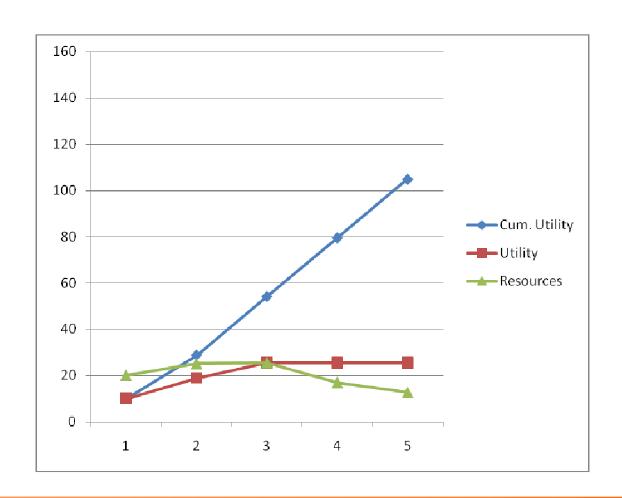
Convex nonlinear model: $\min (1/n) \sum_{i} (u_i - \overline{u})^2$ $\overline{u} = (1/n) \sum_{i} u_i$ $u_i = a_i x_i, \ 0 \le x_i \le b_i, \ \text{all } i, \quad \sum_{i} x_i = B$



Relative Mean Deviation



Relative Range



$$\frac{\left((1/n)\sum_{i}(u_{i}-\overline{u})^{2}\right)^{1/2}}{\overline{u}}$$

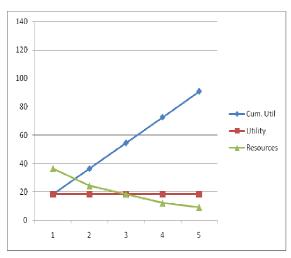
Rationale:

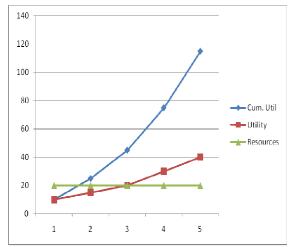
- Similar to variance.
- Invariant with respect to scaling of utilities.

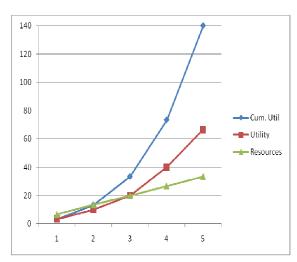
Problems:

- When minimizing inequality, there is an incentive to reduce average utility.
- Should be minimized only for fixed total utility.

Equality Measures: Comparison







Relative range: 0

Rel. mean dev.: 0

Coeff. of variation: 0

1.30

0.42

0.46

2.26

0.72

0.81

$$\frac{\left((1/n)\sum_{i}(u_{i}-\overline{u})^{2}\right)^{1/2}}{\overline{u}}$$

Again use change of variable u = u'/z and fix denominator to 1.

$$\min \frac{\left((1/n)\sum_{i}(u_{i}-\overline{u})^{2}\right)^{1/2}}{\overline{u}} \quad \min \left((1/n)\sum_{i}(u_{i}'-1)^{2}\right)^{1/2}$$

$$\Delta u \geq b \quad \text{Can drop}$$

$$\Delta u \geq 0 \quad (1/n)\sum_{i}u_{i}' = 1 \quad \text{to make}$$

$$u \geq 0 \quad u' \geq 0 \quad \text{convex}$$

$$\frac{\left((1/n)\sum_{i}(u_{i}-\overline{u})^{2}\right)^{1/2}}{\overline{u}}$$

Fractional nonlinear model:

$$\max \frac{\left((1/n)\sum_{i}(u_{i}-\overline{u})^{2}\right)^{n2}}{\overline{u}}$$

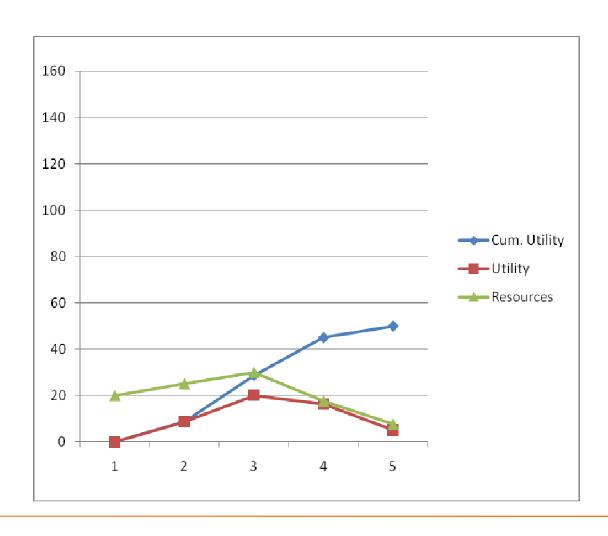
$$\overline{u} = (1/n)\sum_{i}u_{i}$$

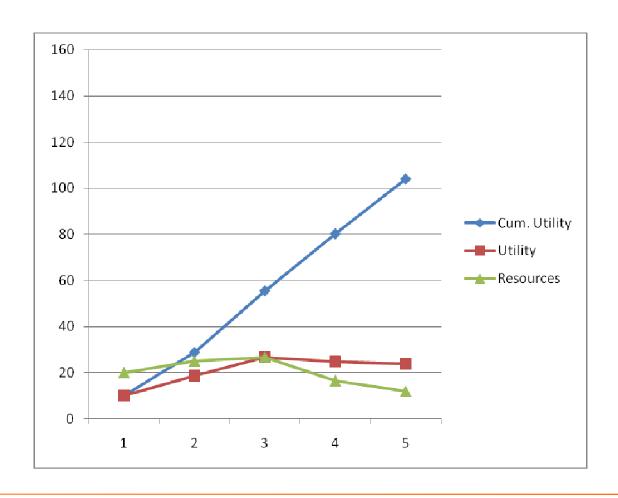
$$u_{i} = a_{i}x_{i}, \ 0 \le x_{i} \le b_{i}, \ \text{all } i, \quad \sum_{i}x_{i} = B$$

Convex nonlinear model:

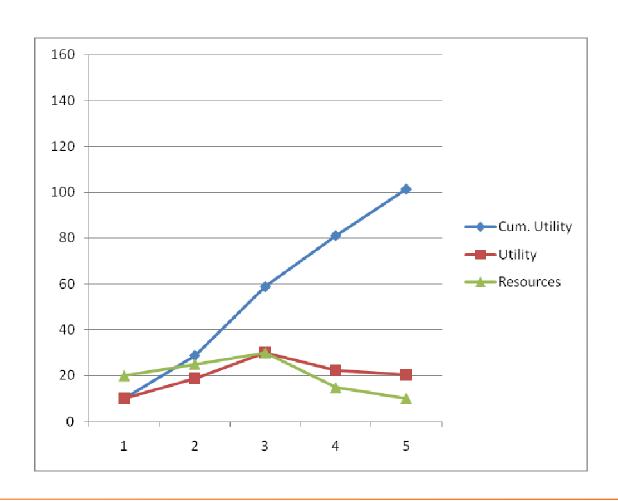
min
$$(1/n)\sum_{i} (u'_{i} - 1)^{2}$$

 $(1/n)\sum_{i} u'_{i} = 1$
 $u'_{i} = a_{i}x'_{i}, \ 0 \le x'_{i} \le b_{i}z, \ \text{all } i, \ \sum_{i} x'_{i} = Bz$





Relative Mean Deviation



$$\frac{(1/2)\sum_{i:u_i< m}u_i}{\overline{u}}$$

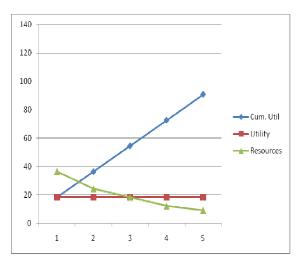
Rationale:

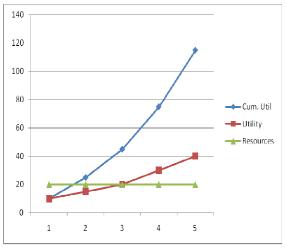
- Ratio of average utility below median to overall average.
- No one wants to be "below average."
- Pushes average up while pushing inequality down.

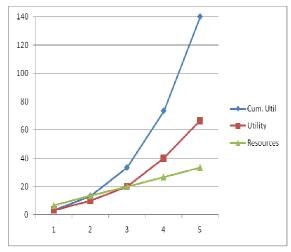
Problems:

- Violates Pigou-Dalton condition.
- Insensitive to upper half.

Equality Measures: Comparison







Relative range: 0

Rel. mean dev.: 0

Coeff. of variation: 0

McLoone:

1.30

0.42

0.46

0.54

2.26

0.72

0.81

0.23

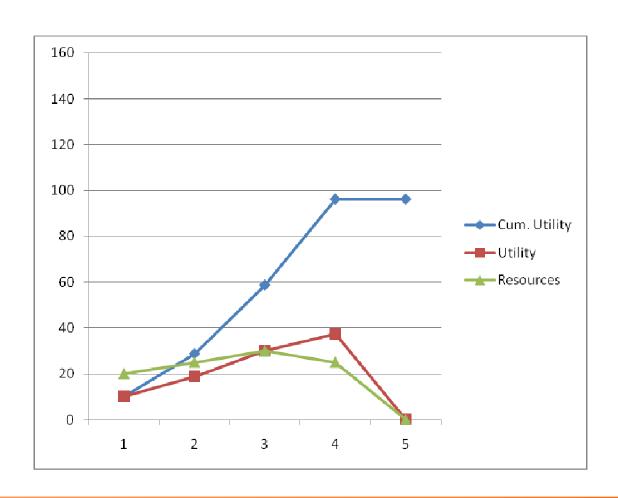
$$\frac{(1/2)\sum_{i:u_i< m}u_i}{\overline{u}}$$

Fractional MILP model:
$$\sum_{i=1}^{n} v_{i}$$
 Defines median $m \longrightarrow m-My_{i} \leq u_{i} \leq m+M(1-y_{i})$, all i Defines $v_{i}=u_{i}$ if
$$v_{i} \leq u_{i}, v_{i} \leq My_{i}, \text{ all } i$$
 Defines we median
$$\sum_{i=1}^{n} y_{i} < n/2$$
 Half of utilities
$$u_{i} = a_{i}x_{i}, \ 0 \leq x_{i} \leq b_{i}, \text{ all } i, \sum_{i=1}^{n} x_{i} = B$$
 are below median
$$y_{i} \in \{0,1\}, \text{ all } i$$
 Selects utilities below median

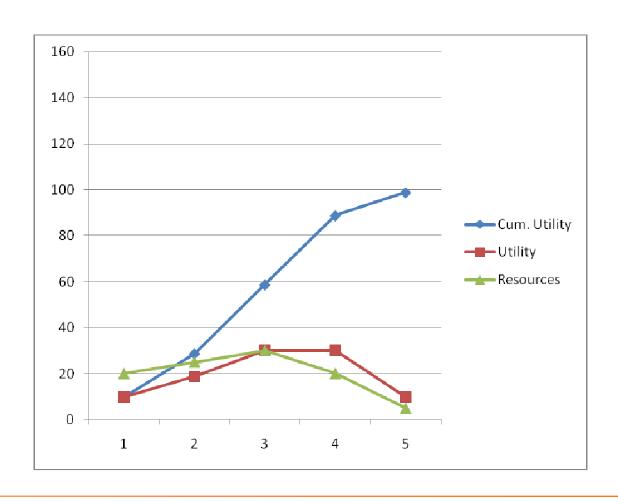
$$\frac{(1/2)\sum_{i:u_i < m} u_i}{\overline{u}}$$

MILP model:
$$\max \sum_{i} v'_{i}$$

 $m' - My_{i} \le u'_{i} \le m' + M(1 - y_{i})$, all i
 $v'_{i} \le u'_{i}, v'_{i} \le My_{i}$, all i
 $\sum_{i} y_{i} < n/2$
 $u'_{i} = a_{i}x'_{i}, \ 0 \le x'_{i} \le b_{i}z$, all i , $\sum_{i} x'_{i} = Bz$
 $y_{i} \in \{0,1\}$, all i



Relative Min



Gini Coefficient

$$\frac{(1/n^2)\sum_{i,j}\left|u_i-u_j\right|}{2\overline{u}}$$

Rationale:

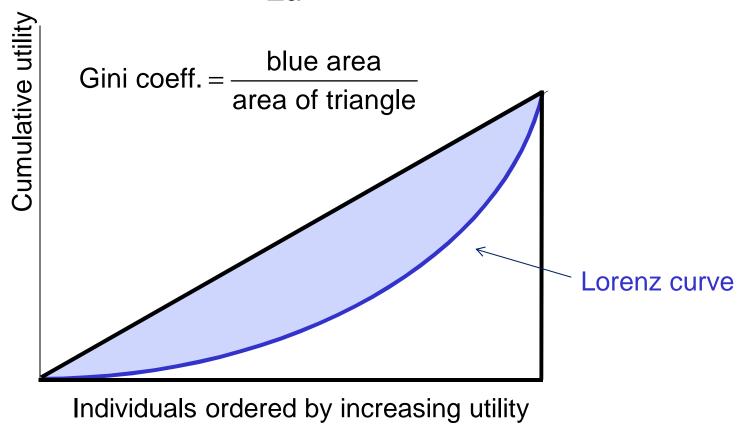
- Relative mean difference between all pairs.
- Takes all differences into account.
- Related to area above cumulative distribution (Lorenz curve).
- Satisfies Pigou-Dalton condition.

Problems:

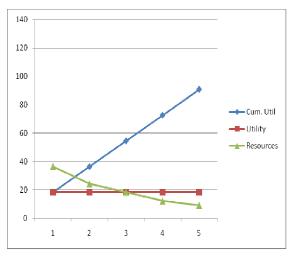
Insensitive to shape of Lorenz curve, for a given area.

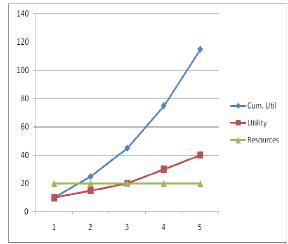
Gini Coefficient

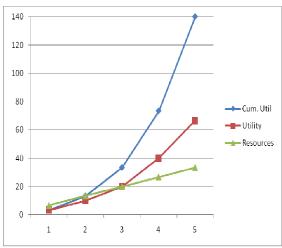
$$\frac{(1/n^2)\sum_{i,j}\left|u_i-u_j\right|}{2\overline{u}}$$



Equality Measures: Comparison







Relative range:	0	
Rel. mean dev.:	0	
Coeff. of variation:	0	
McLoone:	1	
Gini:	0	

1.30	
0.42	
0.46	
0.54	
0.26	

_		_	
	2.26	3	
	0.72	2	
	0.8	1	
	0.23	3	
	0.4	5	

Gini Coefficient

$$\frac{(1/n^2)\sum_{i,j}\left|u_i-u_j\right|}{2\overline{u}}$$

Fractional LP model:
$$\max \frac{(1/2n^2)\sum_{ij}(u_{ij}^+ + u_{ij}^-)}{\overline{u}}$$
$$u_{ij}^+ \ge u_i - u_j, \ u_{ij}^- \ge u_j - u_i, \ \text{all } i, j$$
$$\overline{u} = (1/n)\sum_i u_i$$
$$u_i = a_i x_i, \ 0 \le x_i \le b_i, \ \text{all } i, \ \sum_i x_i = B$$

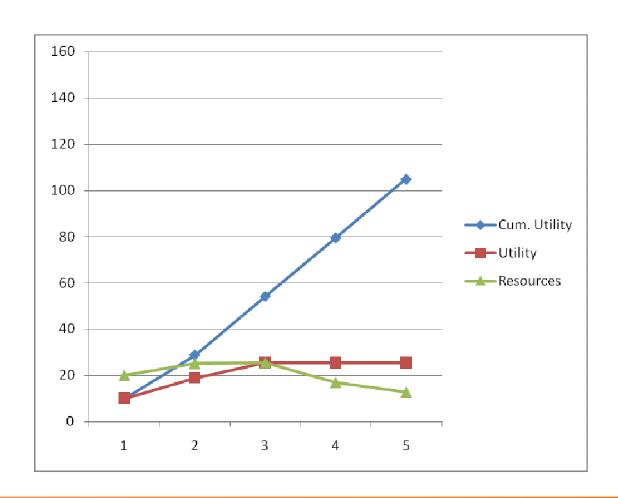
LP model:
$$\max_{i,j} (1/2n^2) \sum_{i,j} (u_{ij}^+ + u_{ij}^-)$$

$$u_{ij}^+ \ge u_i' - u_j', \ u_{ij}^- \ge u_j' - u_i', \ \text{all } i, j$$

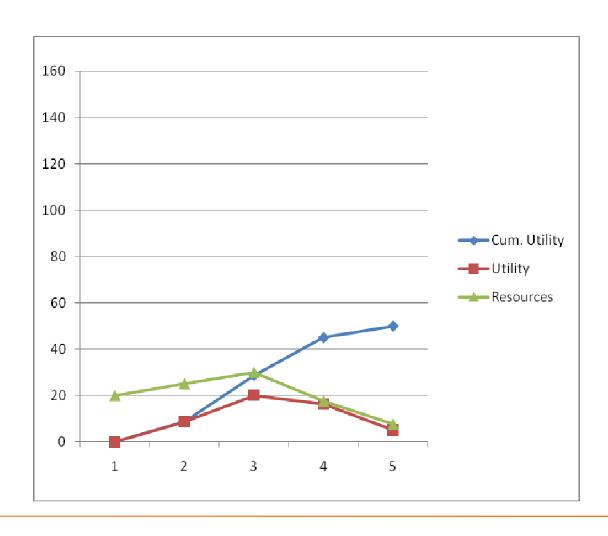
$$(1/n) \sum_{i} u_i' = 1$$

$$u_i' = a_i x_i', \ 0 \le x_i' \le b_i z, \ \text{all } i, \ \sum_{i} x_i' = Bz$$

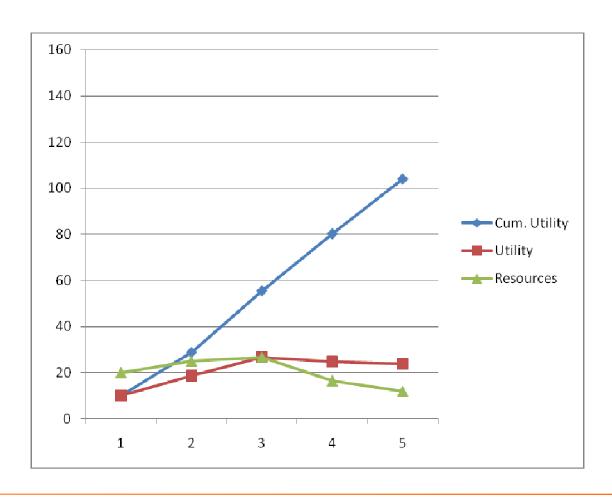
Gini Coefficient



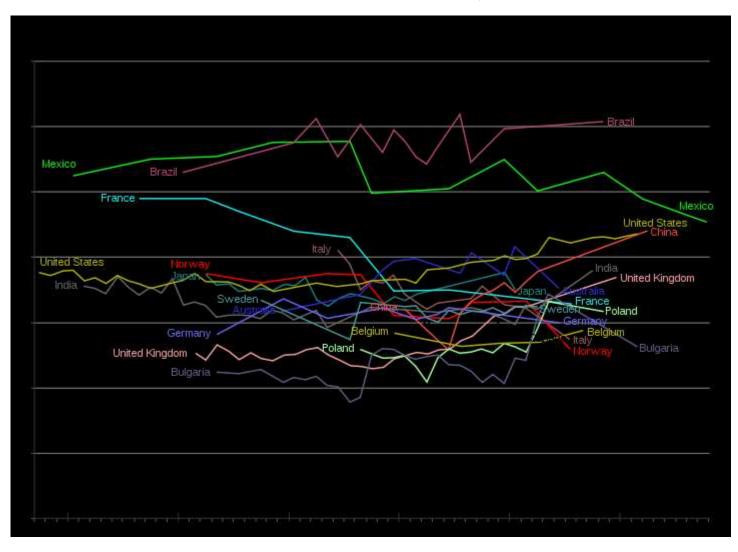
Coefficient of Variation



Variance



Historical Gini Coefficient, 1945-2010



$$1 - \left(\frac{(1/n) \sum_{i} \left(\frac{X_{i}}{\overline{X}} \right)^{p}}{1/p} \right)^{1/p}$$

Rationale:

- Best seen as measuring inequality of resources x_i.
- Assumes allotment y of resources results in utility y^p
- This is average utility per individual.

$$1 - \left((1/n) \sum_{i} \left(\frac{x_{i}}{\overline{x}} \right)^{\rho} \right)^{1/\rho}$$

Rationale:

- Best seen as measuring inequality of resources x_i.
- Assumes allotment y of resources results in utility y^p
- This is average utility per individual.
- This is equal resource allotment to each individual that results in same total utility.

$$1 - \left((1/n) \sum_{i} \left(\frac{X_{i}}{\overline{X}} \right)^{p} \right)^{1/p}$$

Rationale:

- Best seen as measuring inequality of resources x_i.
- Assumes allotment y of resources results in utility y^p
- This is average utility per individual.
- This is equal resource allotment to each individual that results in same total utility.
- This is additional resources per individual necessary to sustain inequality.

$$1 - \left((1/n) \sum_{i} \left(\frac{x_{i}}{\overline{x}} \right)^{p} \right)^{1/p}$$

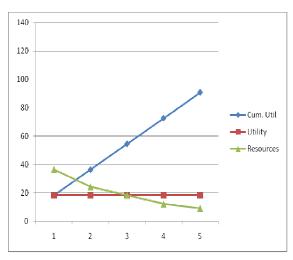
Rationale:

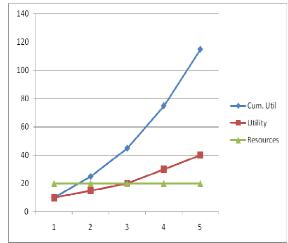
- p indicates "importance" of equality.
- Similar to L_p norm
- p = 1 means inequality has no importance
- p = 0 is Rawlsian (measures utility of worst-off individual).

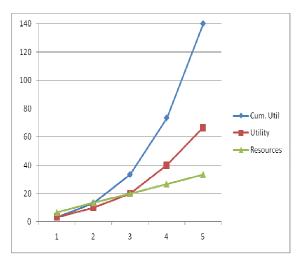
Problems:

- Measures utility, not equality.
- Doesn't evaluate distribution of utility, only of resources.
- p describes utility curve, not importance of equality.

Equality Measures: Comparison







Relative range:	0	1.30	2.26
Rel. mean dev.:	0	0.42	0.72
Coeff. of variation:	0	0.46	0.81
McLoone:	1	0.54	0.23
Gini:	0	0.26	0.45
Atkinson	0.06	0	0.06

$$1 - \left((1/n) \sum_{i} \left(\frac{x_{i}}{\overline{x}} \right)^{p} \right)^{1/p}$$

To minimize index, solve fractional problem

$$\max \sum_{i} \left(\frac{x_{i}}{\overline{x}} \right)^{\rho} = \frac{\sum_{i} x_{i}^{\rho}}{\overline{x}^{\rho}}$$

$$Ax \ge b, \quad x \ge 0$$

After change of variable $x_i = x_i'/z$, this becomes

$$\max \sum_{i} x_{i}^{\prime p}$$

$$(1/n) \sum_{i} x_{i}^{\prime} = 1$$

$$Ax^{\prime} \ge bz, \quad x^{\prime} \ge 0$$

$$1 - \left((1/n) \sum_{i} \left(\frac{x_{i}}{\overline{x}} \right)^{p} \right)^{1/p}$$

Fractional nonlinear model:

$$\max \frac{\sum_{i} x_{i}^{p}}{\overline{x}^{p}}$$

$$\overline{x} = (1/n) \sum_{i} x_{i}$$

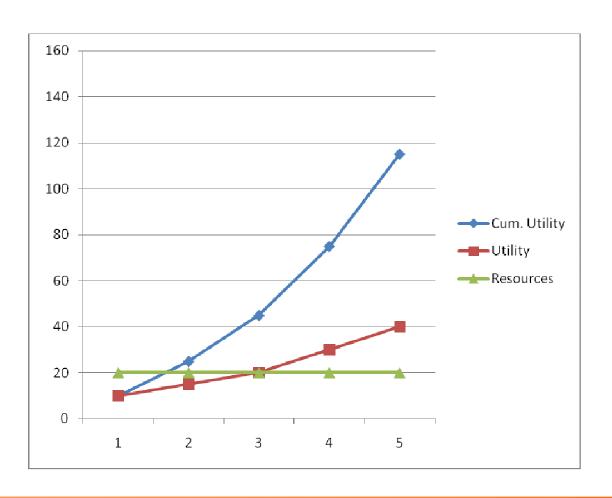
$$\sum_{i} x_{i} = B, \quad x \ge 0$$

Concave nonlinear model:

$$\max \sum_{i} x_{i}^{\prime p}$$

$$(1/n)\sum_{i} x_{i}^{\prime} = 1$$

$$\sum_{i} x_{i}^{\prime} = Bz, \quad x^{\prime} \ge 0$$



Hoover Index

$$(1/2)\frac{\sum_{i}|u_{i}-\overline{u}|}{\sum_{i}u_{i}}$$

Rationale:

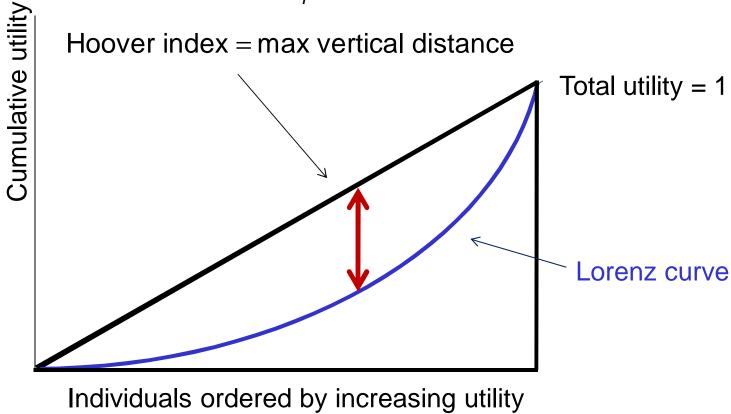
- Fraction of total utility that must be redistributed to achieve total equality.
- Proportional to maximum vertical distance between Lorenz curve and 45° line.
- Originated in regional studies, population distribution, etc. (1930s).
- Easy to calculate.

Problems:

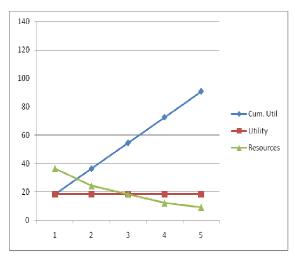
Less informative than Gini coefficient?

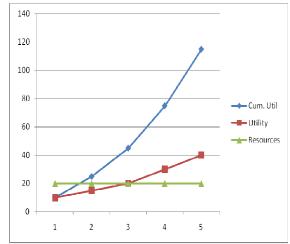
Hoover Index

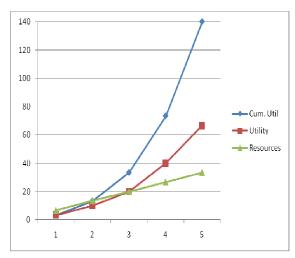
$$\frac{\sum_{i} |u_{i} - \overline{u}|}{\sum_{i} u_{i}}$$



Equality Measures: Comparison

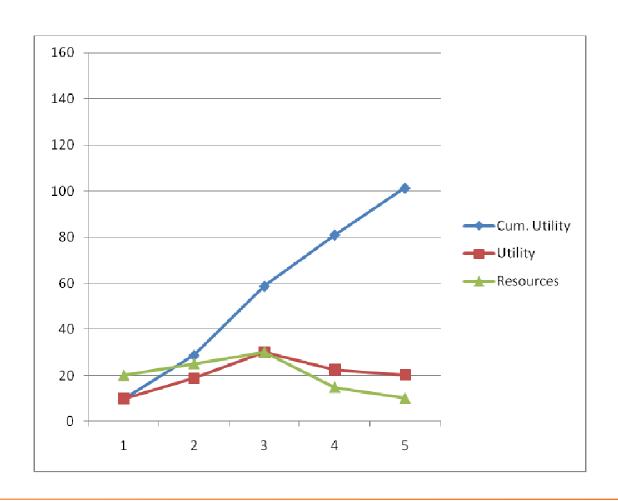




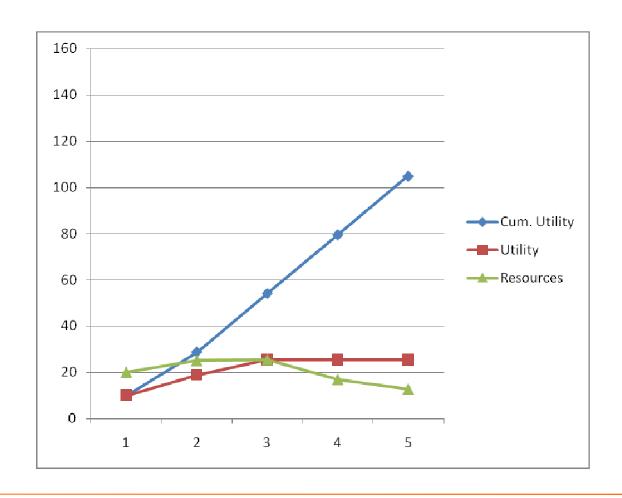


Relative range:	0	1.30	2.26
Rel. mean dev.:	0	0.42	0.72
Coeff. of variation:	0	0.46	0.81
McLoone:	1	0.54	0.23
Gini:	0	0.26	0.45
Atkinson:	0.06	0	0.06
Hoover:	0	0.15	0.28

Hoover Index



Gini Coefficient



Theil Index

$$(1/n)\sum_{i}\left(\frac{u_{i}}{\overline{u}}\ln\frac{u_{i}}{\overline{u}}\right)$$

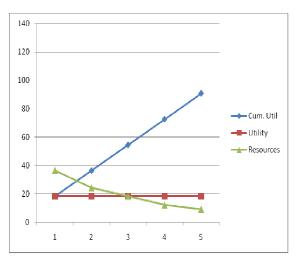
Rationale:

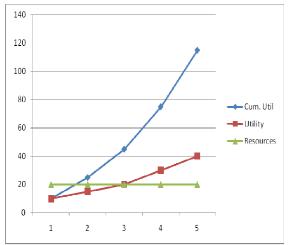
- One of a family of entropy measures of inequality.
- Index is zero for complete equality (maximum entropy)
- Measures nonrandomness of distribution.
- Described as stochastic version of Hoover index.

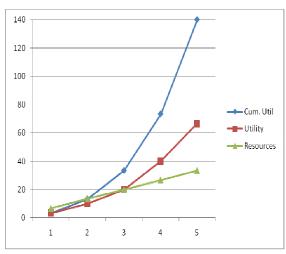
Problems:

- Motivation unclear.
- A. Sen doesn't like it.

Equality Measures: Comparison







Relative range:	0	1.30	2.26
Rel. mean dev.:	0	0.42	0.72
Coeff. of variation:	0	0.46	0.81
McLoone:	1	0.54	0.23
Gini:	0	0.26	0.45
Atkinson:	0.06	0	0.06
Hoover:	0	0.15	0.28
Theil:	0	0.27	0.86

Theil Index

$$(1/n)\sum_{i}\left(\frac{u_{i}}{\overline{u}}\ln\frac{u_{i}}{\overline{u}}\right)$$

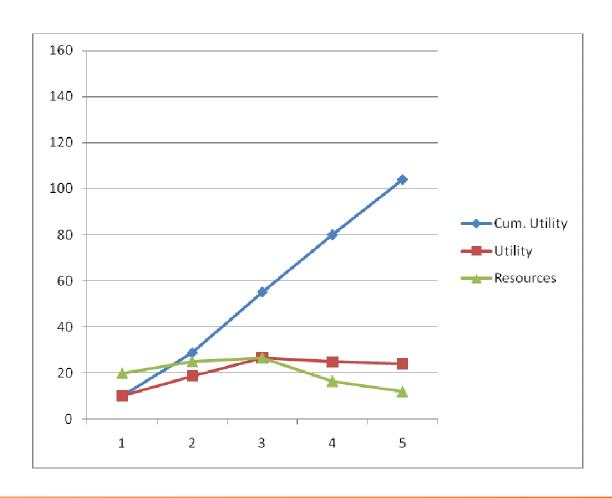
Nasty nonconvex model:

$$\min (1/n) \sum_{i} \left(\frac{u_{i}}{\overline{u}} \ln \frac{u_{i}}{\overline{u}} \right)$$

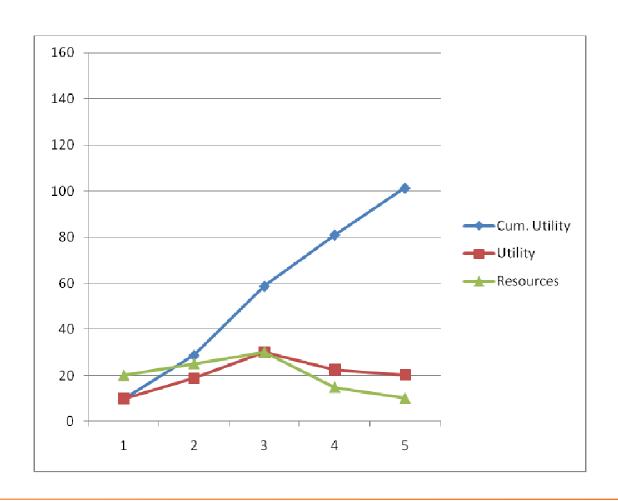
$$\overline{u} = (1/n) \sum_{i} u_{i}$$

$$u_{i} = a_{i} x_{i}, \quad 0 \le x_{i} \le b_{i}, \text{ all } i, \quad \sum_{i} x_{i} = B$$

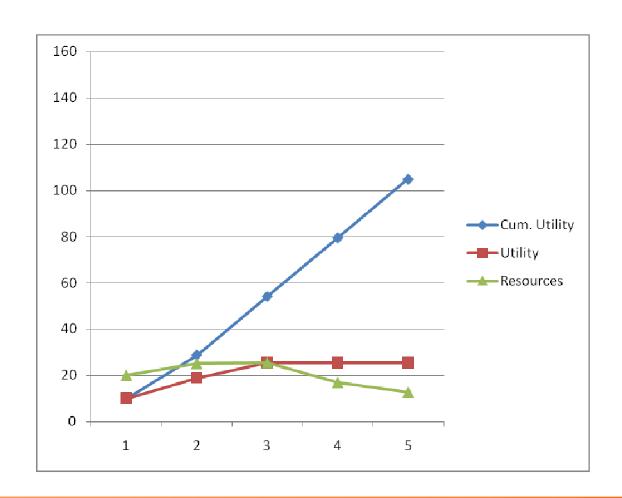
Theil Index



Hoover Index



Gini Coefficient



Outline

- Bargaining:
 - Nash Bargaining Solution
 - Raiffa-Kalai-Smorodinsky Bargaining
 - Disjunctive Modeling
- Combining Equity and Efficiency
 - Health Care Example

An Allocation Problem

- From Yaari and Bar-Hillel, 1983.
- 12 grapefruit and 12 avocados are to be divided between Jones and Smith.
- How to divide justly?

Utility provided by one fruit of each kind

Jones	Smith
100	50
0	50

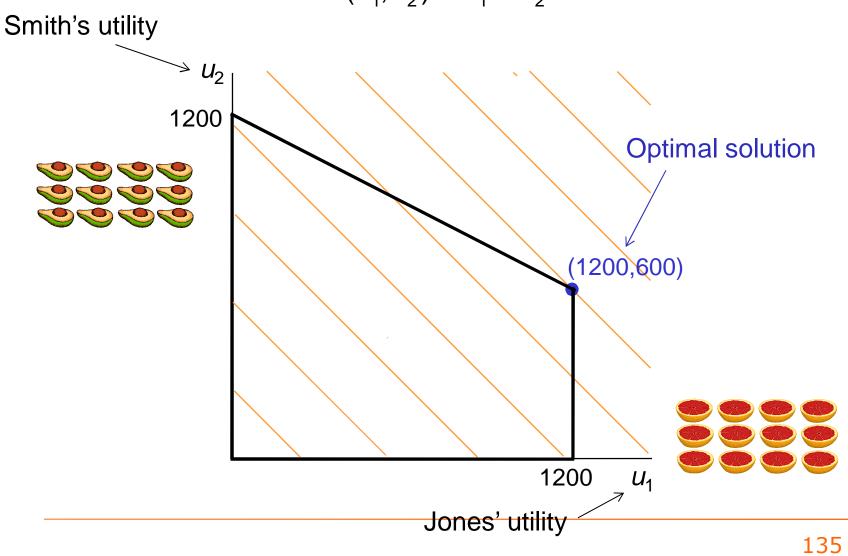
An Allocation Problem

The optimization problem:

```
Social welfare function \max_{u_1} f(u_1, u_2)
u_1 = 100x_{11}, \quad u_2 = 50x_{12} + 50x_{22}
x_{i1} + x_{i2} = 12, \quad i = 1, 2
x_{ij} \ge 0, \quad \text{all } i, j
where u_i = \text{utility for person } i \text{ (Jones, Smith)}
x_{ij} = \text{allocation of fruit } i \text{ (grapefruit, avocados)}
to person j
```

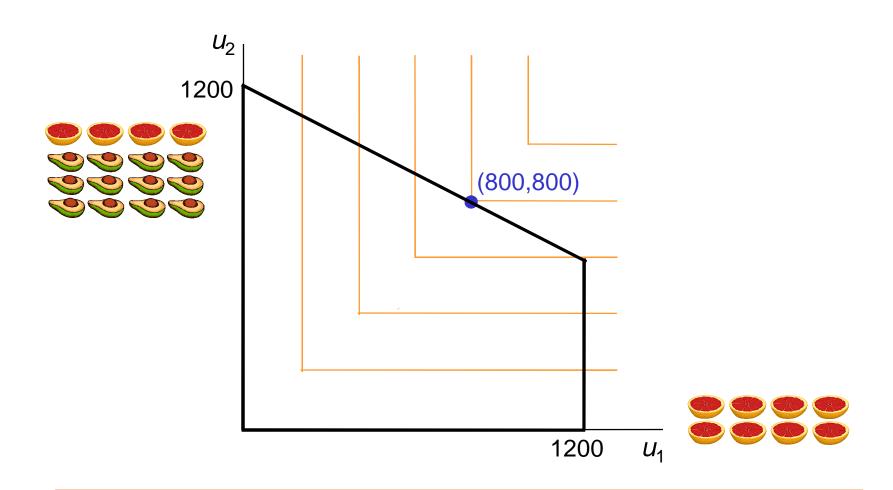
Utilitarian Solution

$$f(u_1, u_2) = u_1 + u_2$$



Rawlsian (maximin) solution

$$f(u_1, u_2) = \min\{u_1, u_2\}$$



Bargaining Solutions

- A bargaining solution is an equilibrium allocation in the sense that none of the parties wish to bargain further.
 - Because all parties are "satisfied" in some sense, the outcome may be viewed as "fair."
 - Bargaining models have a default outcome, which is the result of a failure to reach agreement.
 - The default outcome can be seen as a starting point.

Bargaining Solutions

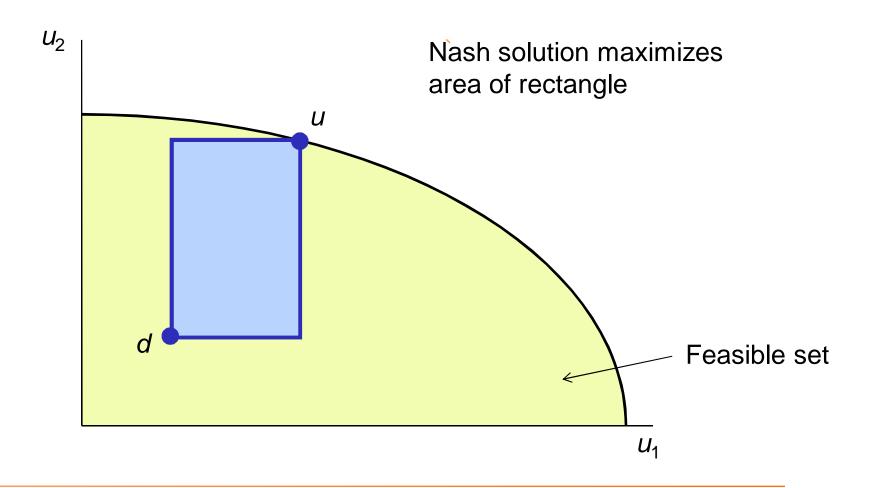
- Several proposals for the default outcome (starting point):
 - **Zero** for everyone. Useful when only the resources being allocated are relevant to fairness of allocation.
 - **Equal split**. Resources (not necessarily utilities) are divided equally. May be regarded as a "fair" **starting point**.
 - Strongly pareto set. Each party receives resources that can benefit no one else. Parties can always agree on this.

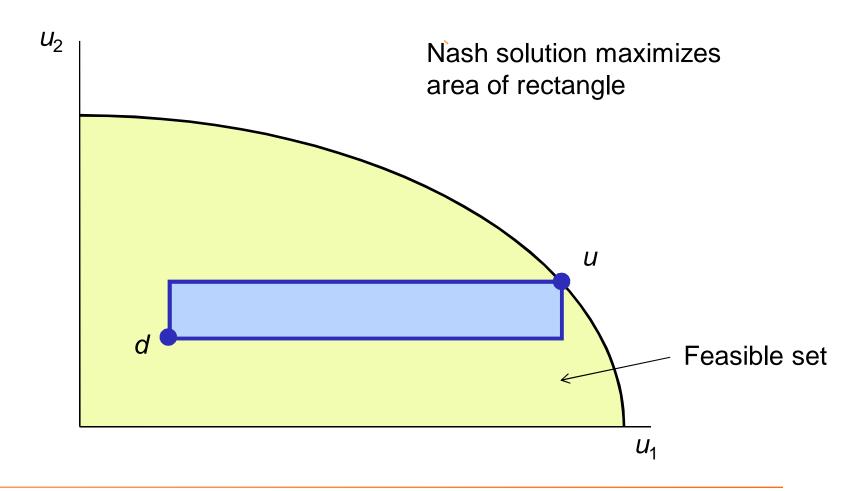
The Nash bargaining solution maximizes the social welfare function

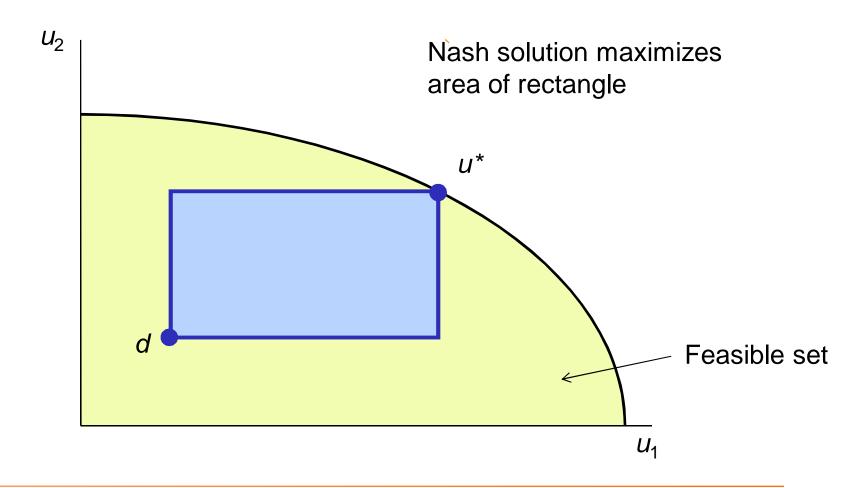
$$f(u) = \prod_{i} (u_i - d_i)$$

where d is the default outcome.

- Not the same as Nash equilibrium.
- It maximizes the **product of the gains** achieved by the bargainers, relative to the fallback position.
- Assume feasible set is **convex**, so that Nash solution is unique (due to strict concavity of *f*).







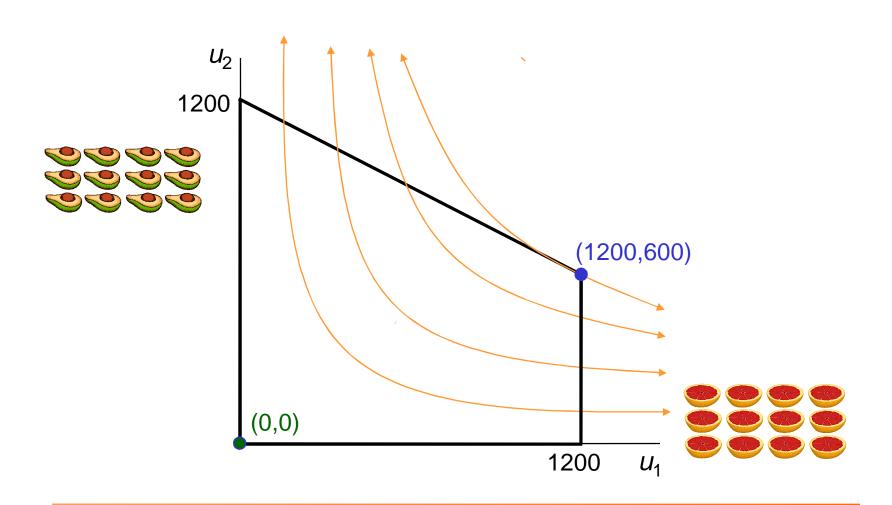
• The **optimization problem** has a concave objective function if we maximize $\log f(u)$.

$$\max \log \prod_{i} (u_{i} - d_{i}) = \sum_{i} \log(u_{i} - d_{i})$$

$$u \in S$$

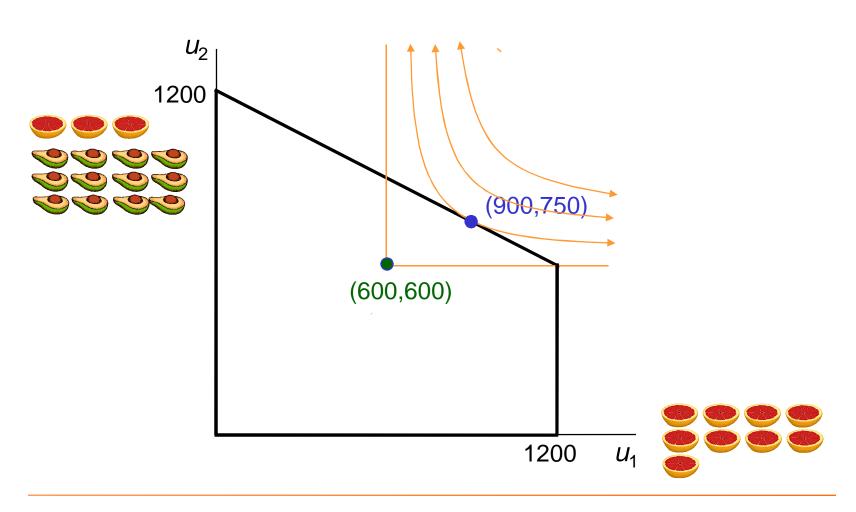
Problem is relatively easy if feasible set S is convex.

From Zero



Nash Bargaining Solution

From Equality



Nash Bargaining Solution

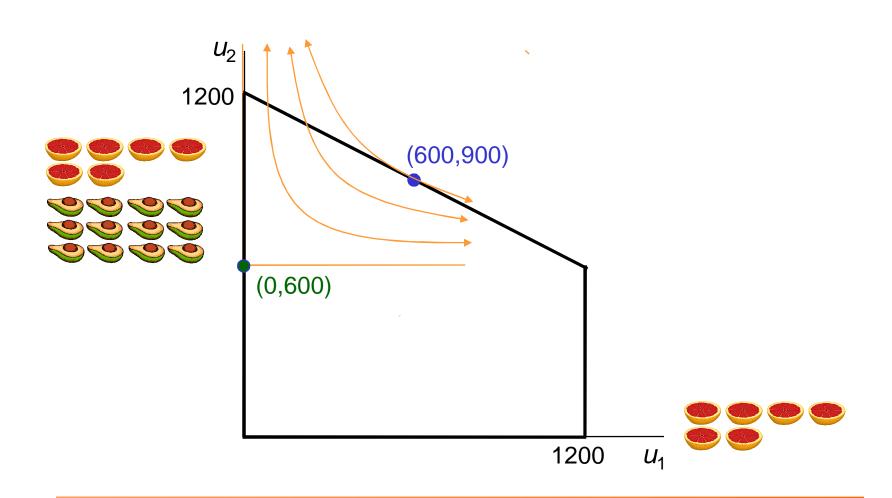
- Strongly pareto set gives Smith all 12 avocados.
 - Nothing for Jones.
 - Results in utility $(u_1, u_2) = (0,600)$

Utility provided by one fruit of each kind

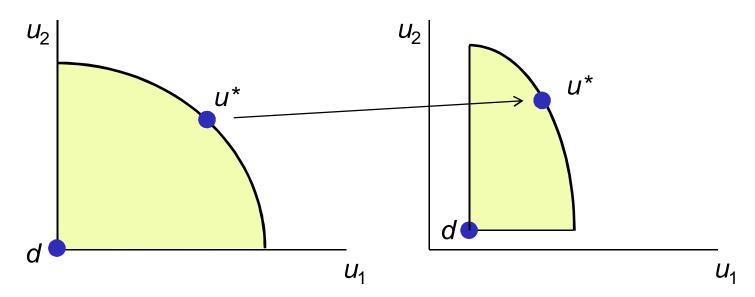
Jones	Smith
100	50
0	50

Nash Bargaining Solution

From Strongly Pareto Set

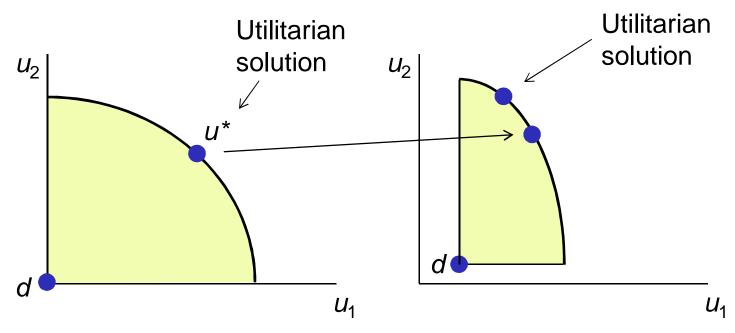


- Axiom 1. Invariance under translation and rescaling.
 - If we map $u_i \rightarrow a_i u_i + b_i$, $d_i \rightarrow a_i d_i + b_i$, then bargaining solution $u_i^* \rightarrow a_i u_i^* + b_i$.



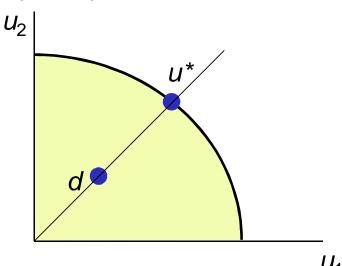
This is cardinal noncomparability.

- Axiom 1. Invariance under translation and rescaling.
 - If we map $u_i \rightarrow a_i u_i + b_i$, $d_i \rightarrow a_i d_i + b_i$, then bargaining solution $u_i^* \rightarrow a_i u_i^* + b_i$.

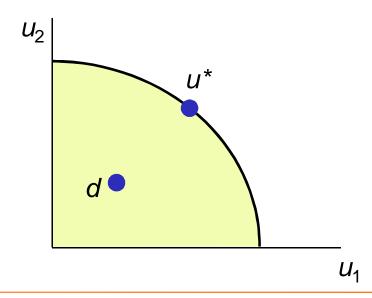


• Strong assumption – failed, e.g., by utilitarian welfare function

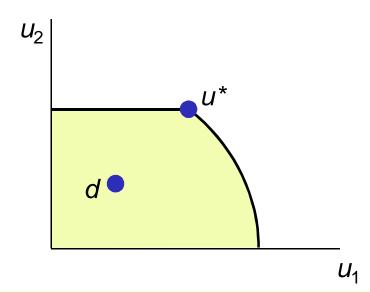
- Axiom 2. Pareto optimality.
 - Bargaining solution is pareto optimal.
- Axiom 3. Symmetry.
 - If all d_i s are equal and feasible set is symmetric, then all u_i^* s are equal in bargaining solution.



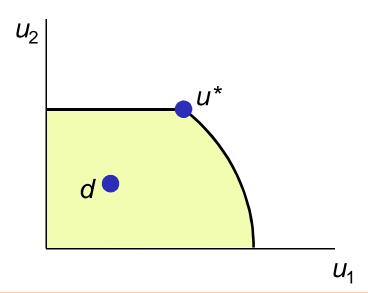
- Axiom 4. Independence of irrelevant alternatives.
 - Not the same as Arrow's axiom.
 - If *u** is a solution with respect to *d*…



- Axiom 4. Independence of irrelevant alternatives.
 - Not the same as Arrow's axiom.
 - If u^* is a solution with respect to d, then it is a solution in a smaller feasible set that contains u^* and d.



- Axiom 4. Independence of irrelevant alternatives.
 - Not the same as Arrow's axiom.
 - If u^* is a solution with respect to d, then it is a solution in a smaller feasible set that contains u^* and d.
 - This basically says that the solution behaves like an optimum.



Theorem. Exactly one solution satisfies Axioms 1-4, namely the Nash bargaining solution.

Proof (2 dimensions).

First show that the Nash solution satisfies the axioms.

Axiom 1. Invariance under transformation. If

$$\prod_{i} (u_{i}^{*} - d_{1}) \ge \prod_{i} (u_{i} - d_{1})$$
then
$$\prod_{i} ((a_{i}u_{i}^{*} + b_{i}) - (a_{i}d_{i} + b_{i})) \ge \prod_{i} ((a_{i}u_{i} + b_{i}) - (a_{i}d_{i} + b_{i}))$$

Axiom 2. Pareto optimality. Clear because social welfare function is strictly monotone increasing.

Axiom 3. Symmetry. Obvious.

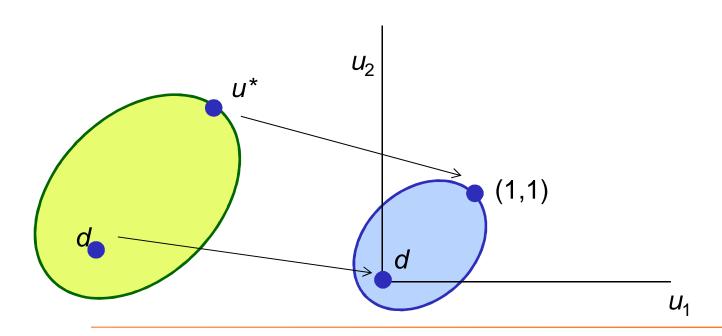
Axiom 4. Independence of irrelevant alternatives. Follows from the fact that u^* is an optimum.

Now show that **only** the Nash solution satisfies the axioms...

Let *u** be the Nash solution for a given problem. Then it satisfies the axioms with respect to *d*. Select a transformation that sends

$$(u_1, u_2) \rightarrow (1,1), \quad (d_1, d_2) \rightarrow (0,0)$$

The transformed problem has Nash solution (1,1), by Axiom 1:

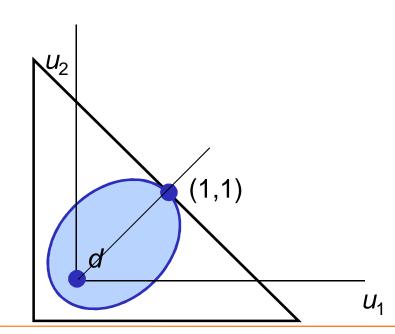


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By Axioms 2 & 3, (1,1) is the **only** bargaining solution in the triangle:

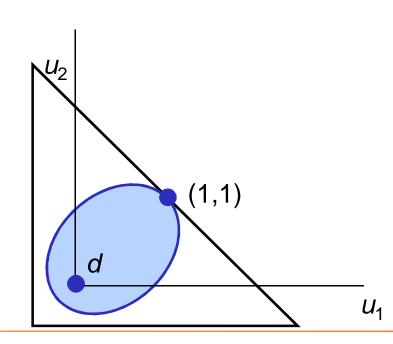


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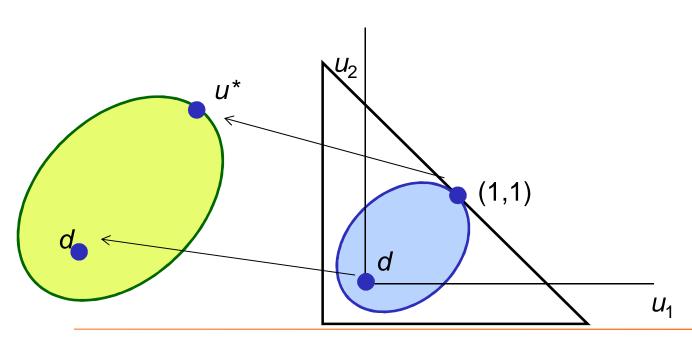


So by Axiom 4, (1,1) is the only bargaining solution in blue set.

Let *u** be the Nash solution for a given problem. Then it satisfies the axioms with respect to *d*. Select a transformation that sends

$$(u_1, u_2) \rightarrow (1,1), \quad (d_1, d_2) \rightarrow (0,0)$$

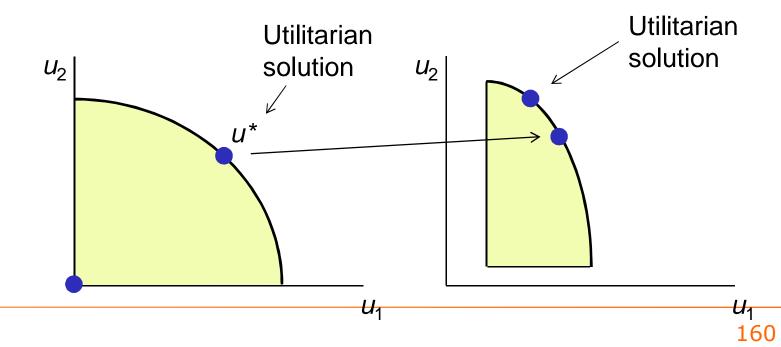
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So by Axiom 4, (1,1) is the only bargaining solution in blue set.

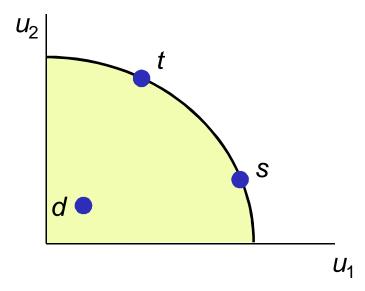
By Axiom 1, *u** is the only bargaining solution in the original problem.

- Problems with axiomatic justification.
 - Axiom 1 (invariance under transformation) is very strong.
 - Axiom 1 denies interpersonal comparability.
 - So how can it reflect moral concerns?



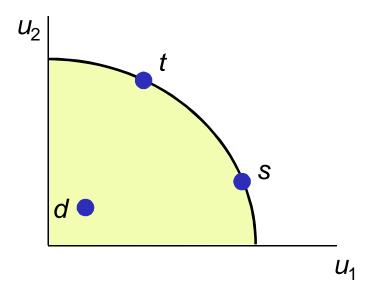
- Problems with axiomatic justification.
 - Axiom 1 (invariance under transformation) is very strong.
 - Axiom 1 denies interpersonal comparability.
 - So how can it reflect moral concerns?
- Most attention has been focused on Axiom 4 (independence of irrelevant alternatives).
 - Will address this later.

Players 1 and 2 make offers s, t.



Players 1 and 2 make offers s, t.

Let p = P(player 2 will reject s), as estimated by player 1.

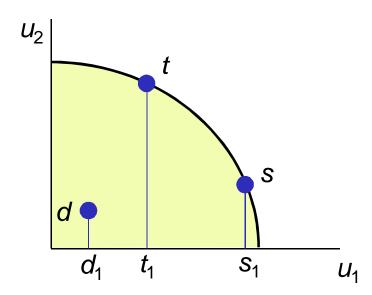


Players 1 and 2 make offers s, t.

Let p = P(player 2 will reject s), as estimated by player 1.

Then player 1 will stick with s, rather than make a counteroffer, if

$$(1-p)s_1+pd_1\geq t_1$$

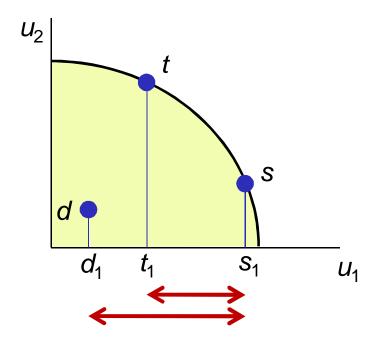


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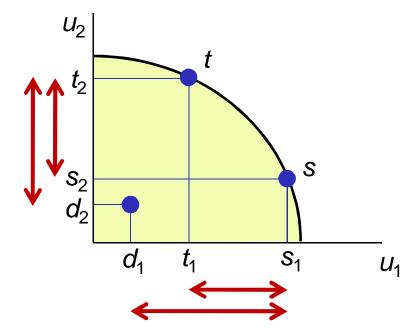


So player 1 will stick with s if

$$p \le \frac{s_1 - t_1}{s_1 - d_1} = r_1$$

It is rational for player 1 to make a counteroffer s', rather than player 2, if

$$r_1 = \frac{s_1 - t_1}{s_1 - d_1} \le \frac{t_2 - s_2}{t_2 - d_2} = r_2$$



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$$r_1 = \frac{s_1 - t_1}{s_1 - d_1} \le \frac{t_2 - s_2}{t_2 - d_2} = r_2$$

 u_2

It is rational for player 2 to make the next counteroffer if

$$r_1' = \frac{s_1' - t_1}{s_1' - d_1} \ge \frac{t_2 - s_2'}{t_2 - d_2} = r_2'$$

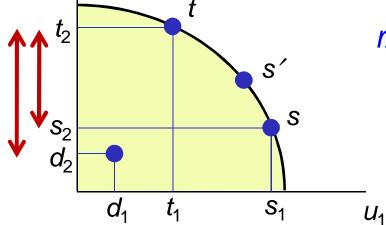
 U_1

It is rational for player 1 to make a counteroffer s', rather than player 2, if

$$r_1 = \frac{s_1 - t_1}{s_1 - d_1} \le \frac{t_2 - s_2}{t_2 - d_2} = r_2$$

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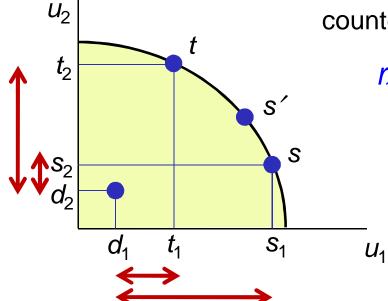
 U_2

But
$$\frac{S_1 - t_1}{S_1 - d_1} \le \frac{t_2 - S_2}{t_2 - d_2}$$

It is rational for player 1 to make a counteroffer s', rather than player 2, if

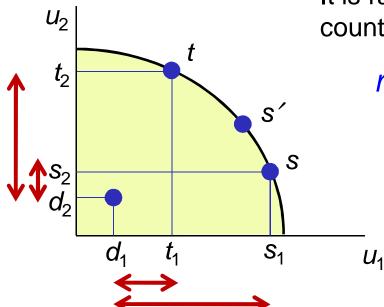
$$r_1 = \frac{s_1 - t_1}{s_1 - d_1} \le \frac{t_2 - s_2}{t_2 - d_2} = r_2$$

$$r_1' = \frac{s_1' - t_1}{s_1' - d_1} \ge \frac{t_2 - s_2'}{t_2 - d_2} = r_2'$$



But
$$\frac{s_1 - t_1}{s_1 - d_1} \le \frac{t_2 - s_2}{t_2 - d_2}$$
 $\iff \frac{t_1 - d_1}{s_1 - d_1} \ge \frac{s_2 - d_2}{t_2 - d_2}$

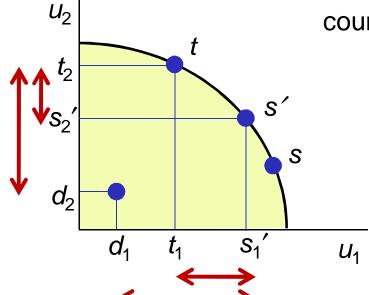
So we have
$$(s_1 - d_1)(s_2 - d_2) \le (t_1 - d_1)(t_2 - d_2)$$



$$r_1' = \frac{s_1' - t_1}{s_1' - d_1} \ge \frac{t_2 - s_2'}{t_2 - d_2} = r_2'$$

But
$$\frac{s_1 - t_1}{s_1 - d_1} \le \frac{t_2 - s_2}{t_2 - d_2}$$
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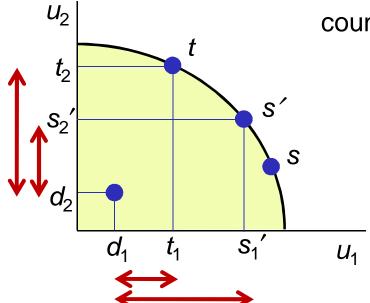
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$$r_1' = \frac{s_1' - t_1}{s_1' - d_1} \ge \frac{t_2 - s_2'}{t_2 - d_2} = r_2'$$

Similarly
$$\frac{s_1' - t_1}{s_1' - d_1} \ge \frac{t_2 - s_2'}{t_2 - d_2'}$$

So we have
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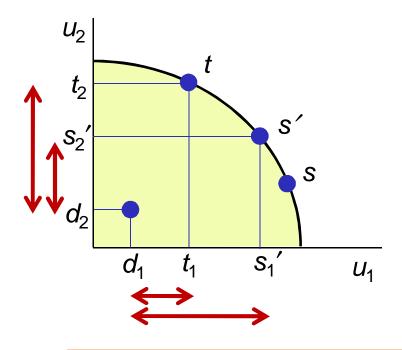
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$$\frac{s_1' - t_1}{s_1' - d_1} \ge \frac{t_2 - s_2'}{t_2 - d_2}$$

$$\iff \frac{t_1 - d_1}{s_1' - d_1} \le \frac{s_2' - d_2}{t_2 - d_2}$$

So we have
$$(s_1 - d_1)(s_2 - d_2) \le (t_1 - d_1)(t_2 - d_2)$$

and we have $(t_1 - d_1)(t_2 - d_2) \le (s_1' - d_1)(s_2' - d_2)$

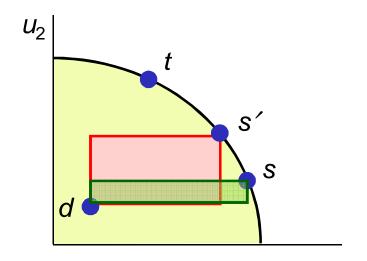


Similarly
$$\frac{s'_1 - t_1}{s'_1 - d_1} \ge \frac{t_2 - s'_2}{t_2 - d_2}$$

$$\iff \frac{t_1 - d_1}{s'_1 - d_1} \le \frac{s'_2 - d_2}{t_2 - d_2}$$

So we have
$$(s_1 - d_1)(s_2 - d_2) \le (t_1 - d_1)(t_2 - d_2)$$

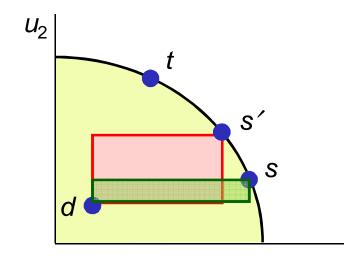
and we have $(t_1 - d_1)(t_2 - d_2) \le (s_1' - d_1)(s_2' - d_2)$



This implies an improvement in the Nash social welfare function

So we have
$$(s_1 - d_1)(s_2 - d_2) \le (t_1 - d_1)(t_2 - d_2)$$

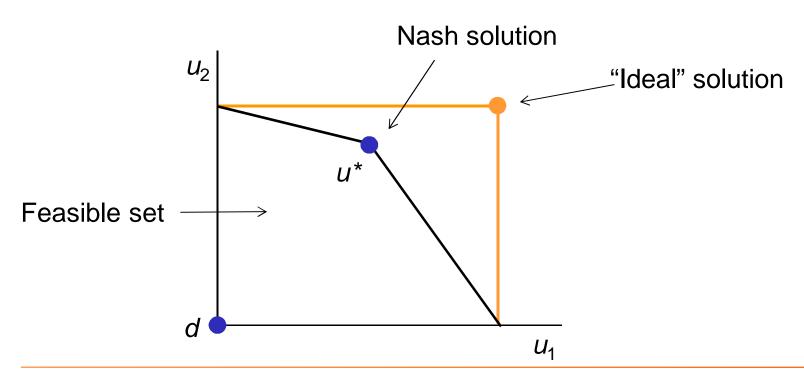
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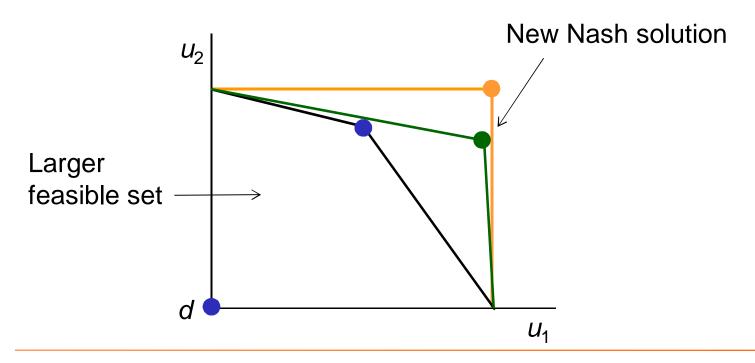
This implies an improvement in the Nash social welfare function.

Given a minimum distance between offers, continued bargaining converges to Nash solution.

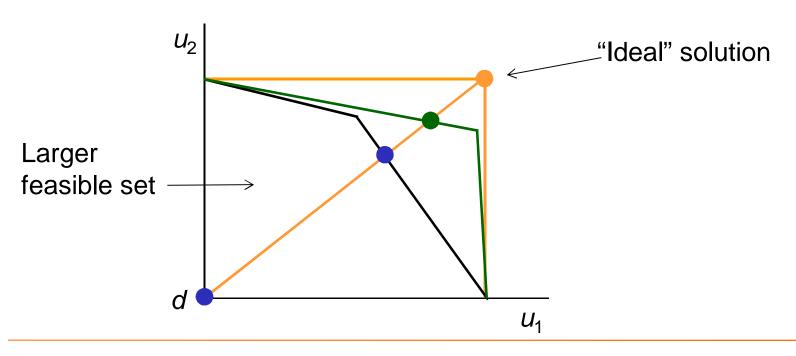
 This approach begins with a critique of the Nash bargaining solution.



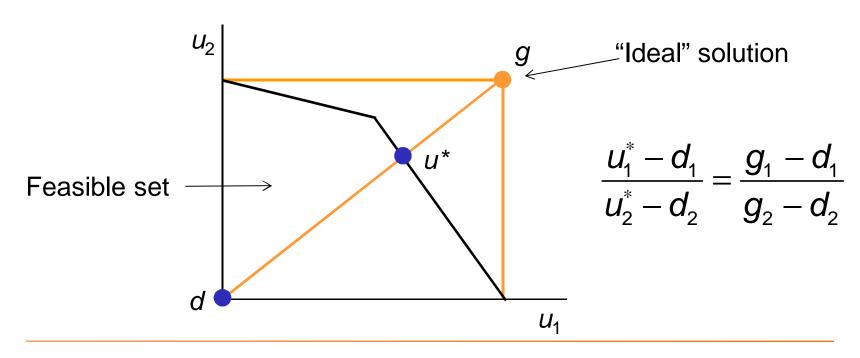
- This approach begins with a critique of the Nash bargaining solution.
 - The new Nash solution is worse for player 2 even though the feasible set is larger.



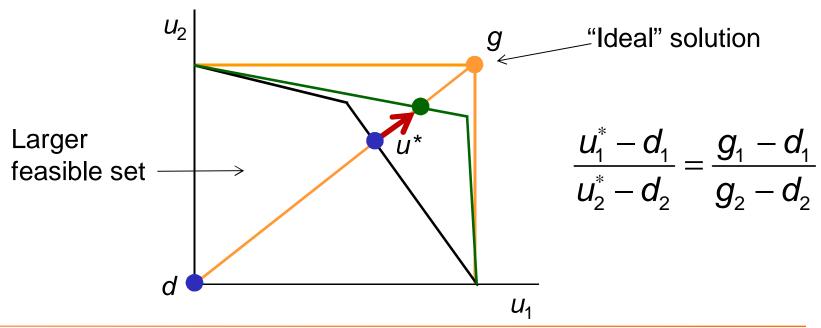
• **Proposal**: Bargaining solution is pareto optimal point on line from *d* to ideal solution.



- Proposal: Bargaining solution is pareto optimal point on line from d to ideal solution.
 - The players receive an equal fraction of their possible utility gains.



- Proposal: Bargaining solution is pareto optimal point on line from d to ideal solution.
 - Replace Axiom 4 with Axiom 4' (Monotonicity): A larger feasible set with same ideal solution results in a bargaining solution that is better (or no worse) for all players.



Optimization model.

- Not an optimization problem over original feasible set (we gave up Axiom 4).
- But it is an optimization problem (pareto optimality) over the line segment from *d* to ideal solution.

$$\max \sum_{i} u_{i}$$

$$(g_{1} - d_{1})(u_{i} - d_{i}) = (g_{i} - d_{i})(u_{1} - d_{1}), \text{ all } i$$

$$u \in S$$

$$\frac{u_{1}^{*} - d_{1}}{u_{2}^{*} - d_{2}} = \frac{g_{1} - d_{1}}{g_{2} - d_{2}}$$

Optimization model.

- Not an optimization problem over original feasible set (we gave up Axiom 4).
- But it is an optimization problem (pareto optimality) over the line segment from *d* to ideal solution.

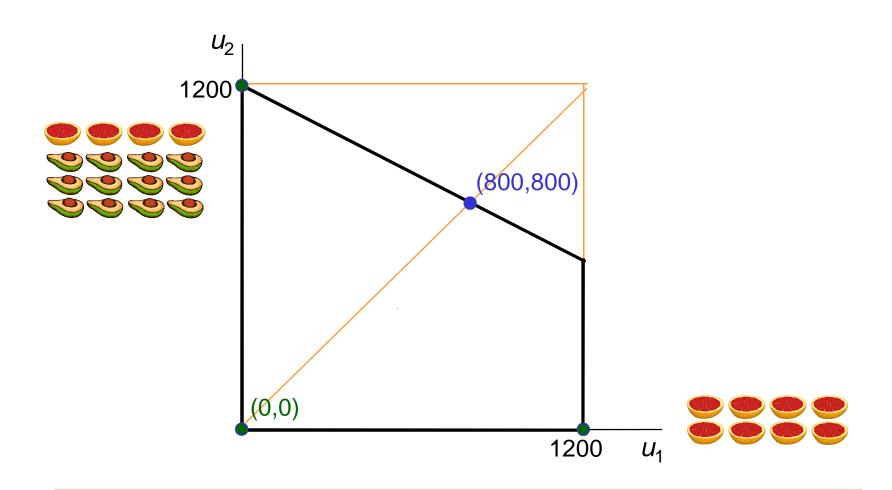
constants
$$\max_{i} \sum_{i} u_{i}$$

$$(g_{1} - d_{1})(u_{i} - d_{i}) = (g_{i} - d_{i})(u_{1} - d_{1}), \text{ all } i$$

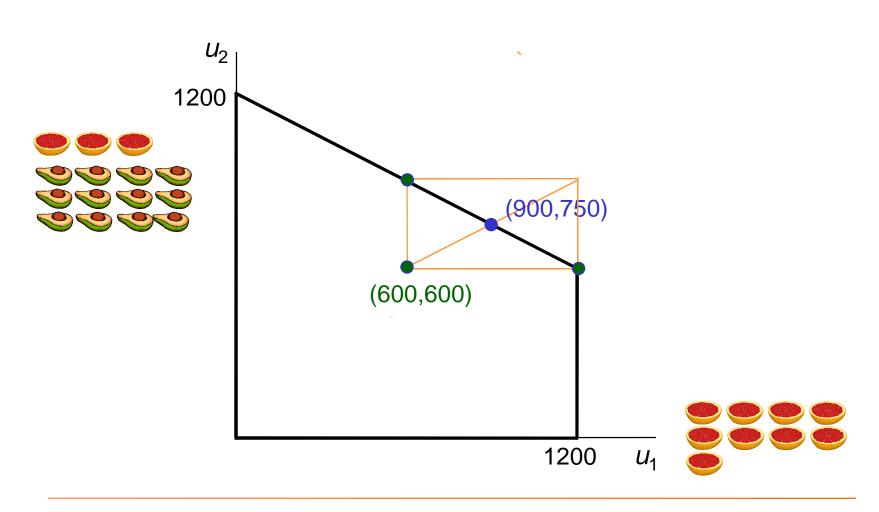
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Optimization model.

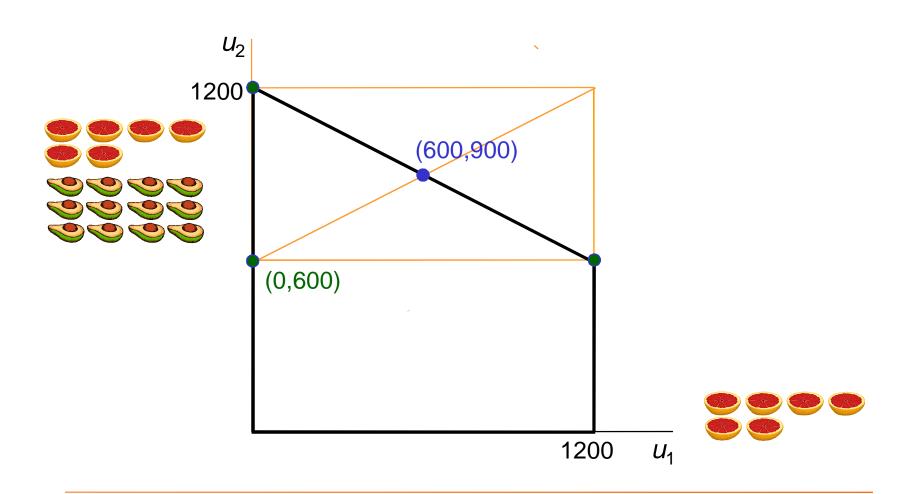
- Not an optimization problem over original feasible set (we gave up Axiom 4).
- But it is an optimization problem (pareto optimality) over the line segment from *d* to ideal solution.



From Equality



From Strong Pareto Set



- Axiom 1. Invariance under transformation.
- Axiom 2. Pareto optimality.
- Axiom 3. Symmetry.
- Axiom 4'. Monotonicity.

Theorem. Exactly one solution satisfies Axioms 1-4', namely the RKS bargaining solution.

Proof (2 dimensions).

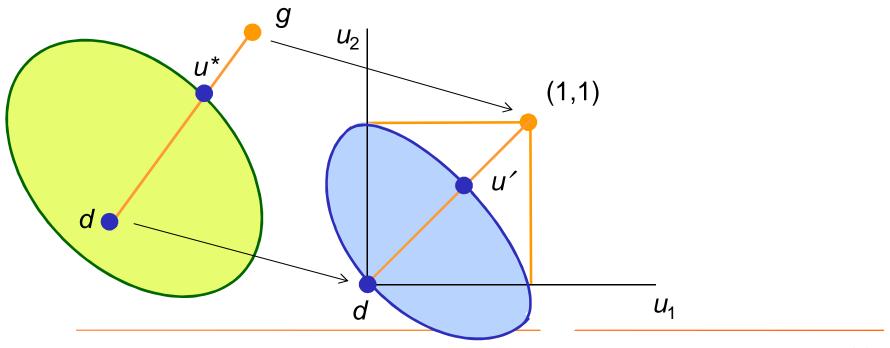
Easy to show that RKS solution satisfies the axioms.

Now show that **only** the RKS solution satisfies the axioms.

Let *u** be the RKS solution for a given problem. Then it satisfies the axioms with respect to *d*. Select a transformation that sends

$$(g_1,g_2) \rightarrow (1,1), \quad (d_1,d_2) \rightarrow (0,0)$$

The transformed problem has RKS solution u', by Axiom 1:

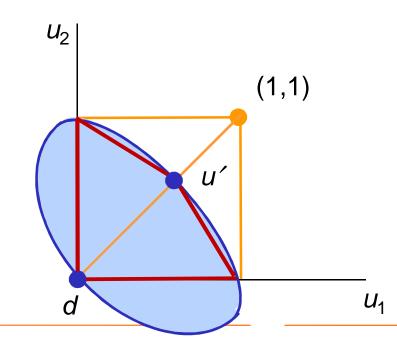


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The transformed problem has RKS solution u', by Axiom 1:

By Axioms 2 & 3, u' is the **only** bargaining solution in the red polygon:

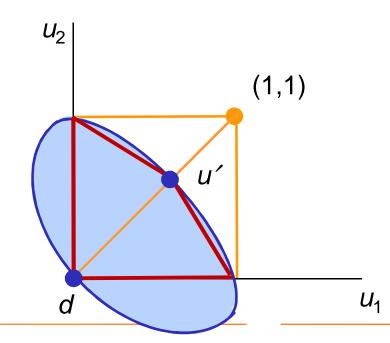


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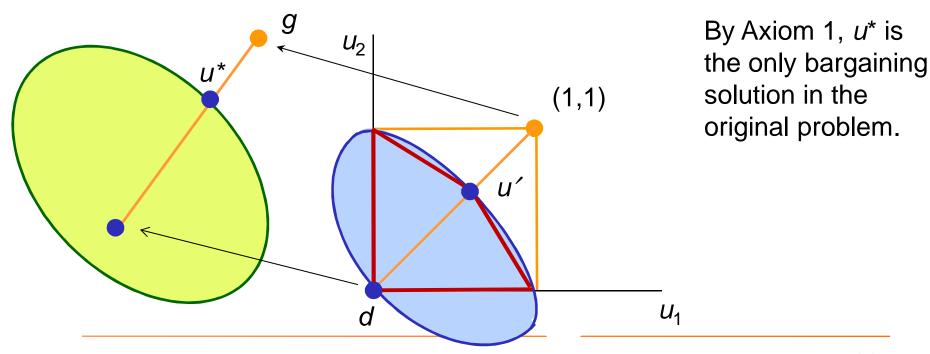


The red polygon lies inside blue set. So by Axiom 4', its bargaining solution is no better than bargaining solution on blue set. So u' is the only bargaining solution on blue set.

Let *u** be the RKS solution for a given problem. Then it satisfies the axioms with respect to *d*. Select a transformation that sends

$$(g_1,g_2) \rightarrow (1,1), \quad (d_1,d_2) \rightarrow (0,0)$$

The transformed problem has RKS solution u', by Axiom 1:



- Problems with axiomatic justification.
 - Axiom 1 is still in effect.
 - It denies interpersonal comparability.
 - Dropping Axiom 4 sacrifices optimization of a social welfare function.
 - This may not be necessary if Axiom 1 is rejected.
 - Needs modification for > 2 players (more on this shortly).

Resistance to an agreement s depends on sacrifice relative to sacrifice under no agreement. Here, player 2 is making a larger

relative sacrifice:

$$\frac{g_1 - s_1}{g_1 - d_1} \le \frac{g_2 - s_2}{g_2 - d_2}$$

 $\begin{array}{c} g_2 \\ g_2 \\ g_2 \\ d_1 \\ g_1 \\ g_1 \\ g_1 \\ g_1 \\ g_1 \\ g_2 \\ g_2 \\ g_3 \\ g_4 \\ g_4 \\ g_5 \\ g_6 \\ g_7 \\ g_8 \\$

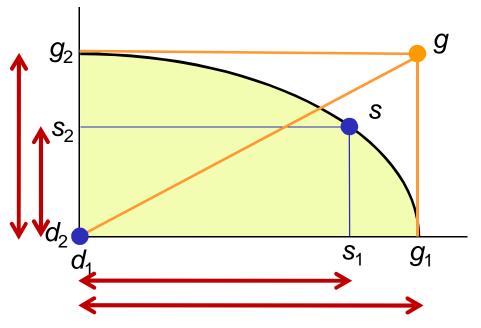
Minimizing resistance to agreement requires minimizing

$$\max_{i} \left\{ \frac{g_{i} - s_{i}}{g_{i} - d_{i}} \right\}$$

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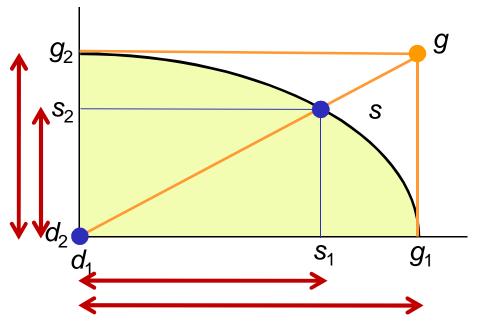
or equivalently, maximizing

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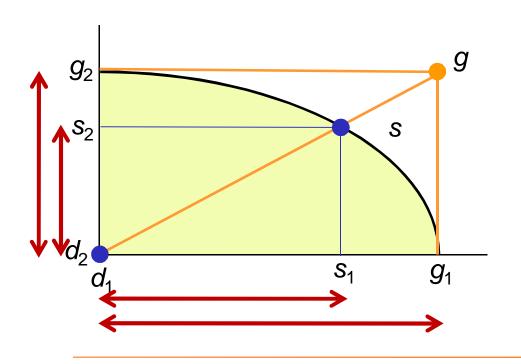
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which is achieved by RKS point.

This is the **Rawlsian social contract** argument applied to **gains** relative to the ideal.



Minimizing resistance to agreement requires minimizing

$$\max_{i} \left\{ \frac{g_{i} - s_{i}}{g_{i} - d_{i}} \right\}$$

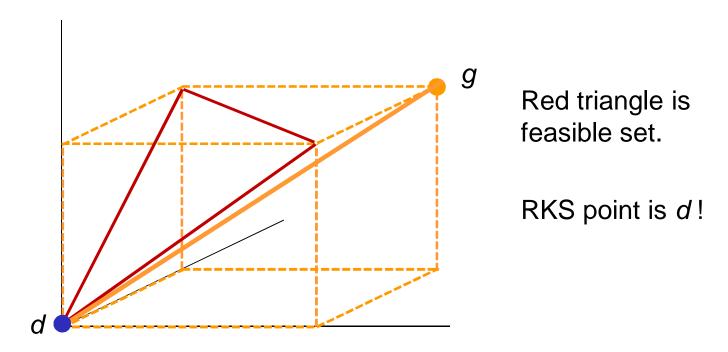
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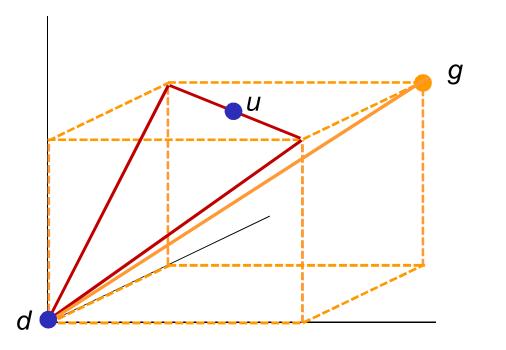
Problem with RKS Solutioon

- However, the RKS solution is Rawlsian only for 2 players.
 - In fact, RKS leads to counterintuitive results for 3 players.



Problem with RKS Solutioon

- However, the RKS solution is Rawlsian only for 2 players.
 - In fact, KLS leads to counterintuitive results for 3 players.

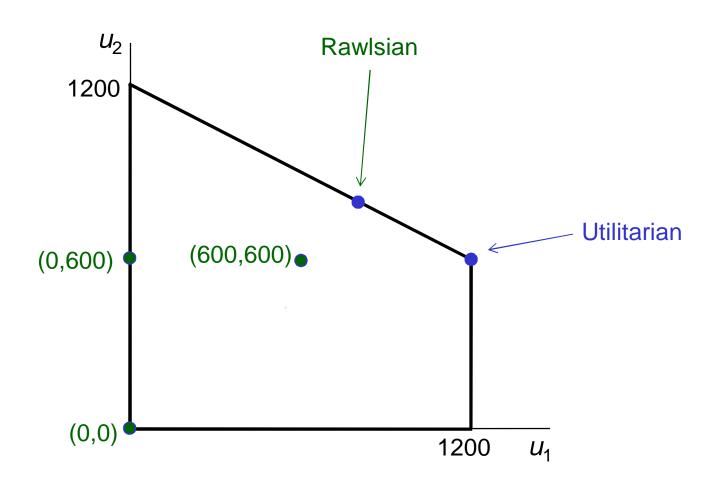


Red triangle is feasible set.

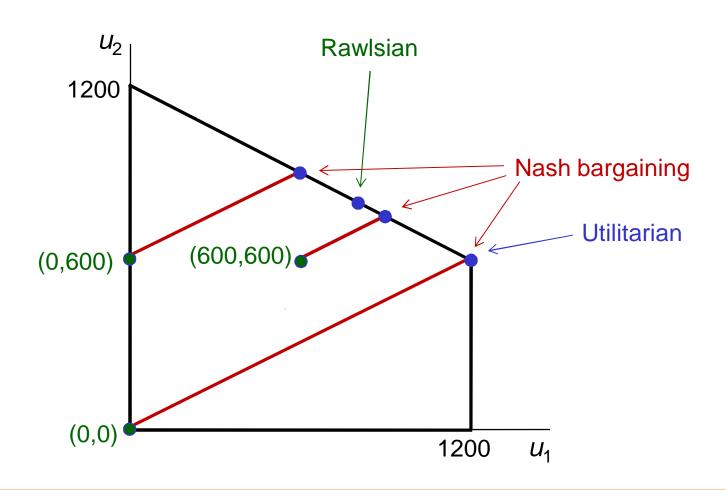
RKS point is *d*!

Rawlsian point is u.

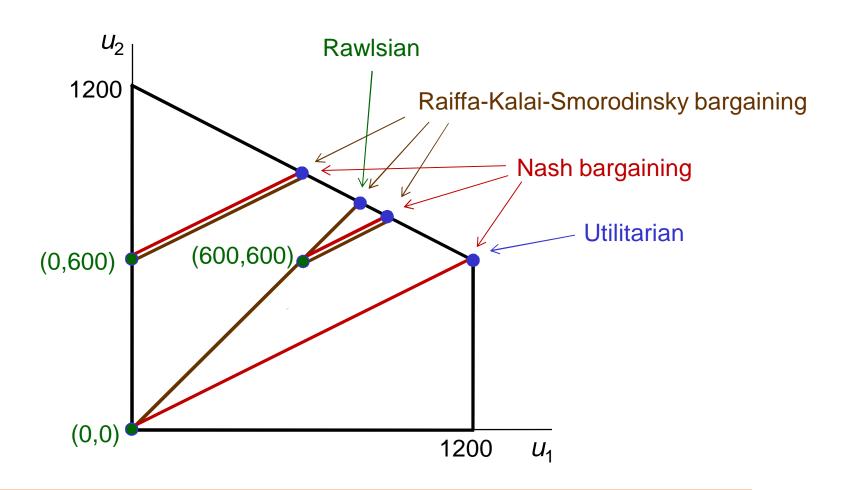
Summary



Summary



Summary



Mixed Integer Linear Modeling

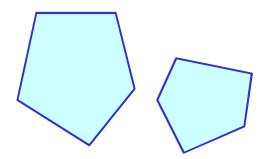
- MILP modeling is basically disjunctive modeling.
- A problem has an MILP model if and only if it represents a union of polyhedra with the same recession cone.
- One can always write an MILP model by expressing the problem as a disjunction of linear systems that describe polyhedra with the same recession cone.
- In fact, one can write a convex hull (sharp) MILP model in this fashion.

Disjunctions of linear systems

A disjunction of linear systems represents a union of polyhedra.

$$\min cx$$

$$\bigvee_{k} (A^{k} x \ge b^{k})$$



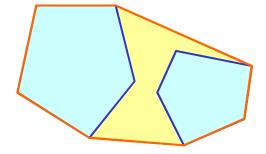
Disjunction of linear systems

A disjunction of linear systems represents a union of polyhedra.

We want a model with a convex hull relaxation (tightest linear relaxation).

$$\min cx$$

$$\bigvee_{k} (A^{k} x \ge b^{k})$$



Disjunction of linear systems

The closure of the convex hull of

$$\bigvee_{k} (A^{k} x \geq b^{k})$$

$$A^k x^k \ge b^k y_k$$
, all k

$$\sum_{k} y_{k} = 1$$

$$X = \sum_{k} X^{k}$$

$$0 \le y_k \le 1$$

Why?

To derive convex hull relaxation of a disjunction...

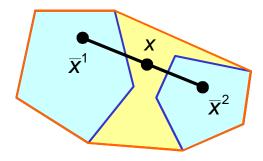
write each solution as a convex combination of points in the polyhedron
$$\min cx$$

$$A^k \overline{x}^k \ge b^k, \text{ all } k$$

$$\sum_k y_k = 1$$

$$x = \sum_k y_k \overline{x}^k$$

$$0 \le y_k \le 1$$



Convex hull relaxation (tightest linear relaxation)

Why?

To derive convex hull relaxation of a disjunction...

Write each solution as a convex combination of points in the

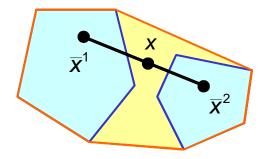
polyhedron

min cx $A^{k}\overline{x}^{k} \ge b^{k}, \text{ all } k$ $\sum_{k} y_{k} = 1$ $x = \sum_{k} y_{k} \overline{x}^{k}$

 $0 \le y_k \le 1$

Change of variable $X = Y_k \overline{X}^k$

min cx $A^{k}x^{k} \ge b^{k}y_{k}, \text{ all } k$ $\sum_{k} y_{k} = 1$ $x = \sum_{k} x^{k}$ $0 \le y_{k} \le 1$



Convex hull relaxation (tightest linear relaxation)

MILP Representability

A subset S of \mathbb{R}^n is MILP representable if it is the projection onto x of some MILP constraint set of the form

$$Ax + Bu + Dy \ge b$$

$$x, y \ge 0$$

$$x \in \mathbb{R}^n, \ u \in \mathbb{R}^m, \ y_k \in \{0,1\}$$

MILP Representability

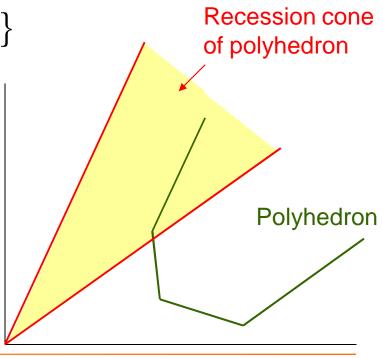
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$$Ax + Bu + Dy \ge b$$

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$$x \in \mathbb{R}^n$$
, $u \in \mathbb{R}^m$, $y_k \in \{0,1\}$

Theorem. $S \subset \mathbb{R}^n$ is MILP representable if and only if S is the union of finitely many polyhedra having the same recession cone.

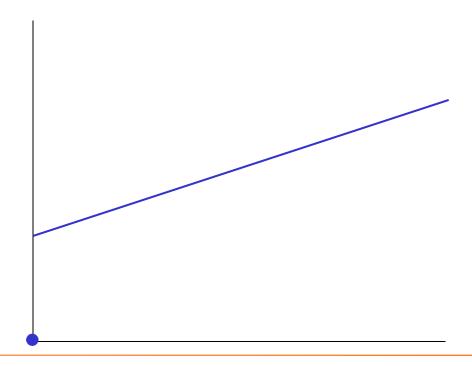


Example: Fixed charge function

$$\min x_2$$

$$x_2 \ge \begin{cases} 0 & \text{if } x_1 = 0 \\ f + cx_1 & \text{if } x_1 > 0 \end{cases}$$

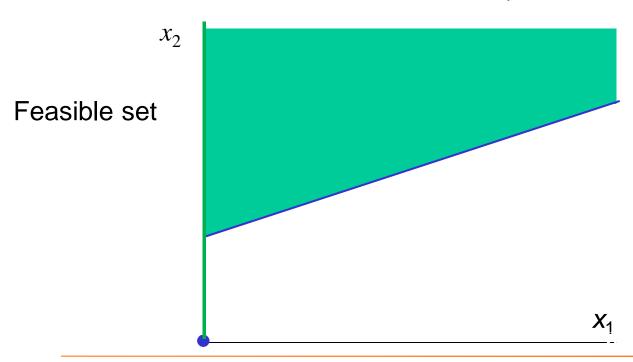
$$x_1 \ge 0$$



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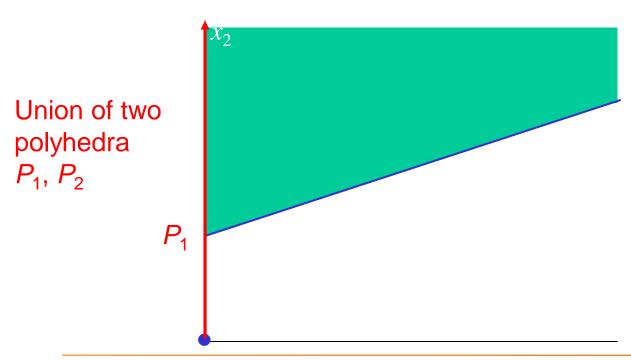
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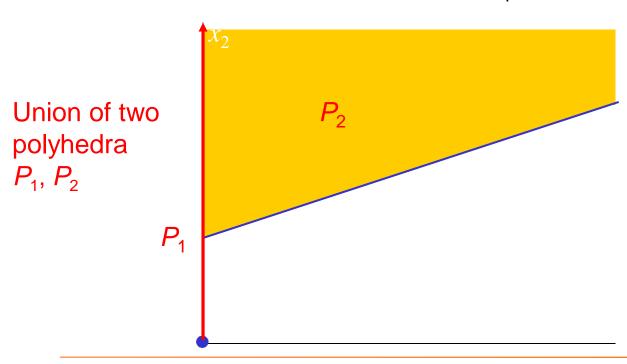
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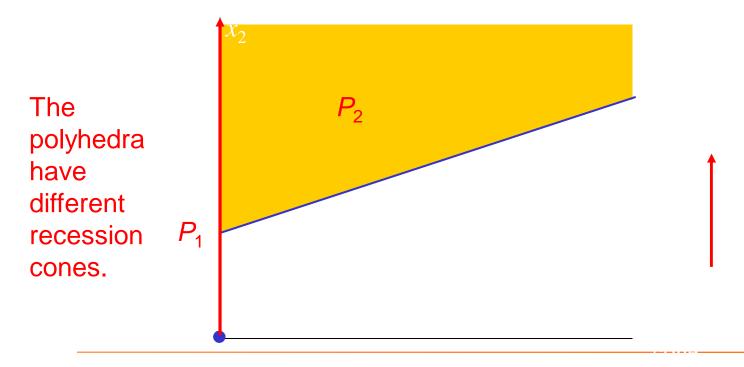
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$$\min x_2$$

$$x_2 \ge \begin{cases} 0 & \text{if } x_1 = 0 \\ f + cx_1 & \text{if } x_1 > 0 \end{cases}$$

$$x_1 \ge 0$$



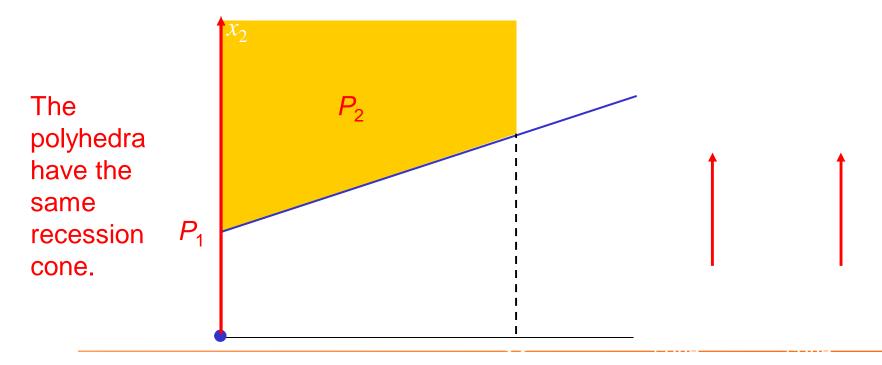
Minimize a fixed charge function:

Add an upper bound on x_1

$$\min x_2$$

$$x_2 \ge \begin{cases} 0 & \text{if } x_1 = 0 \\ f + cx_1 & \text{if } x_1 > 0 \end{cases}$$

$$0 \le x_1 \le M$$



Modeling a union of polyhedra

Start with a disjunction of linear systems to represent the union of polyhedra.

The *k*th polyhedron is $\{x \mid A^k x \ge b\}$

Introduce a 0-1 variable y_k that is 1 when x is in polyhedron \underline{k} .

Disaggregate x to create an x^k for each k.

min *cx*

$$\bigvee_{k} (A^{k} x \geq b^{k})$$

min cx

$$A^k x^k \ge b^k y_k$$
, all k
 $\sum_k y_k = 1$

$$\mathbf{x} = \sum_{k} \mathbf{x}^{k}$$

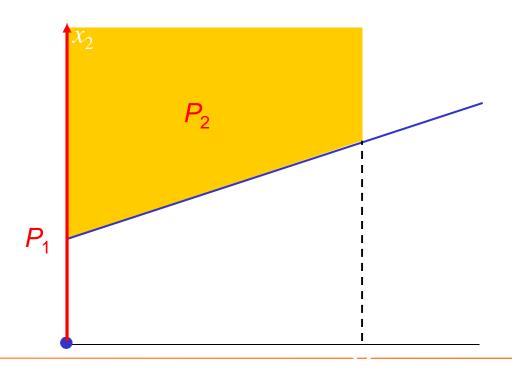
$$y_k \in \{0,1\}$$

Example

Start with a disjunction of linear systems to represent the union of polyhedra

$$\min x_2$$

$$\begin{pmatrix} x_1 = 0 \\ x_2 \ge 0 \end{pmatrix} \lor \begin{pmatrix} 0 \le x_1 \le M \\ x_2 \ge f + cx_1 \end{pmatrix}$$



Example

Start with a disjunction of linear systems to represent the union of polyhedra

Introduce a 0-1 variable y_k that is 1 when x is in polyhedron \underline{k} .

Disaggregate x to create an x^k for each k.

$$\min x_2$$

$$\begin{pmatrix} x_1 = 0 \\ x_2 \ge 0 \end{pmatrix} \lor \begin{pmatrix} 0 \le x_1 \le M \\ x_2 \ge f + cx_1 \end{pmatrix}$$

min
$$cx$$

 $x_1^1 = 0, x_2^1 \ge 0$
 $0 \le x_1^2 \le My_2, -cx_1^2 + x_2^2 \ge fy_2$
 $y_1 + y_2 = 1, y_k \in \{0,1\}$
 $x = x^1 + x^2$

Example

Replace x_1^2 with x_1 .

Replace x_2^2 with x_2 .

Replace y_2 with y.

min
$$x_2$$

$$x_1^1 = 0, \quad x_2^1 \ge 0$$

$$0 \le x_1^2 \le My_2, -cx_1^2 + x_2^2 \ge fy_2$$

$$y_1 + y_2 = 1, y_k \in \{0,1\}$$

$$\mathbf{X} = \mathbf{X}^1 + \mathbf{X}^2$$

This yields

min
$$x_2$$

$$0 \le x_1 \le My$$

$$x_2 \ge fy + cx_1$$

$$y \in \{0,1\}$$

or

min
$$fy + cx$$

$$0 \le x \le My$$

$$y \in \{0,1\} \quad \text{"Big } M\text{"}$$

$$y \in \{0,1\}$$

- Utilitarian and Rawlsian distributions seem too extreme in practice.
 - How to combine them?

- Utilitarian and Rawlsian distributions seem too extreme in practice.
 - How to combine them?

One proposal:

- Maximize welfare of worst off (Rawlsian)...
- ...until this requires undue sacrifice from others
- Seems appropriate in health care allocation.
- Joint work with H. P. Williams, to appear in *Management Science*.

- In particular:
 - Switch from **Rawlsian** to **utilitarian** when **inequality** exceeds Δ .

• In particular:

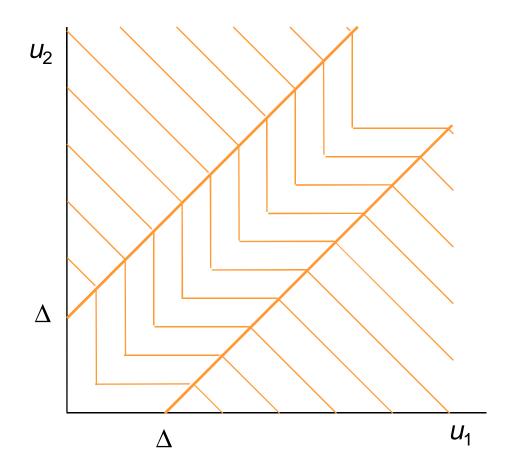
- Switch from **Rawlsian** to **utilitarian** when **inequality** exceeds Δ .
- Build mixed integer programming model.
- Let u_i = utility allocated to person i

• For 2 persons:

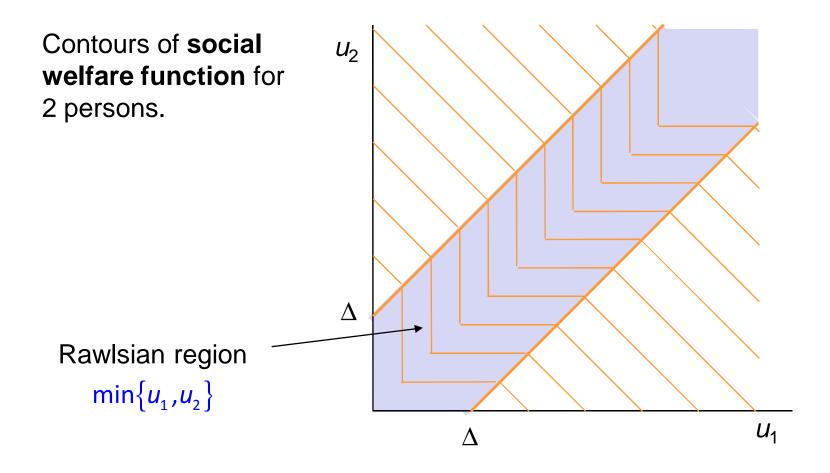
- Maximize $\min_{i} \{u_1, u_2\}$ (Rawlsian) when $|u_1 u_2| \le \Delta$
- Maximize $u_1 + u_2$ (utilitarian) when $|u_1 u_2| > \Delta$

Two-person Model

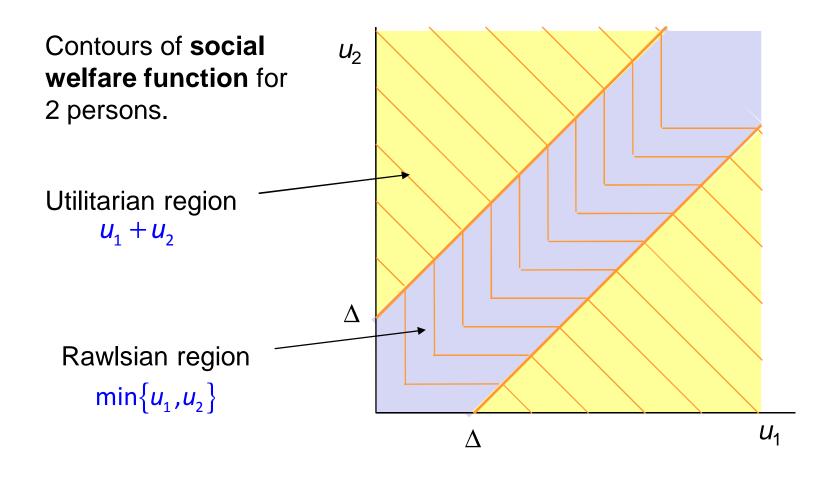
Contours of **social** welfare function for 2 persons.

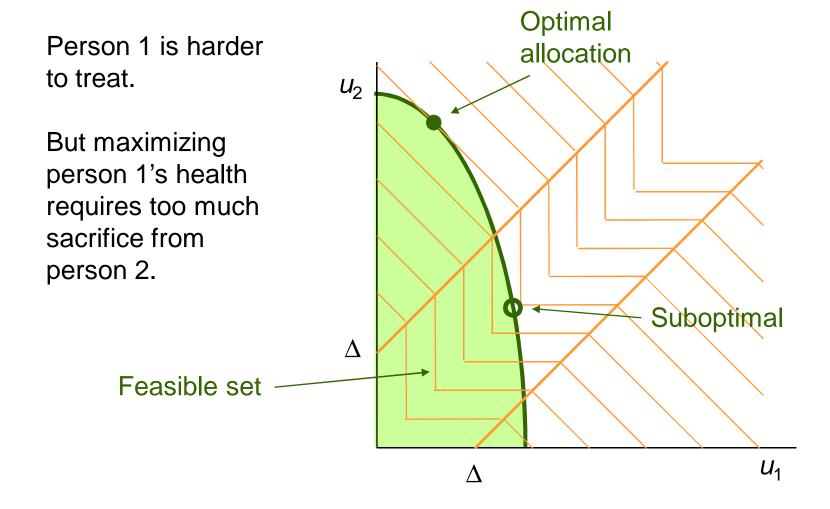


Two-person Model



Two-person Model



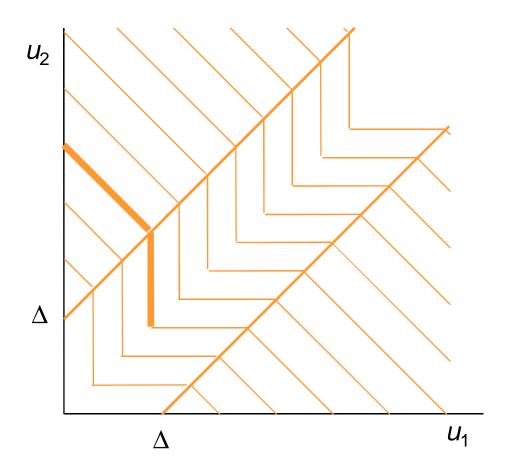


Advantages

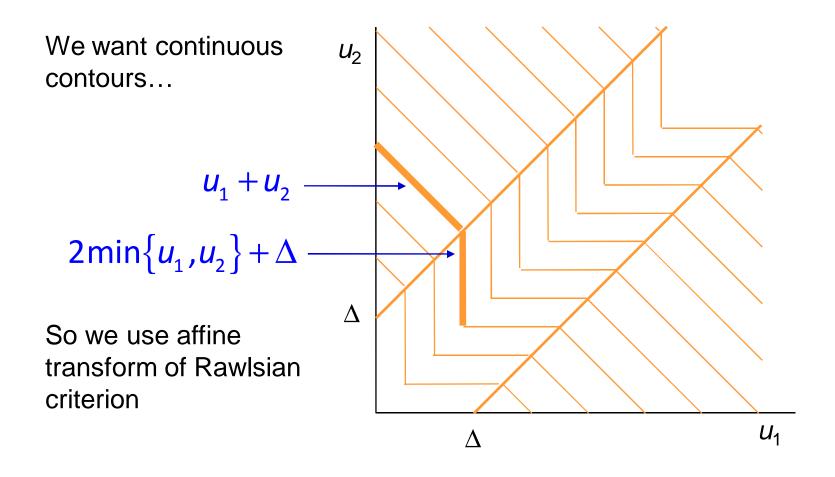
- Only one parameter ∆
 - Focus for debate.
 - $-\Delta$ has **intuitive meaning** (unlike weights)
 - Examine **consequences** of different settings for Δ
 - Find least objectionable setting
 - Results in a consistent policy

Social Welfare Function

We want continuous contours...



Social Welfare Function



Social Welfare Function

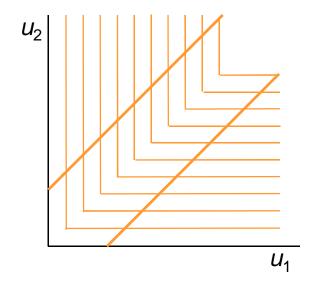
The social welfare problem becomes

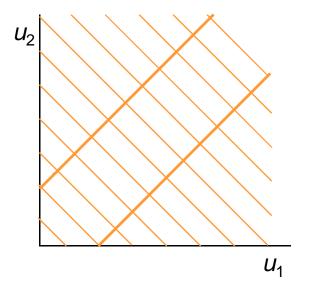
max z

$$z \leq \begin{cases} 2\min\{u_1, u_2\} + \Delta, & \text{if } |u_1 - u_2| \leq \Delta \\ u_1 + u_2, & \text{otherwise} \end{cases}$$

constraints on feasible set

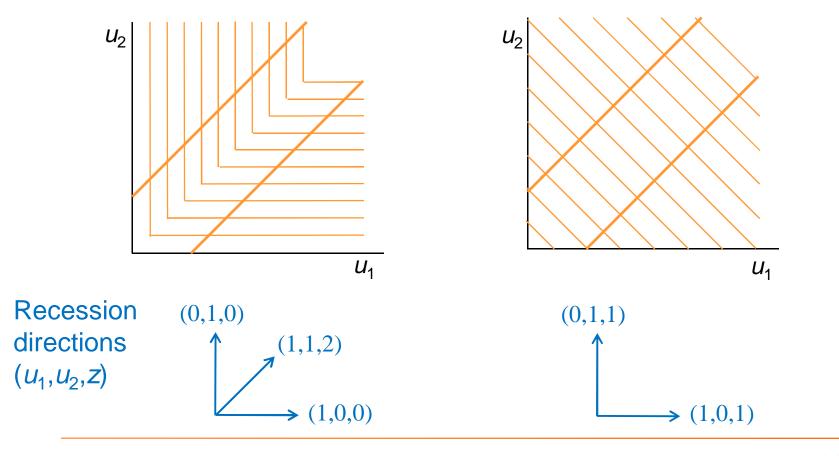
Epigraph is union of 2 polyhedra.



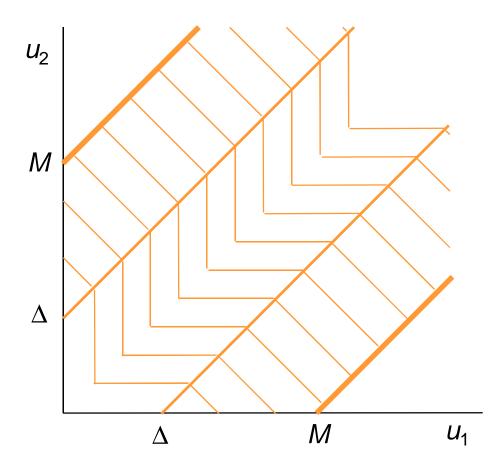


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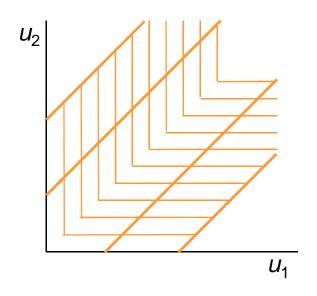
Because they have different recession cones, there is no MILP model.

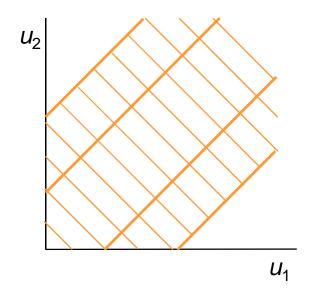


Impose constraints $|u_1 - u_2| \le M$



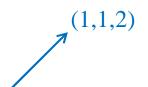
This equalizes recession cones.





Recession directions (u_1, u_2, z)





We have the model...

```
\max z
z \leq 2u_i + \Delta + (M - \Delta)\delta, \quad i = 1, 2
z \leq u_1 + u_2 + \Delta(1 - \delta)
u_1 - u_2 \leq M, \quad u_2 - u_1 \leq M
u_1, u_2 \geq 0
\delta \in \{0, 1\}
constraints on feasible set
```

 U_1

We have the model...

$$\max z$$

$$z \le 2u_i + \Delta + (M - \Delta)\delta, \quad i = 1, 2$$

$$z \le u_1 + u_2 + \Delta(1 - \delta)$$

$$u_1 - u_2 \le M, \quad u_2 - u_1 \le M$$

$$u_1, u_2 \ge 0$$

$$\delta \in \{0, 1\}$$

 U_1

This is a convex hull formulation.

n-person Model

Rewrite the 2-person social welfare function as...

$$\Delta + 2u_{\min} + \left(u_1 - u_{\min} - \Delta\right)^+ + \left(u_2 - u_{\min} - \Delta\right)^+$$

$$\min\{u_1, u_2\}$$

$$\alpha^+ = \max\{0, \alpha\}$$

n-person Model

Rewrite the 2-person social welfare function as...

$$\Delta + 2u_{\min} + \left(u_1 - u_{\min} - \Delta\right)^+ + \left(u_2 - u_{\min} - \Delta\right)^+ \times$$

$$\alpha^+ = \max\{0, \alpha\}$$

This can be generalized to *n* persons:

$$(n-1)\Delta + nu_{\min} + \sum_{j=1}^{n} \left(u_{j} - u_{\min} - \Delta\right)^{+}$$

n-person Model

Rewrite the 2-person social welfare function as...

$$\Delta + 2u_{\min} + (u_1 - u_{\min} - \Delta)^+ + (u_2 - u_{\min} - \Delta)^+$$

$$\alpha^+ = \max\{0, \alpha\}$$

This can be generalized to *n* persons:

$$(n-1)\Delta + nu_{\min} + \sum_{j=1}^{n} \left(u_{j} - u_{\min} - \Delta\right)^{+}$$

Epigraph is a union of n! polyhedra with same recession direction (u,z) = (1,...,1,n) if we require $|u_i - u_j| \le M$

So there is an MILP model...

n-person MILP Model

To avoid n! 0-1 variables, add auxiliary variables w_{ii}

```
\max z
z \leq u_i + \sum_{j \neq i} w_{ij}, \text{ all } i
w_{ij} \leq \Delta + u_i + \delta_{ij} (M - \Delta), \text{ all } i, j \text{ with } i \neq j
w_{ij} \leq u_j + (1 - \delta_{ij})\Delta, \text{ all } i, j \text{ with } i \neq j
u_i - u_j \leq M, \text{ all } i, j
u_i \geq 0, \text{ all } i
\delta_{ij} \in \{0,1\}, \text{ all } i, j \text{ with } i \neq j
```

n-person MILP Model

To avoid n! 0-1 variables, add auxiliary variables w_{ij}

$$\begin{aligned} &\max \ z \\ &z \leq u_i + \sum_{j \neq i} w_{ij}, \ \text{all} \ i \\ &w_{ij} \leq \Delta + u_i + \delta_{ij} (M - \Delta), \ \text{all} \ i,j \ \text{with} \ i \neq j \\ &w_{ij} \leq u_j + (1 - \delta_{ij})\Delta, \ \text{all} \ i,j \ \text{with} \ i \neq j \\ &u_i - u_j \leq M, \ \text{all} \ i,j \\ &u_i \geq 0, \ \text{all} \ i \\ &\delta_{ij} \in \{0,1\}, \ \text{all} \ i,j \ \text{with} \ i \neq j \end{aligned}$$

Theorem. The model is correct (not easy to prove).

n-person MILP Model

To avoid n! 0-1 variables, add auxiliary variables w_{ii}

$$\begin{aligned} &\max \ z \\ &z \leq u_i + \sum_{j \neq i} w_{ij}, \ \text{all} \ i \\ &w_{ij} \leq \Delta + u_i + \delta_{ij} (M - \Delta), \ \text{all} \ i,j \ \text{with} \ i \neq j \\ &w_{ij} \leq u_j + (1 - \delta_{ij})\Delta, \ \text{all} \ i,j \ \text{with} \ i \neq j \\ &u_i - u_j \leq M, \ \text{all} \ i,j \\ &u_i \geq 0, \ \text{all} \ i \\ &\delta_{ij} \in \{0,1\}, \ \text{all} \ i,j \ \text{with} \ i \neq j \end{aligned}$$

Theorem. The model is correct (not easy to prove).

Theorem. This is a convex hull formulation (not easy to prove).

n-group Model

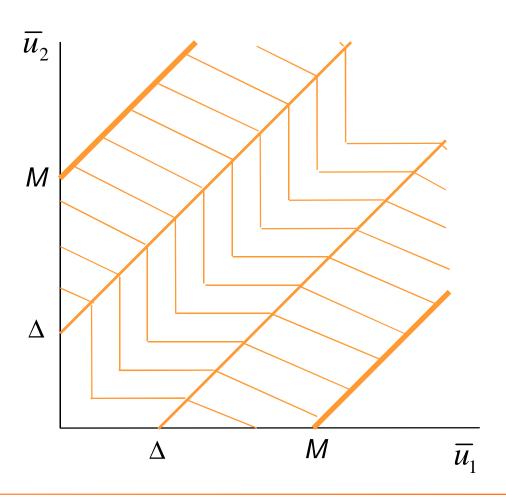
In practice, funds may be allocated to groups of different sizes

For example, disease/treatment categories.

```
Let \bar{u}_i = \text{average utility gained by a person in group } i
n_i = \text{size of group } i
```

n-group Model

2-person case with $n_1 < n_2$. Contours have slope $-n_1/n_2$



n-group MILP Model

Again add auxiliary variables w_{ij}

$$\begin{aligned} &\max \ z \\ &z \leq (n_i - 1)\Delta + n_i \overline{u}_i + \sum_{j \neq i} w_{ij}, \ \text{all } i \\ &w_{ij} \leq n_j (\overline{u}_i + \Delta) + \delta_{ij} n_j (M - \Delta), \ \text{all } i, j \text{ with } i \neq j \\ &w_{ij} \leq \overline{u}_j + (1 - \delta_{ij}) n_j \Delta, \ \text{all } i, j \text{ with } i \neq j \\ &\overline{u}_i - \overline{u}_j \leq M, \ \text{all } i, j \\ &\overline{u}_i \geq 0, \ \text{all } i \\ &\delta_{ii} \in \{0, 1\}, \ \text{all } i, j \text{ with } i \neq j \end{aligned}$$

Theorem. The model is correct.

Theorem. This is a convex hull formulation.

Health Example

Measure utility in QALYs (quality-adjusted life years).

QALY and cost data based on Briggs & Gray, (2000) etc.

Each group is a disease/treatment pair.

Treatments are discrete, so group funding is all-or-nothing.

Divide groups into relatively homogeneous subgroups.

 U_1

Health Example

Add constraints to define feasible set...

```
max z
z \leq (n_i - 1)\Delta + n_i \overline{u}_i + \sum_{i \neq i} w_{ij}, all i
w_{ij} \le n_i(\overline{u}_i + \Delta) + \delta_{ij}n_i(M - \Delta), all i, j with i \ne j
w_{ij} \leq \overline{u}_i + (1 - \delta_{ij})n_i \Delta, all i, j with i \neq j
\overline{u}_i - \overline{u}_i \leq M, all i, j
\overline{u}_i \ge 0, all i
                                                                                            U_1
\delta_{ij} \in \{0,1\}, all i,j with i \neq j
                                                                   y_i indicates
                                                                  whether
                                                                   group i is
                                                                   funded
```

	Intervention	Cost per person c_i (£)	QALYs gained q_i	Cost per QALY (£)	$\begin{array}{c} {\rm QALYs} \\ {\rm without} \\ {\rm intervention} \\ {\alpha_i} \end{array}$	Subgroup size n_i
	Pacemaker for atriove	entricular hea	rt block			
	Subgroup A	3500	3	1167	13	35
	Subgroup B	3500	5	700	10	45
	Subgroup C	3500	10	350	5	35
OALV	Hip replacement					
QALY	Subgroup A	3000	2	1500	3	45
& cost	Subgroup B	3000	4	750	4	45
a cost	Subgroup C	3000	8	375	5	45
data	Valve replacement for	aortic stenos	is			
	Subgroup A	4500	3	1500	2.5	20
	Subgroup B	4500	5	900	3	20
Dont 4	Subgroup C	4500	10	450	3.5	20
Part 1	CABG ¹ for left main	disease				
	Mild angina	3000	1.25	2400	4.75	50
	Moderate angina	3000	2.25	1333	3.75	55
	Severe angina	3000	2.75	1091	3.25	60
	CABG for triple vesse	el disease				
	Mild angina	3000	0.5	6000	5.5	50
	Moderate angina	3000	1.25	2400	4.75	55
	Severe angina	3000	2.25	1333	3.75	60
	CABG for double ves.	sel disease				
	Mild angina	3000	0.25	12,000	5.75	60
	Moderate angina	3000	0.75	4000	5.25	65
	Severe angina	3000	1.25	2400	4.75	70

	Intervention	Cost per person c_i (£)	QALYs gained q_i	Cost per QALY (£)	QALYs without intervention α_i	Subgroup size n_i
		22,500	4.5	5000	1.1	2
	Kidney transplant					
	Subgroup A	15,000	4	3750	1	8
QALY	Subgroup B	15,000	6	2500	1	8
	Kidney dialysis					
& cost	Less than 1 year st	urvival				
data	Subgroup A	5000	0.1	50,000	0.3	8
data	1-2 years survival					
	Subgroup B	12,000	0.4	30,000	0.6	6
Don't O	2-5 years survival					
Part 2	Subgroup C	20,000	1.2	16,667	0.5	4
	Subgroup D	28,000	1.7	16,471	0.7	4
	Subgroup E	36,000	2.3	15,652	0.8	4
	5-10 years survival					
	Subgroup F	46,000	3.3	13,939	0.6	3
	Subgroup G	56,000	3.9	14,359	0.8	2 2
	Subgroup H	66,000	4.7	14,043	0.9	2
	Subgroup I	77,000	5.4	14,259	1.1	2
	At least 10 years s	urvival				
	Subgroup J	88,000	6.5	13,538	0.9	2
	Subgroup K	100,000	7.4	13,514	1.0	1
	Subgroup L	111,000	8.2	13,537	1.2	1

Total budget £3 million

Δ	Pace-	Hip	Aortic	(CABO	3	Heart	Kidney		Ki	dney	dialy	sis
range	$_{\mathrm{maker}}$	repl.	valve	L	3	2	trans.	trans.	< 1	1-2	2-5	5-10	> 10
0-3.3	111	111	111	111	111	111	1	11	0	0	000	0000	000
3.4 – 4.0	111	111	111	111	111	111	0	11	1	0	000	0000	000
4.0 – 4.4	111	111	111	111	111	111	0	01	1	0	000	0000	001
4.5 - 5.01	111	011	111	111	111	111	1	01	1	0	000	0000	011
5.02 - 5.55	111	011	011	111	111	111	0	01	1	0	000	0001	011
5.56 - 5.58	111	011	011	111	111	011	0	01	1	0	000	0001	111
5.59	111	011	011	110	111	111	0	01	1	0	000	0001	111
5.60 - 13.1	111	111	111	101	000	000	1	11	1	0	111	1111	111
13.2 – 14.2	111	011	111	011	000	000	1	11	1	1	111	1111	111
14.3 - 15.4	111	111	111	011	000	000	1	11	1	1	101	1111	111
15.5–up	111	011	111	011	001	000	1	11	1	0	011	1111	111

Utilitarian solution

Δ	Pace-	Hip	Aortic	(CABO	3	Heart	Kidney		Ki	dney	dialy	sis
range	$_{\mathrm{maker}}$	repl.	valve	L	3	2	trans.	trans.	< 1	1-2	2-5	5-10	> 10
0-3.3	111	111	111	111	111	111	1	11	0	0	000	0000	000
3.4 – 4.0	111	111	111	111	111	111	0	11	1	0	000	0000	000
4.0 – 4.4	111	111	111	111	111	111	0	01	1	0	000	0000	001
4.5 - 5.01	111	011	111	111	111	111	1	01	1	0	000	0000	011
5.02 – 5.55	111	011	011	111	111	111	0	01	1	0	000	0001	011
5.56 - 5.58	111	011	011	111	111	011	0	01	1	0	000	0001	111
5.59	111	011	011	110	111	111	0	01	1	0	000	0001	111
5.60 - 13.1	111	111	111	101	000	000	1	11	1	0	111	1111	111
13.2 - 14.2	111	011	111	011	000	000	1	11	1	1	111	1111	111
14.3 - 15.4	111	111	111	011	000	000	1	11	1	1	101	1111	111
15.5-up	111	011	111	011	001	000	1	11	1	0	011	1111	111

Rawlsian solution

Δ	Pace-	Hip	Aortic	(CABO	3	Heart	Kidney		Ki	dney	dialy	sis
range	$_{\mathrm{maker}}$	repl.	valve	L	3	2	trans.	trans.	< 1	1-2	2-5	5-10	> 10
0-3.3	111	111	111	111	111	111	1	11	0	0	000	0000	000
3.4 – 4.0	111	111	111	111	111	111	0	11	1	0	000	0000	000
4.0 – 4.4	111	111	111	111	111	111	0	01	1	0	000	0000	001
4.5 – 5.01	111	011	111	111	111	111	1	01	1	0	000	0000	011
5.02 - 5.55	111	011	011	111	111	111	0	01	1	0	000	0001	011
5.56 - 5.58	111	011	011	111	111	011	0	01	1	0	000	0001	111
5.59	111	011	011	110	111	111	0	01	1	0	000	0001	111
5.60 - 13.1	111	111	111	101	000	000	1	11	1	0	111	1111	111
13.2 - 14.2	111	011	111	011	000	000	1	11	1	1	111	1111	111
$14.3 – 15.4$ \checkmark	111	111	111	011	000	000	1	11	1	1	101	1111	111
15.5–up	111	011	111	011	001	000	1	11	1	0	011	1111	111

Fund for all Δ

	\checkmark		7										
Δ	Pace-	Hip	Aortic	(CABO	3	Heart	Kidney		K	idney	dialy	sis
range	$_{\mathrm{maker}}$	repl.	valve	\mathbf{L}	3	2	trans.	trans.	< 1	1-2	2-5	5-10	> 10
0-3.3	111	111	111	111	111	111	1	11	0	0	000	0000	000
3.4 – 4.0	111	111	111	111	111	111	0	11	1	0	000	0000	000
4.0 – 4.4	111	111	111	111	111	111	0	01	1	0	000	0000	001
4.5 – 5.01	111	011	111	111	111	111	1	01	1	0	000	0000	011
5.02 - 5.55	111	011	011	111	111	111	0	01	1	0	000	0001	011
5.56 - 5.58	111	011	011	111	111	011	0	01	1	0	000	0001	111
5.59	111	011	011	110	111	111	0	01	1	0	000	0001	111
5.60 - 13.1	111	111	111	101	000	000	1	11	1	0	111	1111	111
13.2 - 14.2	111	011	111	011	000	000	1	11	1	1	111	1111	111
14.3 - 15.4	111	111	111	011	000	000	1	11	1	1	101	1111	111
15.5–up	111	011	111	011	001	000	1	11	1	0	011	1111	111

More dialysis with larger Δ , beginning with longer life span

Δ	Pace-	Hip	Aortic	(CABO	3	Heart	Kidney		Kidney dialysis
range	$_{\mathrm{maker}}$	repl.	valve	L	3	2	trans.	trans.	< 1	1 1-2 2-5 5-10 > 10
0-3.3	111	111	111	111	111	111	1	11	0	0 000 0000 000
3.4 – 4.0	111	111	111	111	111	111	0	11	1	0 000 0000 000
4.0 – 4.4	111	111	111	111	111	111	0	01	1	0 000 0000 001
4.5 - 5.01	111	011	111	111	111	111	1	01	1	0 000 0000 011
5.02 - 5.55	111	011	011	111	111	111	0	01	1	0 000 0001 011
5.56 - 5.58	111	011	011	111	111	011	0	01	1	0 000 0001 111
5.59	111	011	011	110	111	111	0	01	1	0 000 0001 111
5.60 - 13.1	111	111	111	101	000	000	1	11	1	0 111 1111 111
13.2 - 14.2	111	011	111	011	000	000	1	11	1	1 111 1111 111
14.3 - 15.4	111	111	111	011	000	000	1	11	1	1 101 1111 111
15.5–up	111	011	111	011	001	000	1	11	1	0 011 1111 111

Abrupt change at $\Delta = 5.60$

Δ	Pace-	Hip	Aortic		CABC			Kidney				dialy	
range	maker	repl.	valve	L	3	2	trans.	trans.	< 1	1-2	2-5	5-10	> 10
0-3.3	111	111	111	111	111	111	1	11	0	0	000	0000	000
3.4 – 4.0	111	111	111	111	111	111	0	11	1	0	000	0000	000
4.0 – 4.4	111	111	111	111	111	111	0	01	1	0	000	0000	001
4.5 - 5.01	111	011	111	111	111	111	1	01	1	0	000	0000	011
5.02 - 5.55	111	011	011	111	111	111	0	01	1	0	000	0001	011
$5.56 – 5.58$ $^{\vee}$	111	011	011	111	111	011	0	01	1	0	000	0001	111
5.59	111	011	011	110	111	111	0	01	1	0	000	0001	111
5.60-13.1	111	111	111	101	000	000	1	11	1	0	111	1111	111
13.2-14.2	111	011	111	011	000	000	1	11	1	1	111	1111	111
14.3 - 15.4	111	111	111	011	000	000	1	11	1	1	101	1111	111
15.5-up	111	011	111	011	001	000	1	11	1	0	011	1111	111

Come and go together

Δ	Pace-	Hip	Aortic	(CABO	3	Heart	Kidney		K	idney	dialy	sis
range	$_{\mathrm{maker}}$	repl.	valve	L	3	2/	trans.	trans.	< 1	1-2	2-5	5-10	> 10
0-3.3	111	111	111	111	111	111	1	11	0	0	000	0000	000
3.4 – 4.0	111	111	111	111	111	111	0	11	1	0	000	0000	000
4.0 – 4.4	111	111	111	111	111	111	0	01	1	0	000	0000	001
4.5 – 5.01	111	011	111	111	111	111	1	01	1	0	000	0000	011
5.02 - 5.55	111	011	011	111	111	111	0	01	1	0	000	0001	011
5.56 - 5.58	111	011	011	111	111	011	0	01	1	0	000	0001	111
5.59	111	011	011	110	111	111	0	01	1	0	000	0001	111
5.60 - 13.1	111	111	111	101	000	000	1	11	1	0	111	1111	111
13.2 - 14.2	111	011	111	011	000	000	1	11	1	1	111	1111	111
14.3 - 15.4	111	111	111	011	000	000	1	11	1	1	101	1111	111
15.5–up	111	011	111	011	001	000	1	11	1	0	011	1111	111

	In-out-in													
				/		`								
Δ	Pace-	Hip	Aortic	/ (CABO	r J	Heart	Kidney		K	idney	dialy	sis	
range	$_{\mathrm{maker}}$	repl.	valve	L	3	2	trans.	trans.	< 1	1-2	2-5	5-10	> 10	
0-3.3	111	111	111	111	111	111	1	11	0	0	000	0000	000	
3.4 – 4.0	111	111	111	111	111	111	0	11	1	0	000	0000	000	
4.0 – 4.4	111	111	111	111	111	111	0	01	1	0	000	0000	001	
4.5 – 5.01	111	011	111	111	111	111	1	01	1	0	000	0000	011	
5.02 - 5.55	111	011	011	111	111	111	0	01	1	0	000	0001	011	
5.56 - 5.58	111	011	011	111	111	011	0	01	1	0	000	0001	111	
5.59	111	011	011	110	111	111	0	01	1	0	000	0001	111	
5.60 - 13.1	111	111	111	101	000	000	1	11	1	0	111	1111	111	
13.2 - 14.2	111	011	111	011	000	000	1	11	1	1	111	1111	111	
14.3 - 15.4	111	111	111	011	000	000	1	11	1	1	101	1111	111	
15.5-up	111	011	111	011	001	000	1	11	1	0	011	1111	111	

Most rapid change. Possible range for politically acceptable compromise

Δm_{range}	Pace- maker	Hip repl.	Aortic valve	L (CABO	G 2		Kidney trans.				dialy 5-10	
0-3.3	111	111	111	111	111	111	1	11	0	0	000	0000	000
3.4 – 4.0	111	111	111	111	111	111	0	11	1	0	000	0000	000
4.0 – 4.4	111	111	111	111	111	111	0	01	1	0	000	0000	001
4.5-5.01	111	011	111	111	111	111	1	01	1	0	000	0000	011
5.02 - 5.55	111	011	011	111	111	111	0	01	1	0	000	0001	011
5.56 - 5.58	111	011	011	111	111	011	0	01	1	0	000	0001	111
5.59	111	011	011	110	111	111	0	01	1	0	000	0001	111
5.60-13.1	111	111	111	101	000	000	1	11	1	0	111	1111	111
13.2-14.2	111	011	111	011	000	000	1	11	1	1	111	1111	111
14.3 - 15.4	111	111	111	011	000	000	1	11	1	1	101	1111	111
15.5–up	111	011	111	011	001	000	1	11	1	0	011	1111	111

32 groups, 1089 integer variables Solution time (CPLEX 12.2) is negligible

Δ	Pace-	Hip	Aortic	(CABC	3	Heart	Kidney		Ki	dney	dialy	sis
range	$_{\mathrm{maker}}$	repl.	valve	L	3	2	trans.	trans.	< 1	1-2	2-5	5-10	> 10
0-3.3	111	111	111	111	111	111	1	11	0	0	000	0000	000
3.4 – 4.0	111	111	111	111	111	111	0	11	1	0	000	0000	000
4.0 – 4.4	111	111	111	111	111	111	0	01	1	0	000	0000	001
4.5 - 5.01	111	011	111	111	111	111	1	01	1	0	000	0000	011
5.02 - 5.55	111	011	011	111	111	111	0	01	1	0	000	0001	011
5.56 - 5.58	111	011	011	111	111	011	0	01	1	0	000	0001	111
5.59	111	011	011	110	111	111	0	01	1	0	000	0001	111
5.60 - 13.1	111	111	111	101	000	000	1	11	1	0	111	1111	111
13.2 - 14.2	111	011	111	011	000	000	1	11	1	1	111	1111	111
14.3 - 15.4	111	111	111	011	000	000	1	11	1	1	101	1111	111
15.5–up	111	011	111	011	001	000	1	11	1	0	011	1111	111

Table 3: Solution times in seconds for m groups and different values of Δ . Instances with more than a few hundred groups seem very unlikely to occur in practice.

	Δ							
m	0	1	2	3	4	5	6	∞
330	0.02	1.2	0.67	0.56	0.50	0.30	0.03	0.02
660	0.03	4.1	1.6	1.6	0.92	0.80	0.05	0.02
990	0.02	5.2	3.1	3.6	1.5	1.5	0.08	0.02
1320	0.00	15	4.3	4.2	2.7	3.0	0.09	0.02
1980	0.02	24	11	11	11	5.4	0.14	0.02
2640	0.00	32	19	14	8.6	8.8	0.19	0.02
3300	0.17	51	43	44	34	13	0.25	0.02