

# Discrete Optimization with Decision Diagrams

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*Joint work with*

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# Goal

- A **general discrete optimization method** that combines:
  - Relaxation and bounding of **mixed integer programming (MIP)**
  - Primal heuristics of **MIP**
  - Propagation of **constraint programming (CP)**
  - Recursive modeling style of **dynamic programming (DP)**
- How?
  - Use binary/multivalued **decision diagrams** as a framework.

# Approach

- Decision diagram as a basic tool:
  - **Relaxed** decision diagrams provide **bounding** mechanism (MIP)
  - **Restricted** decision diagrams provide **primal heuristic** (MIP)
  - **Relaxed** decision diagrams enhanced **propagation** (CP)
  - **State variables** in decision diagram provide **DP modeling style**
- Plus...
  - **Decision diagram** framework provides **novel branching scheme**.

# Decision Diagrams

- **Binary decision diagrams (BDDs)** historically used for circuit design and verification.
  - Lee 1959, Akers 1978, Bryant 1986.

# Decision Diagrams

- **Binary decision diagrams (BDDs)** historically used for circuit design and verification.
  - Lee 1959, Akers 1978, Bryant 1986.
- **Compact** graphical representation of **boolean** function.
  - Can also represent **feasible set** of problem with binary variables.
  - Slight generalization (MDDs) represents finite domain variables.

# Decision Diagrams

- BDD is result of superimposing isomorphic subtrees in a search tree.
  - Unique reduced BDD for given variable ordering.
  - “Caching” is a popular theme in knowledge representation.
- Constraints need not have an inequality representation.

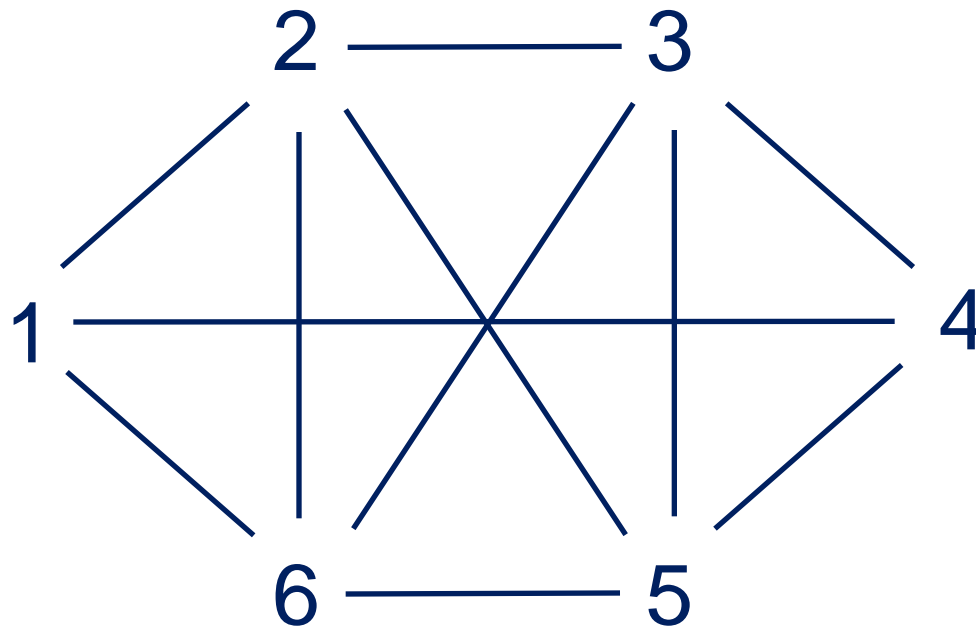
# Binary Decision Diagrams

- BDD can grow exponentially with problem size.
  - So we use a smaller, **relaxed** BDD that represents **superset** of feasible set.
    - Andersen, Hadzic, Hooker, Tiedemann 2007.
    - For alldiff systems, reduced search tree from >1 million nodes to 1 node.
    - Subsequent papers with Hadzic, Hoda, van Hoeve, O'Sullivan.
- We focus on **independent set problem** on a graph...

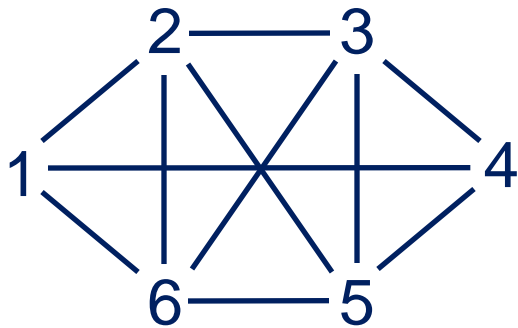
# Independent Set Problem

Let each vertex have weight  $w_i$

Select nonadjacent vertices to maximize  $\sum_i w_i x_i$

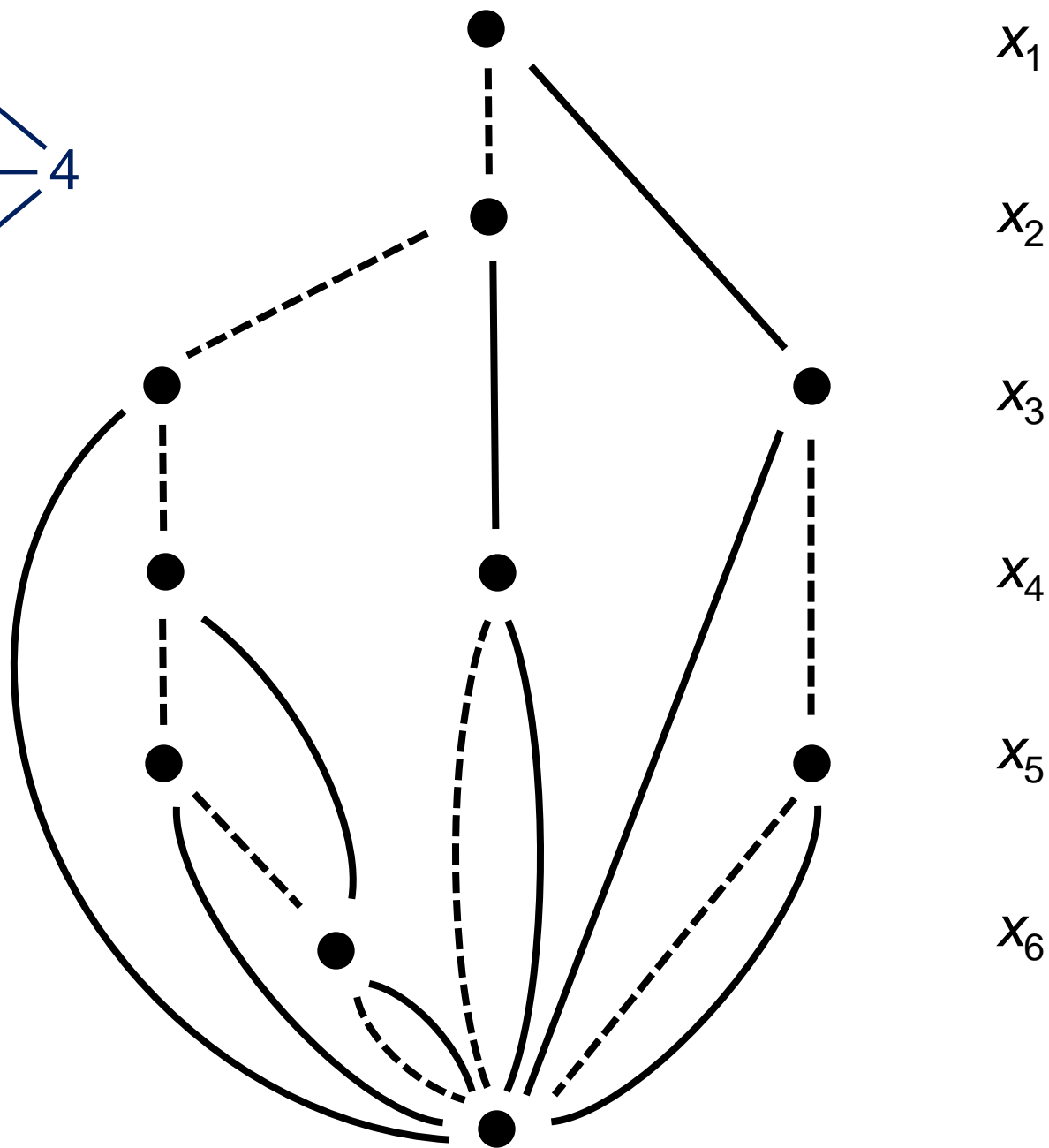


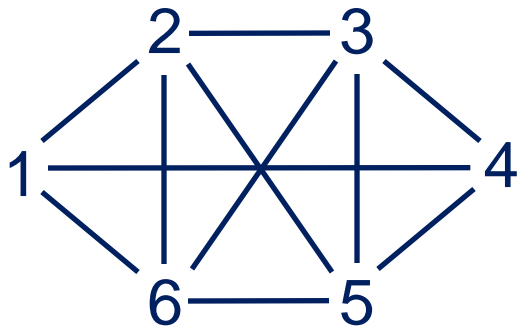




Exact BDD for  
independent set  
problem

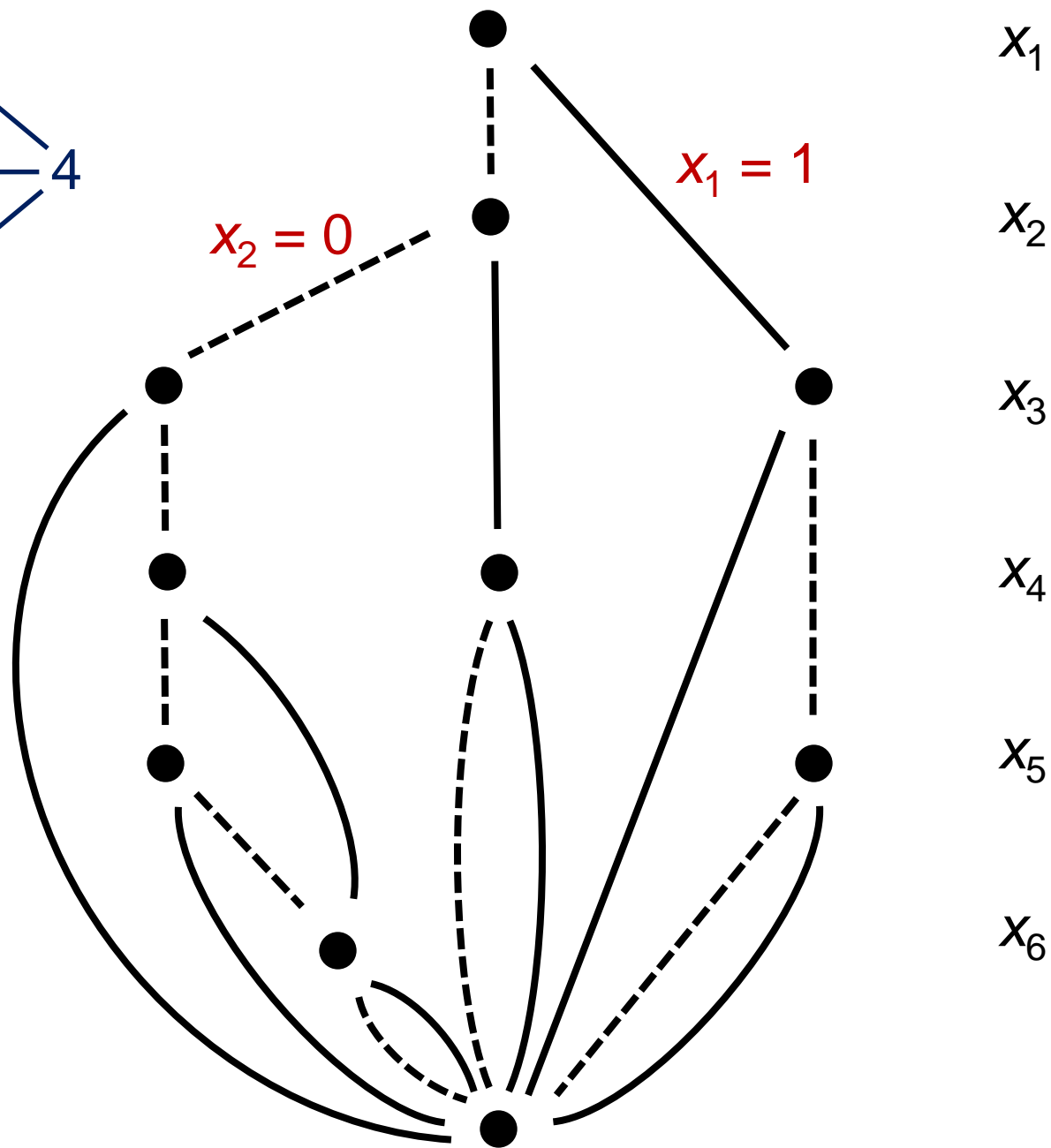
“zero-suppressed”  
BDD

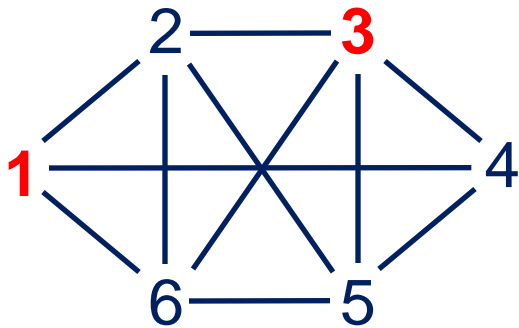




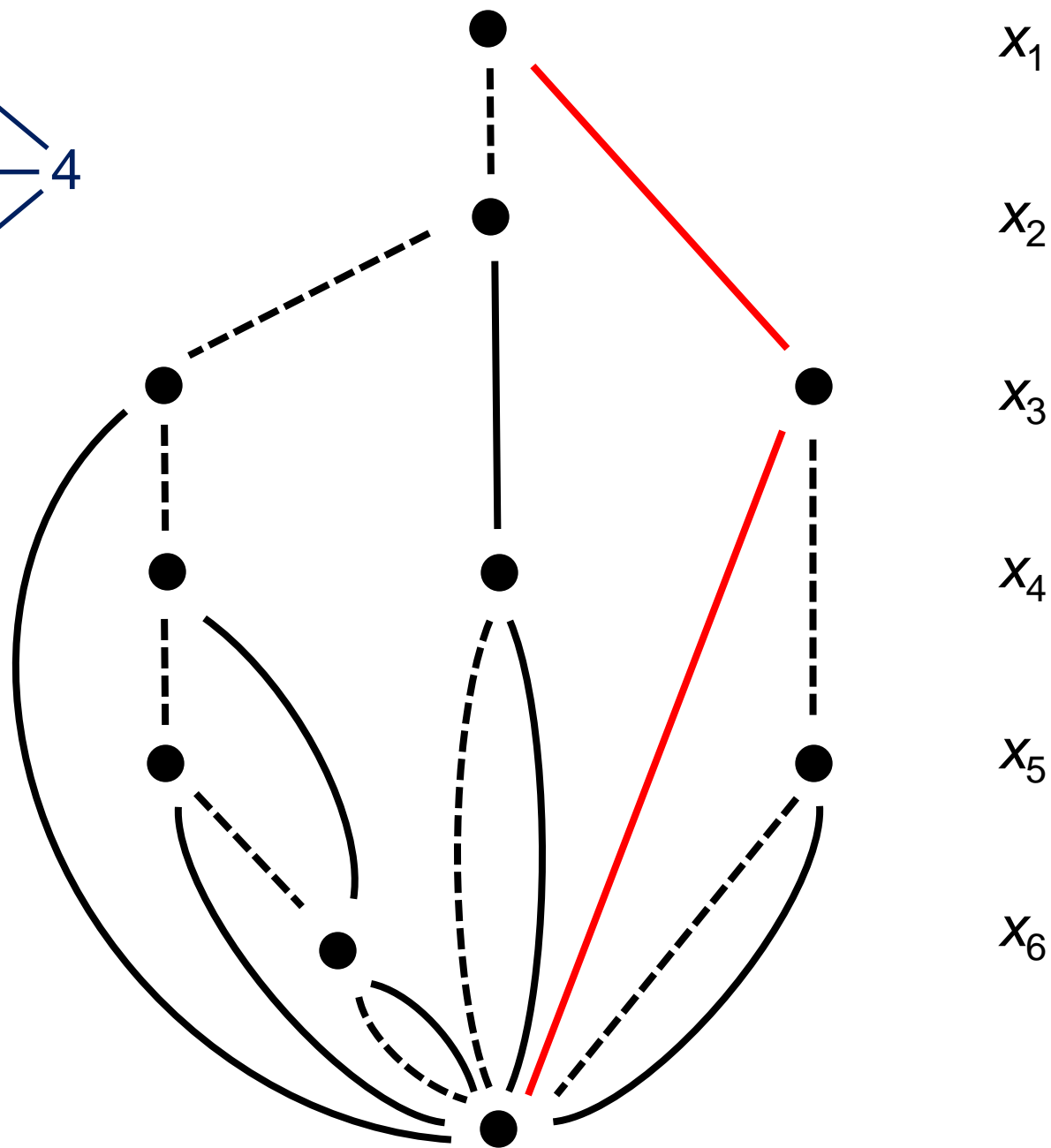
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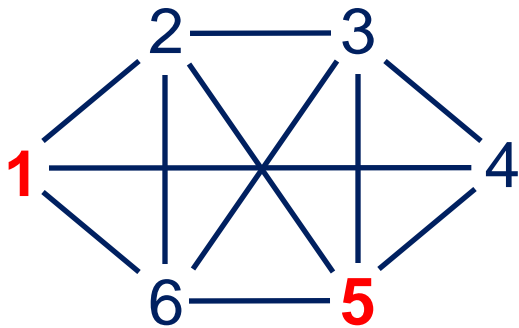
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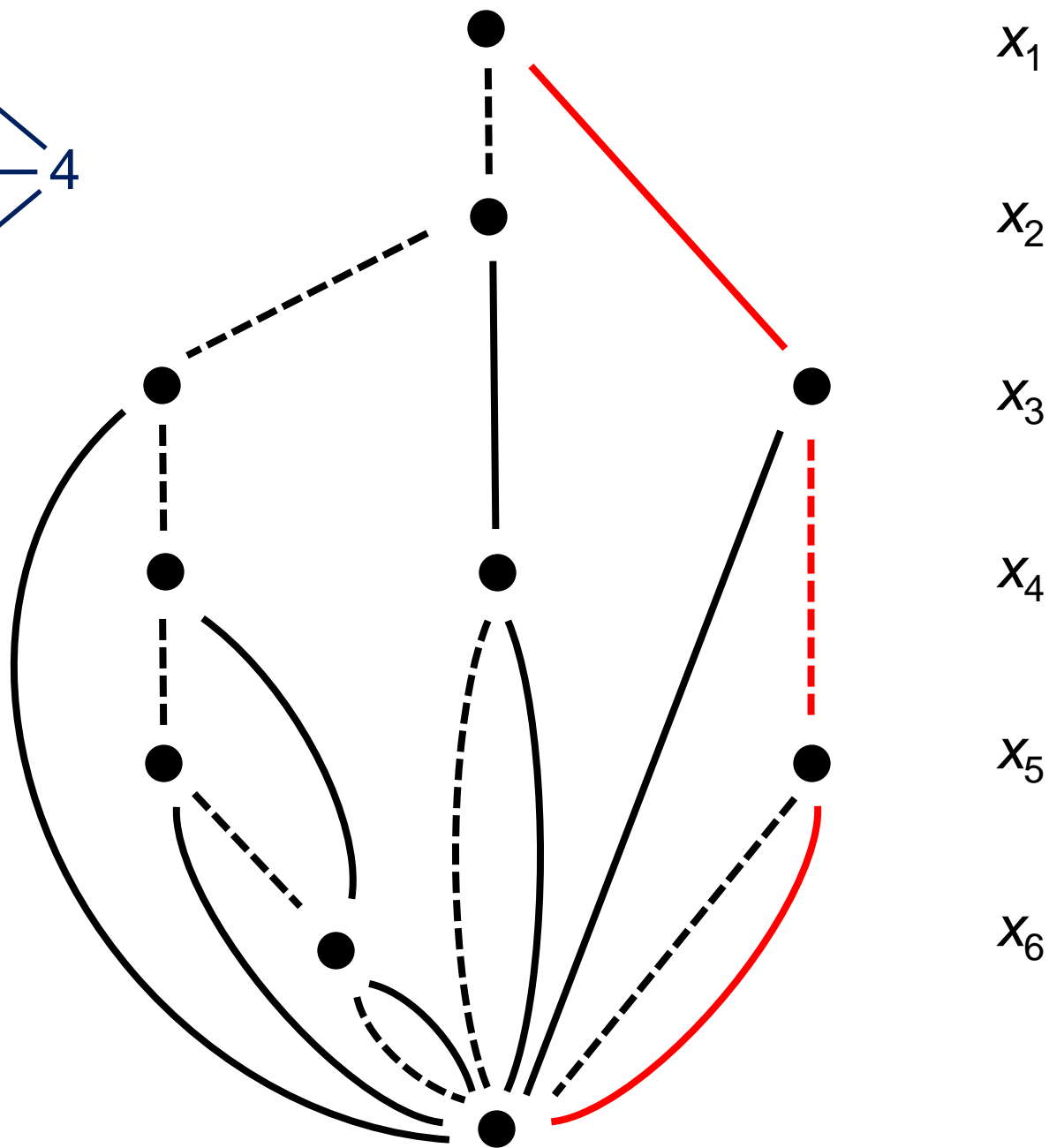


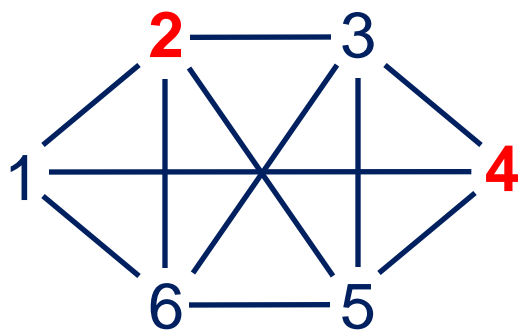
Paths from top to bottom correspond to the 11 feasible solutions



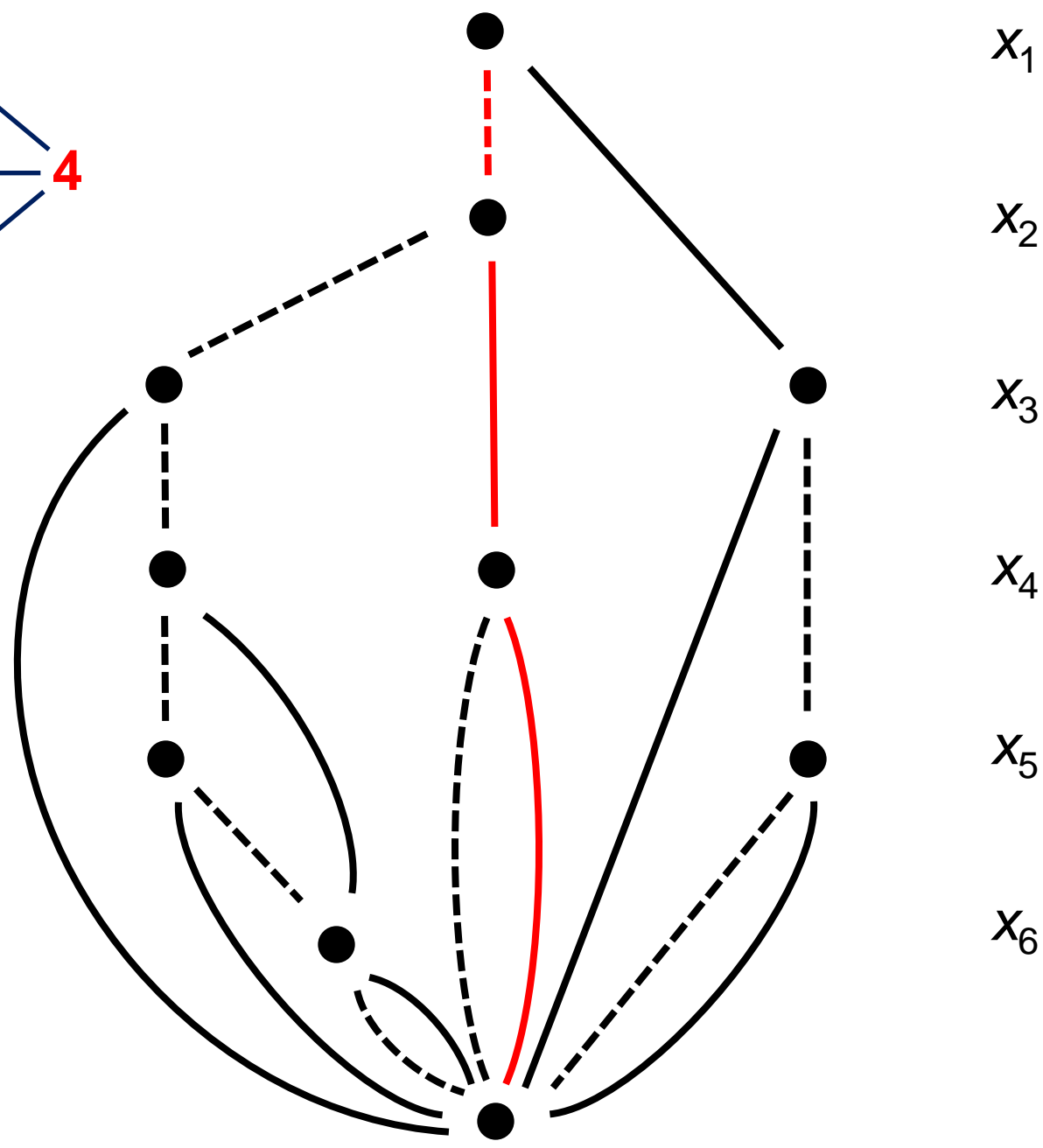


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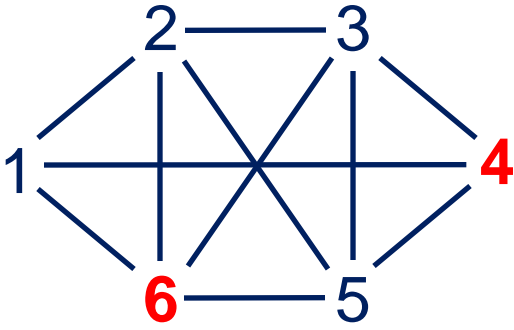




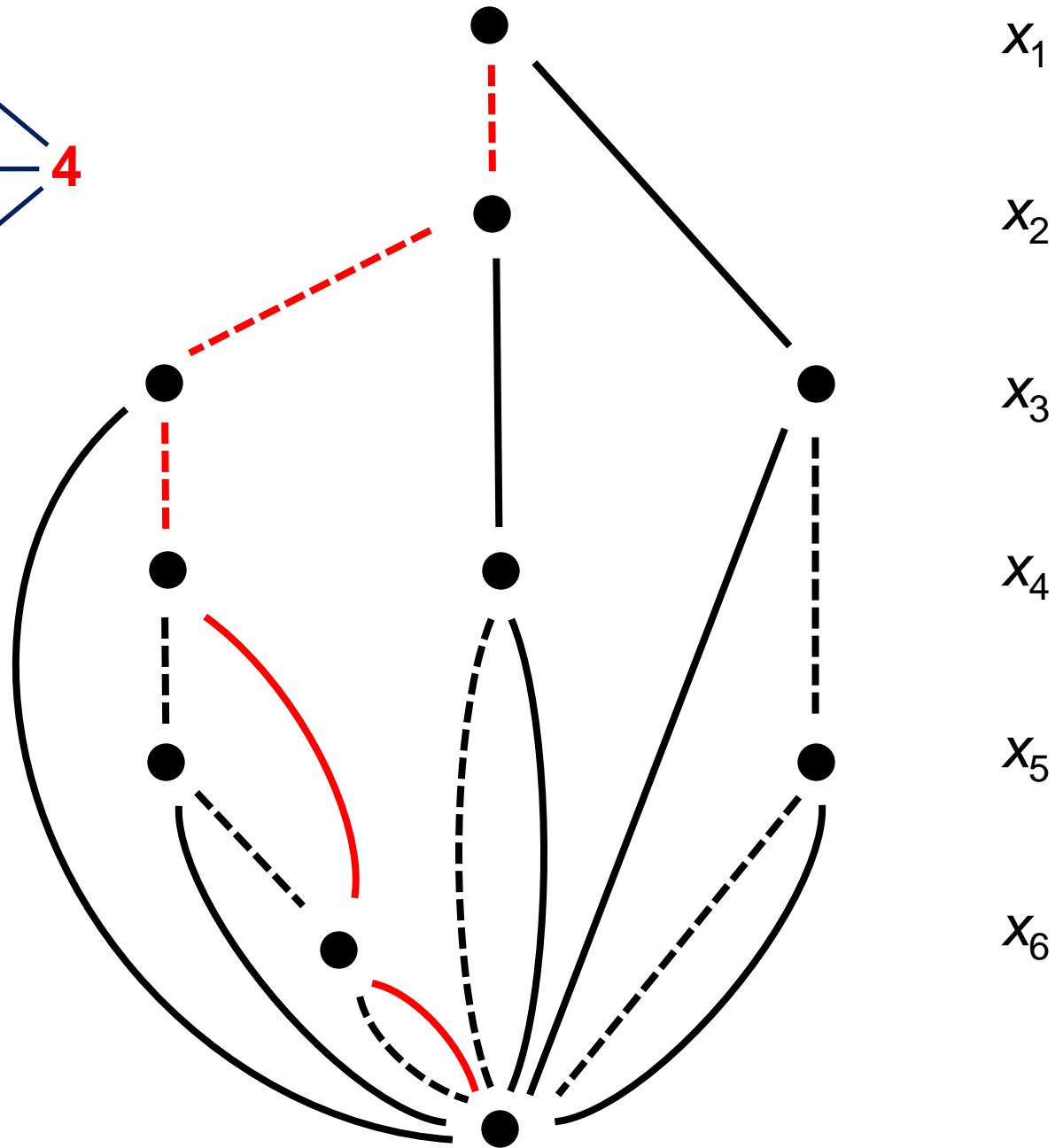
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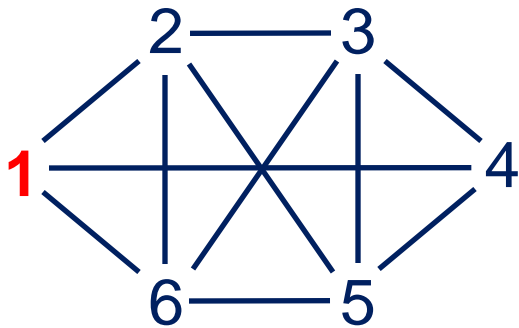


$x_1$   
 $x_2$   
 $x_3$   
 $x_4$   
 $x_5$   
 $x_6$



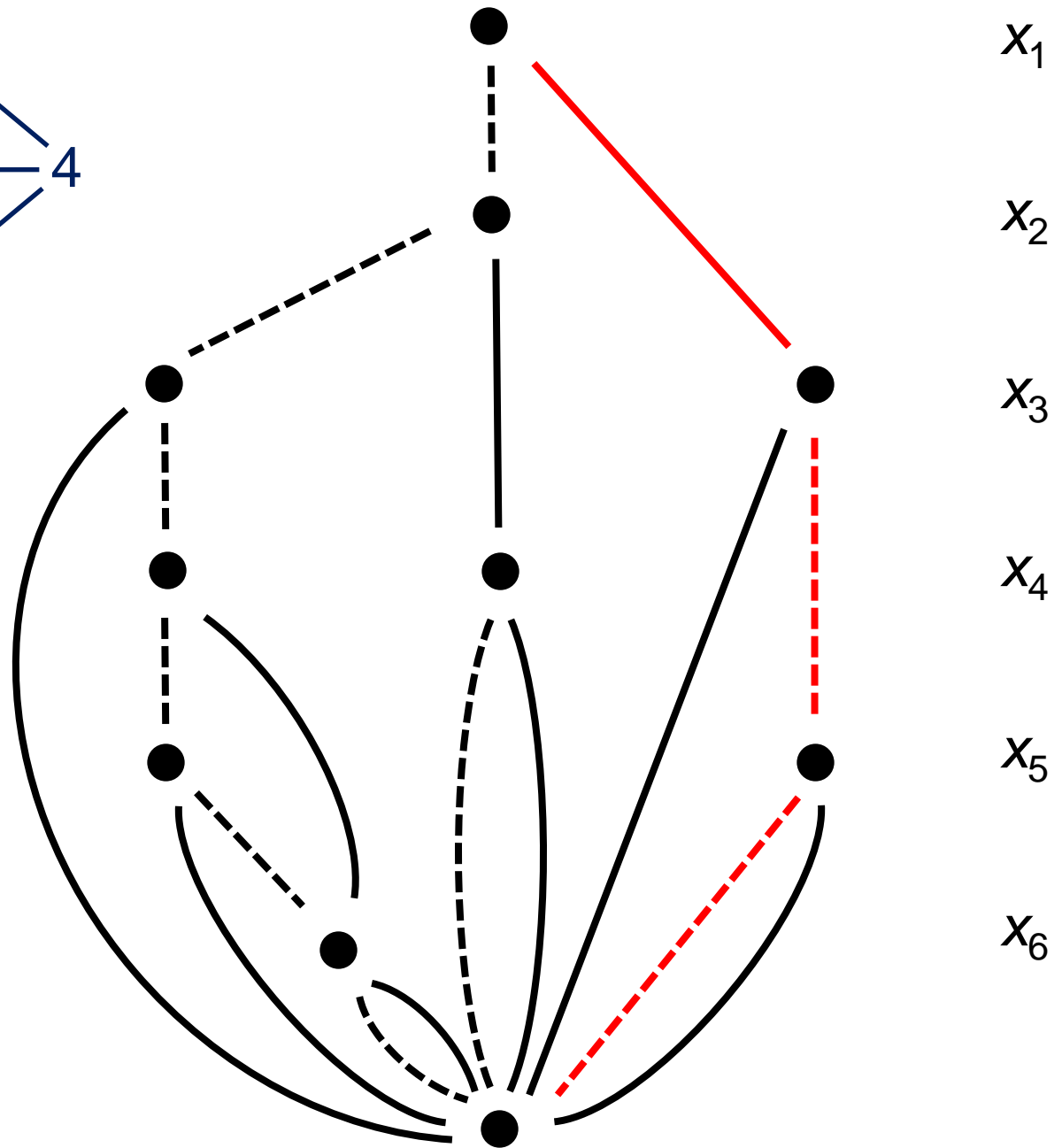
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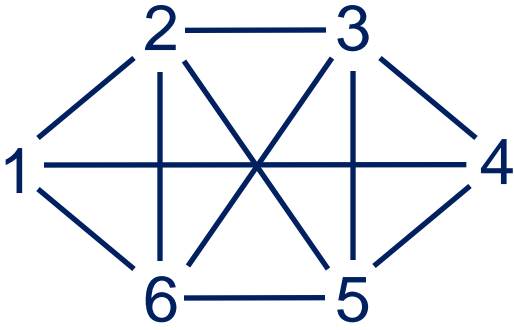




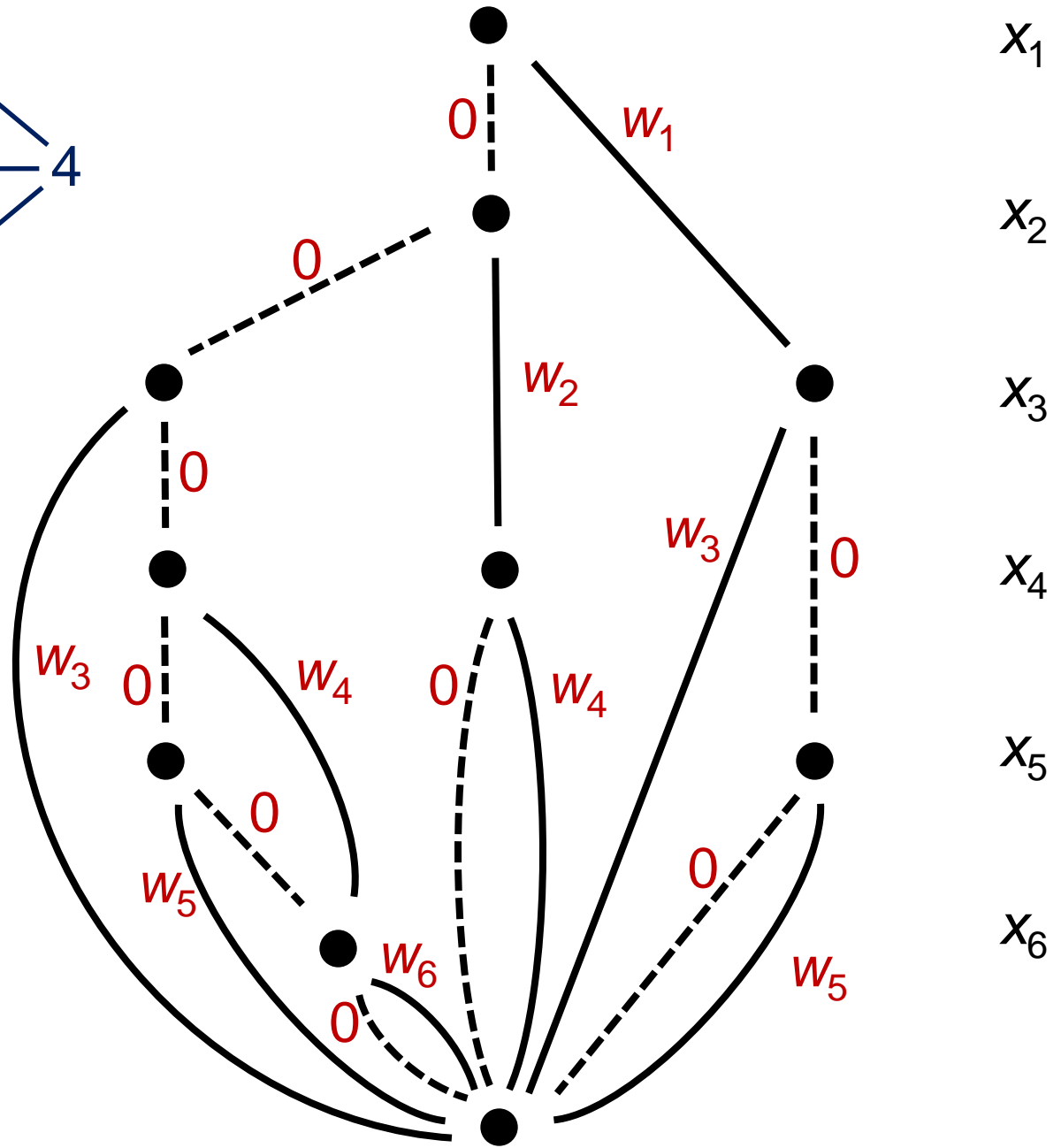
Paths from top to bottom correspond to the 11 feasible solutions

...and so forth



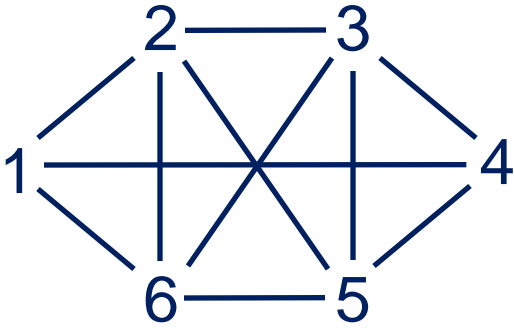


For objective function, associate weights with arcs



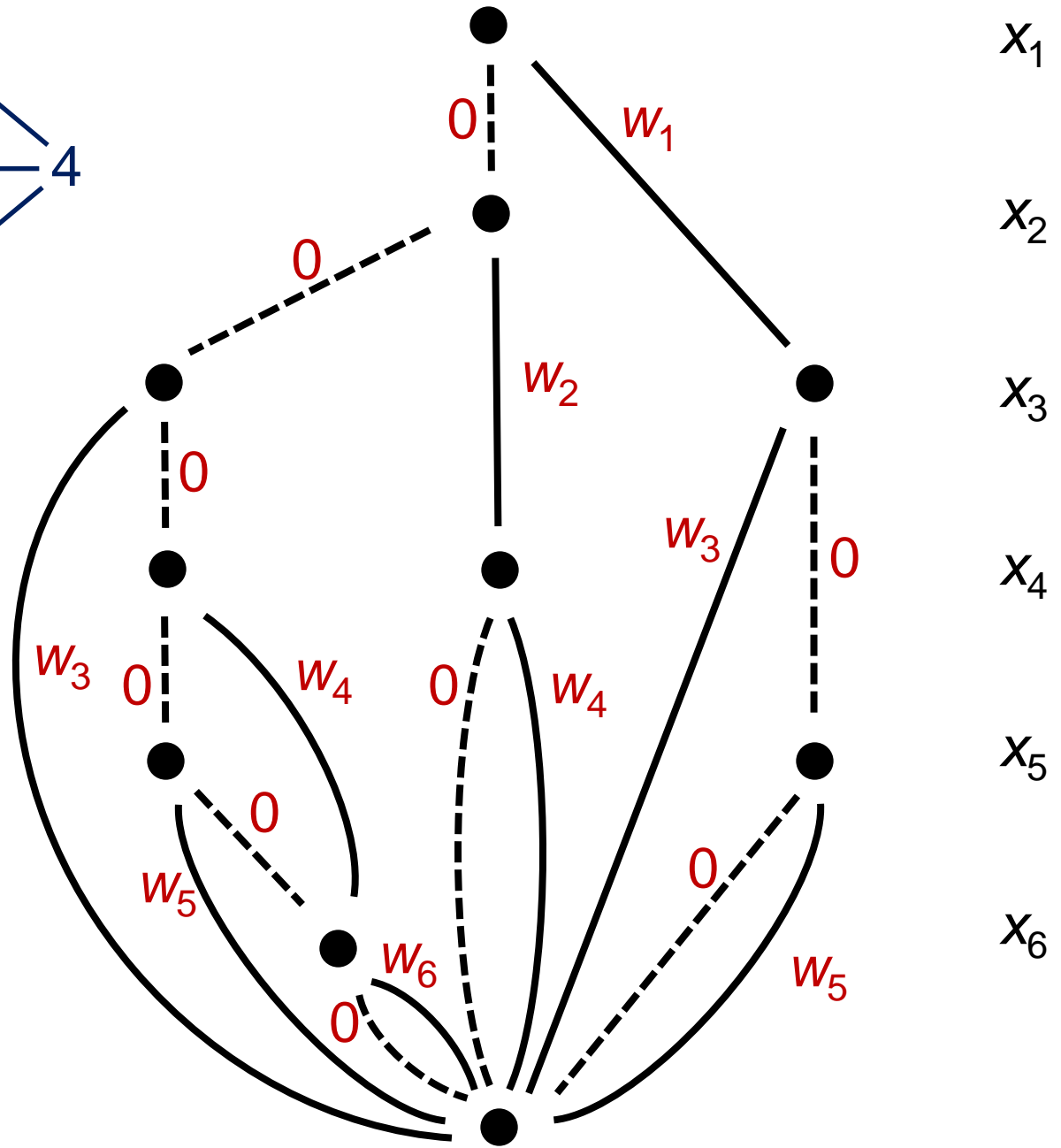
$x_1$   
 $x_2$   
 $x_3$   
 $x_4$   
 $x_5$   
 $x_6$





For objective function, associate weights with arcs

Optimal solution is **longest path**



$x_1$

$x_2$

$x_3$

$x_4$

$x_5$

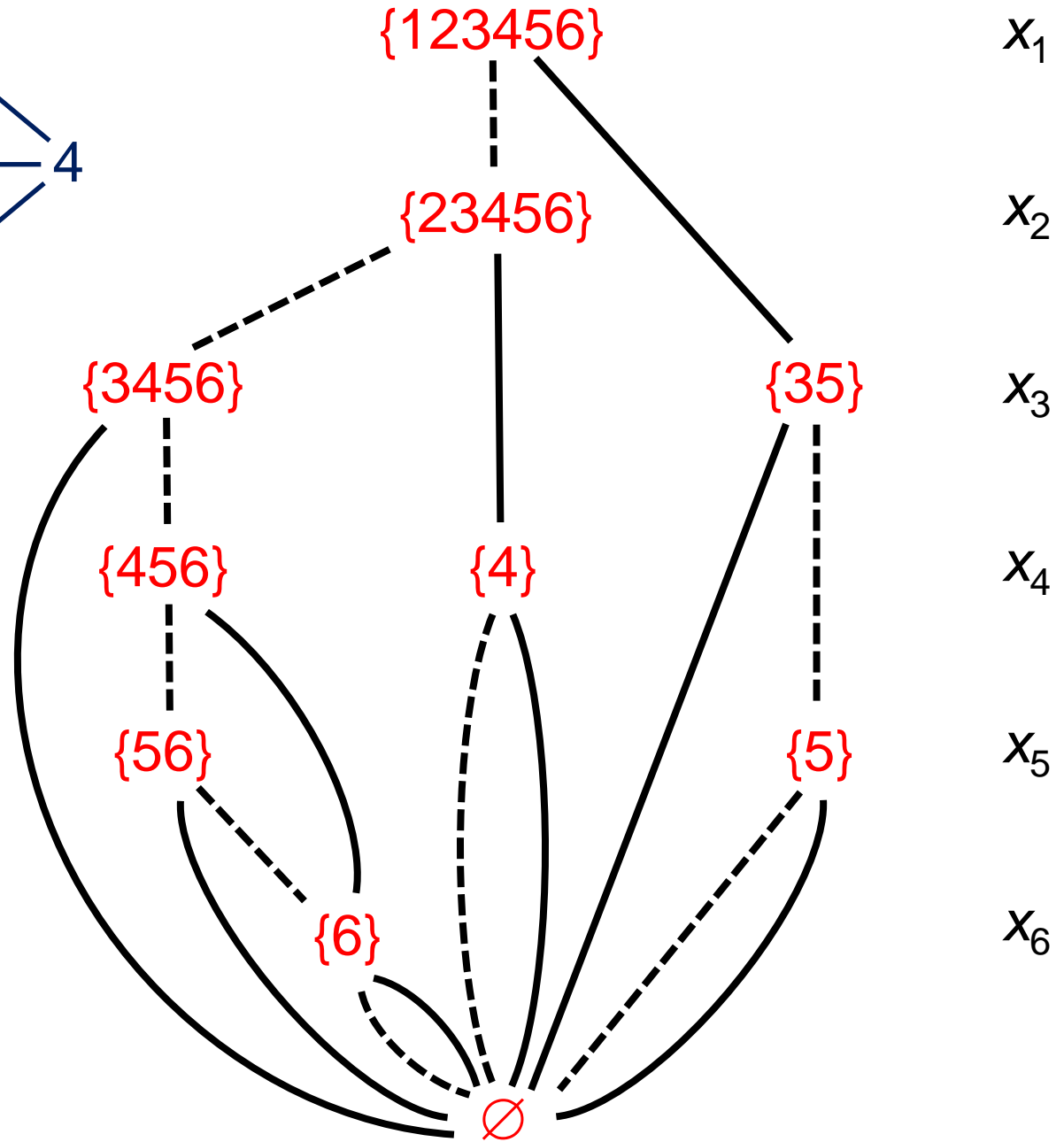
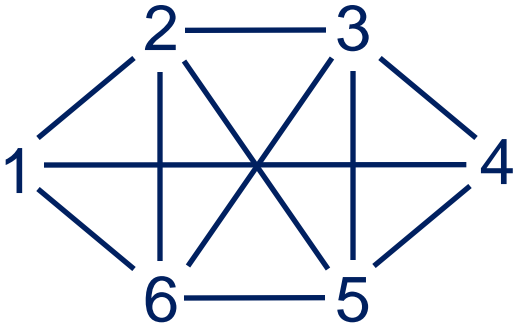
$x_6$

# Objective Function

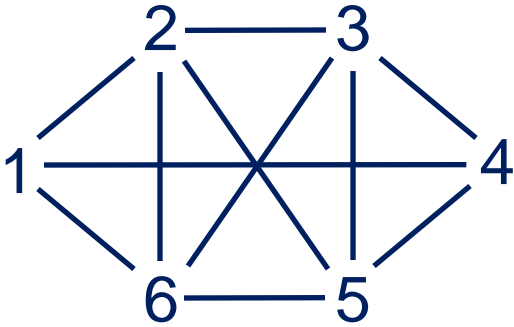
- In general, objective function can be any **separable function**.
  - Linear or nonlinear, convex or nonconvex
- BDDs can be generalized to **nonseparable** objective functions.
  - There is a unique reduced BDD with **canonical** edge costs.

# DP-Style Modeling

- To build a decision diagram, we need information about the problem.
  - We model the problem with **state variables**.
  - Rather than inequalities as in MIP.
  - Or global constraints as in CP.



To build BDD,  
associate **state**  
with each node



{123456}

$x_1$

$x_2$

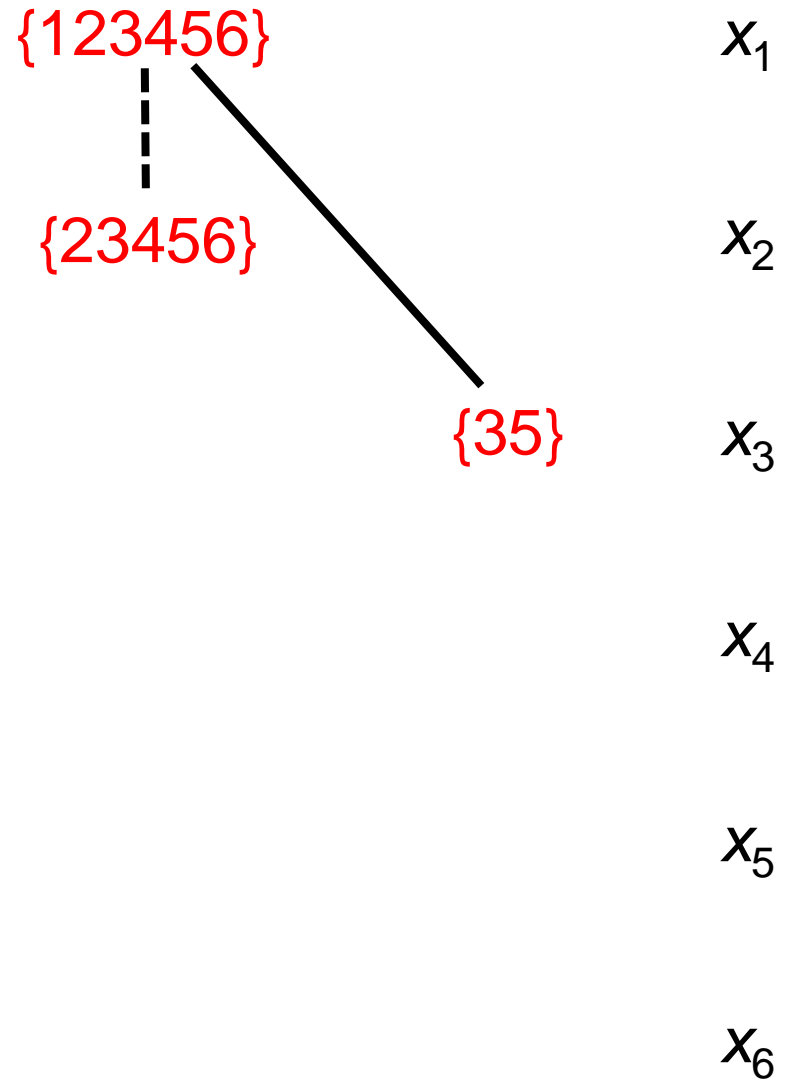
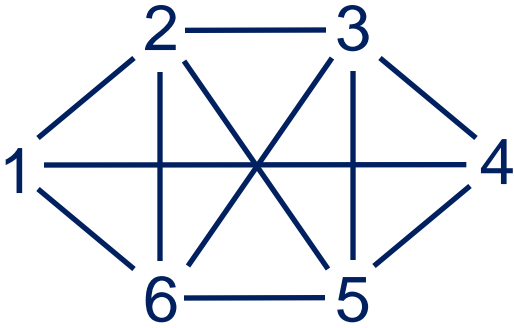
$x_3$

$x_4$

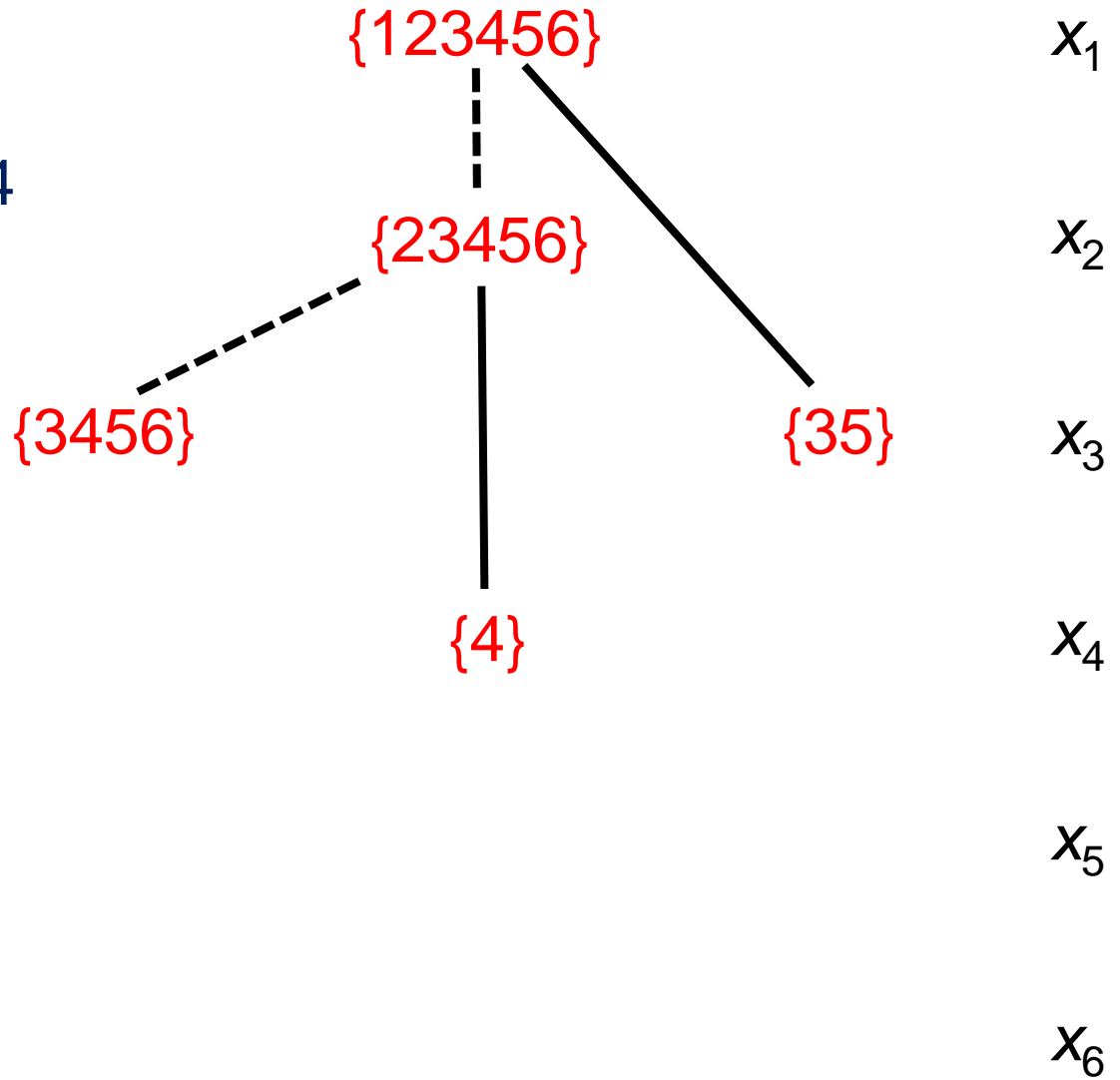
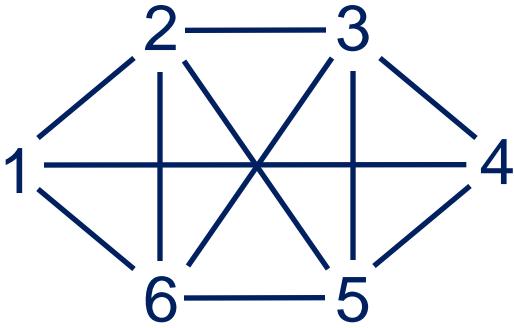
$x_5$

$x_6$

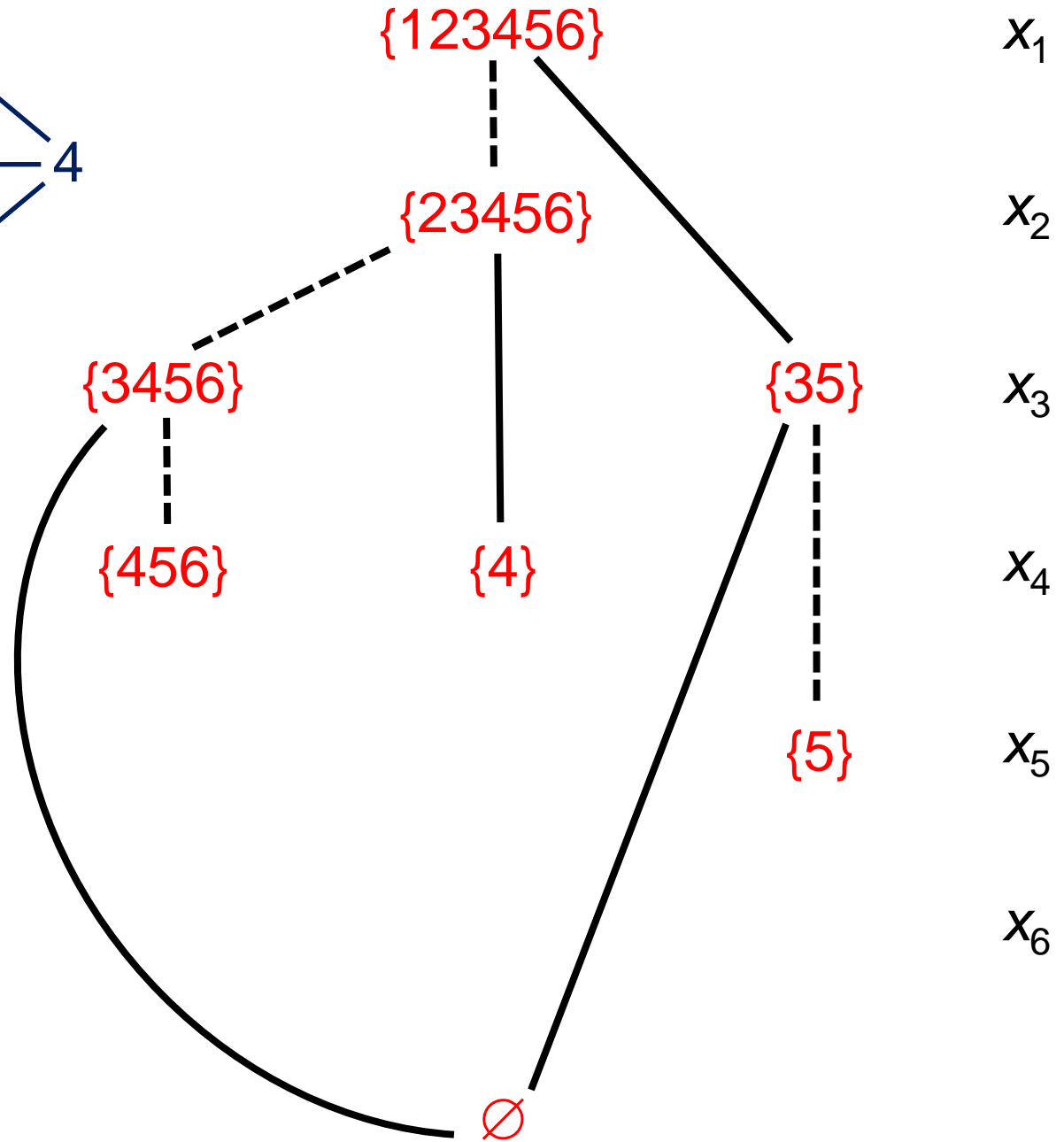
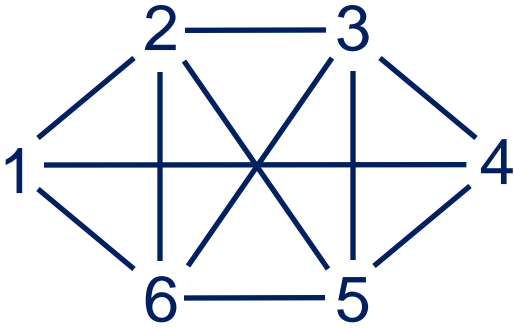
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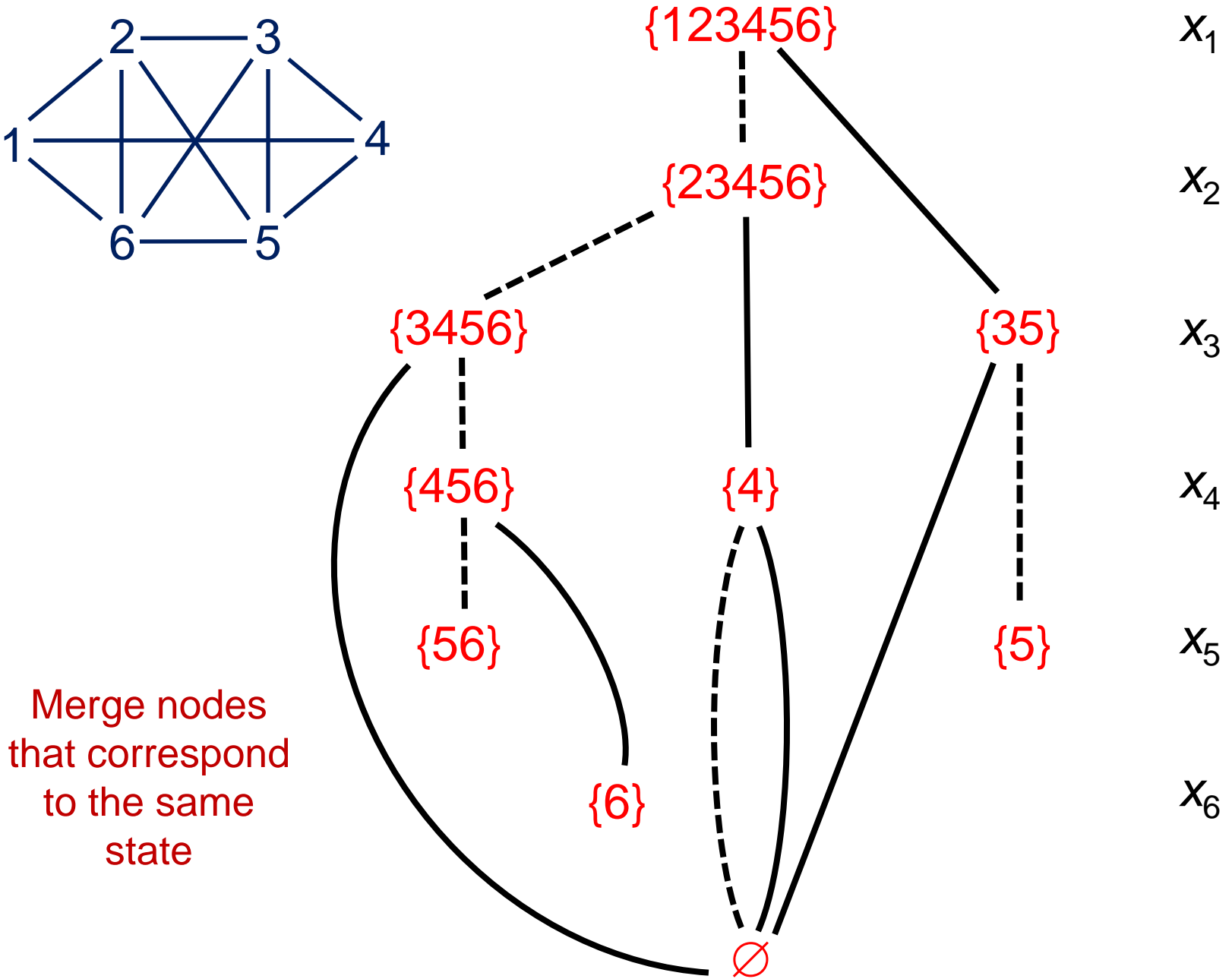
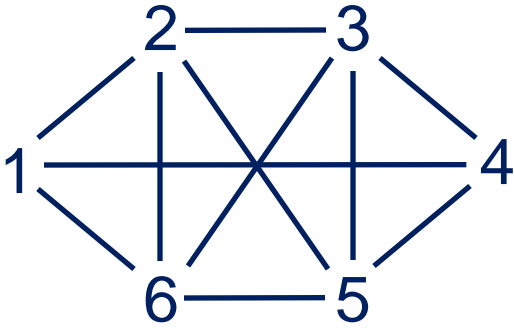


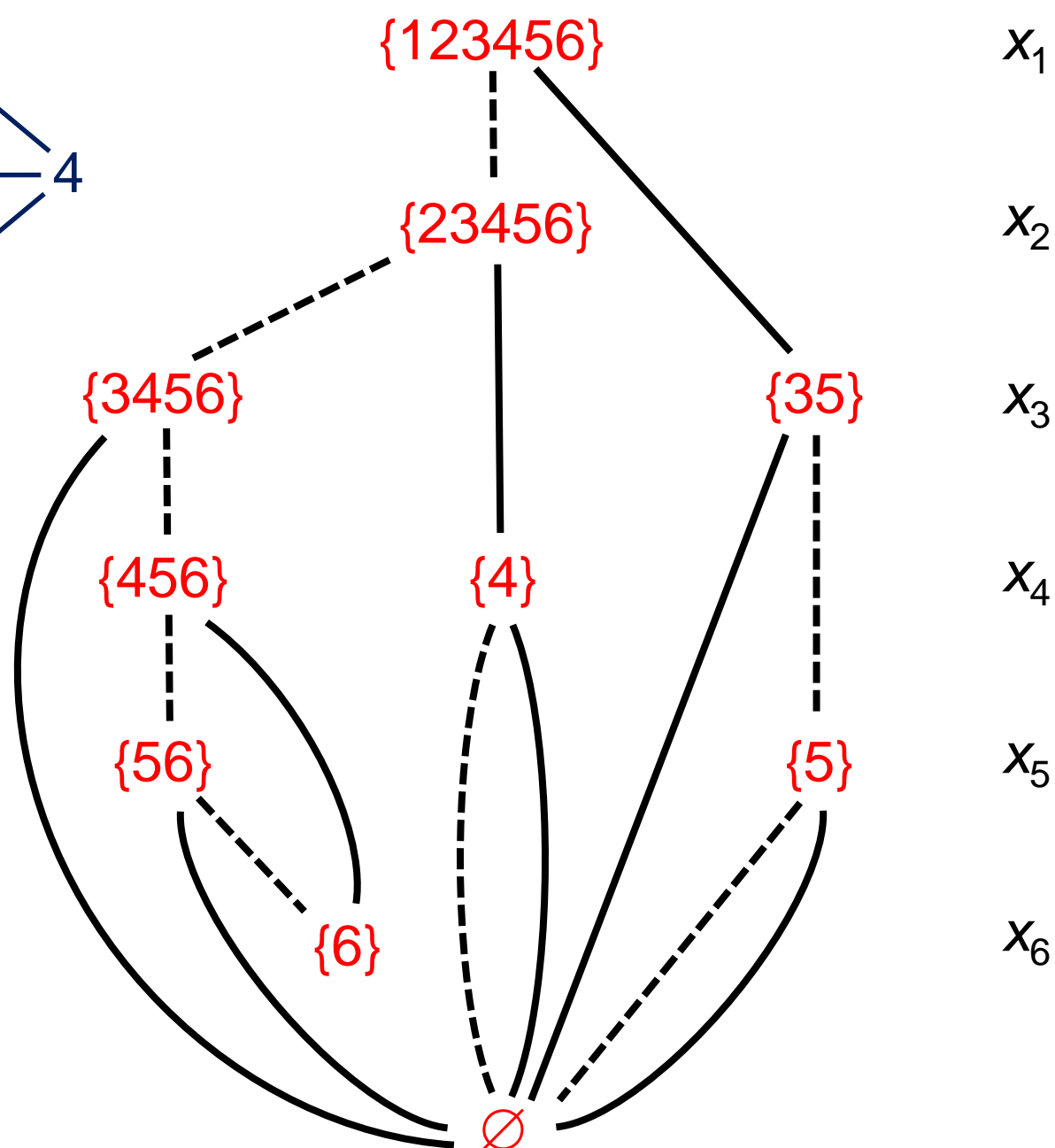
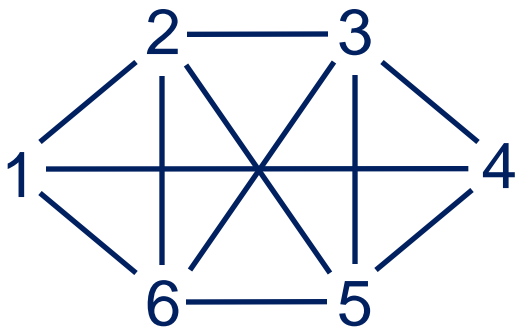
To build BDD,  
 associate **state**  
 with each node



Merge nodes  
that correspond  
to the same  
state

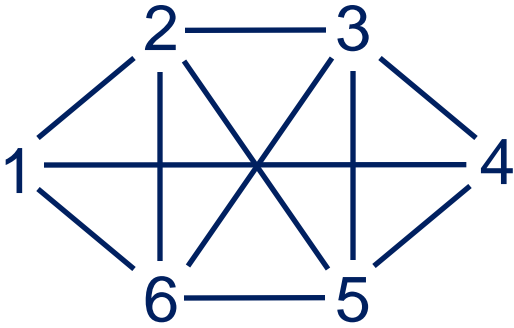






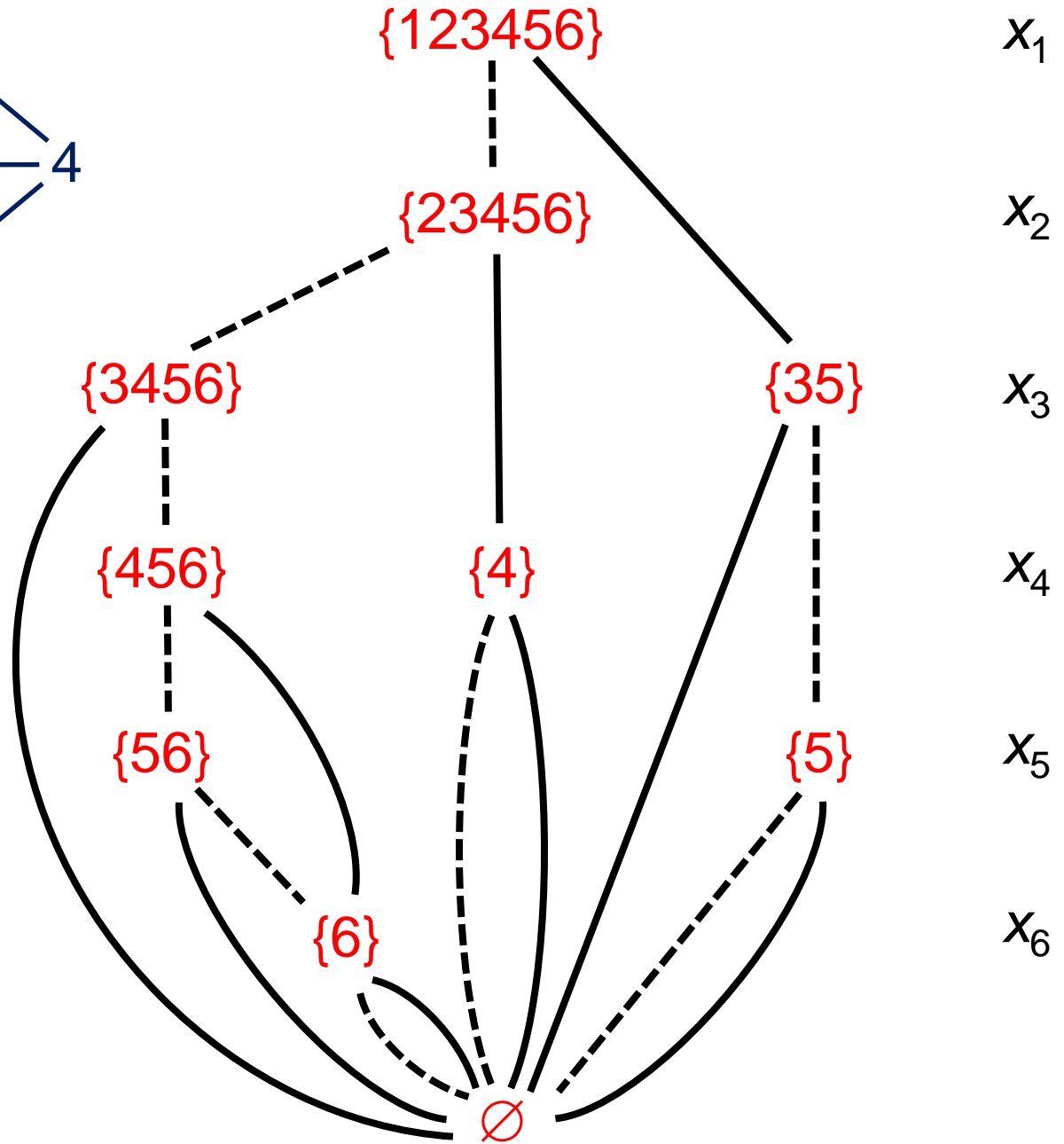
Merge nodes  
that correspond  
to the same  
state

$x_1$   
 $x_2$   
 $x_3$   
 $x_4$   
 $x_5$   
 $x_6$



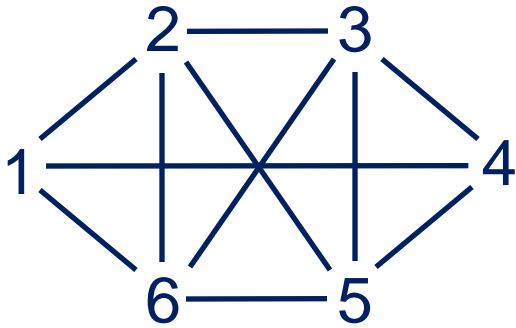
Width = 2

Merge nodes  
that correspond  
to the same  
state



# Relaxation Bounding

- To obtain a bound on the objective function:
  - Use a **relaxed** decision diagram
  - Analogous to linear programming relaxation in MIP
  - This relaxation is **discrete**.
  - Doesn't require the linear inequality formulation of MIP.



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$x_1$

$x_2$

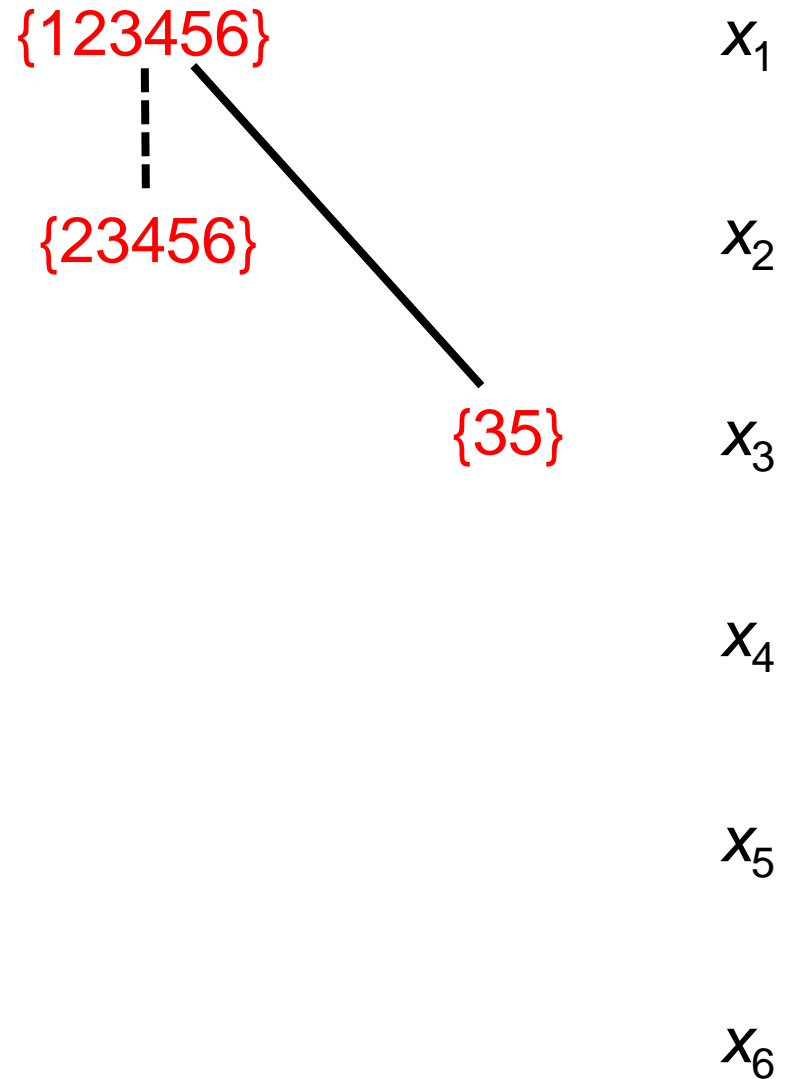
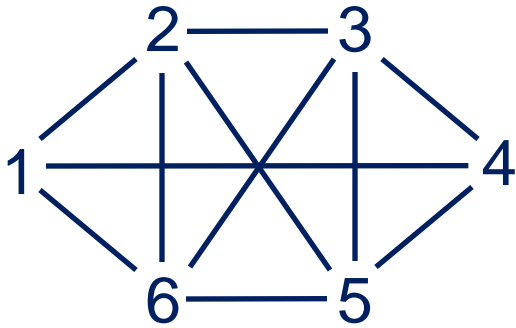
$x_3$

$x_4$

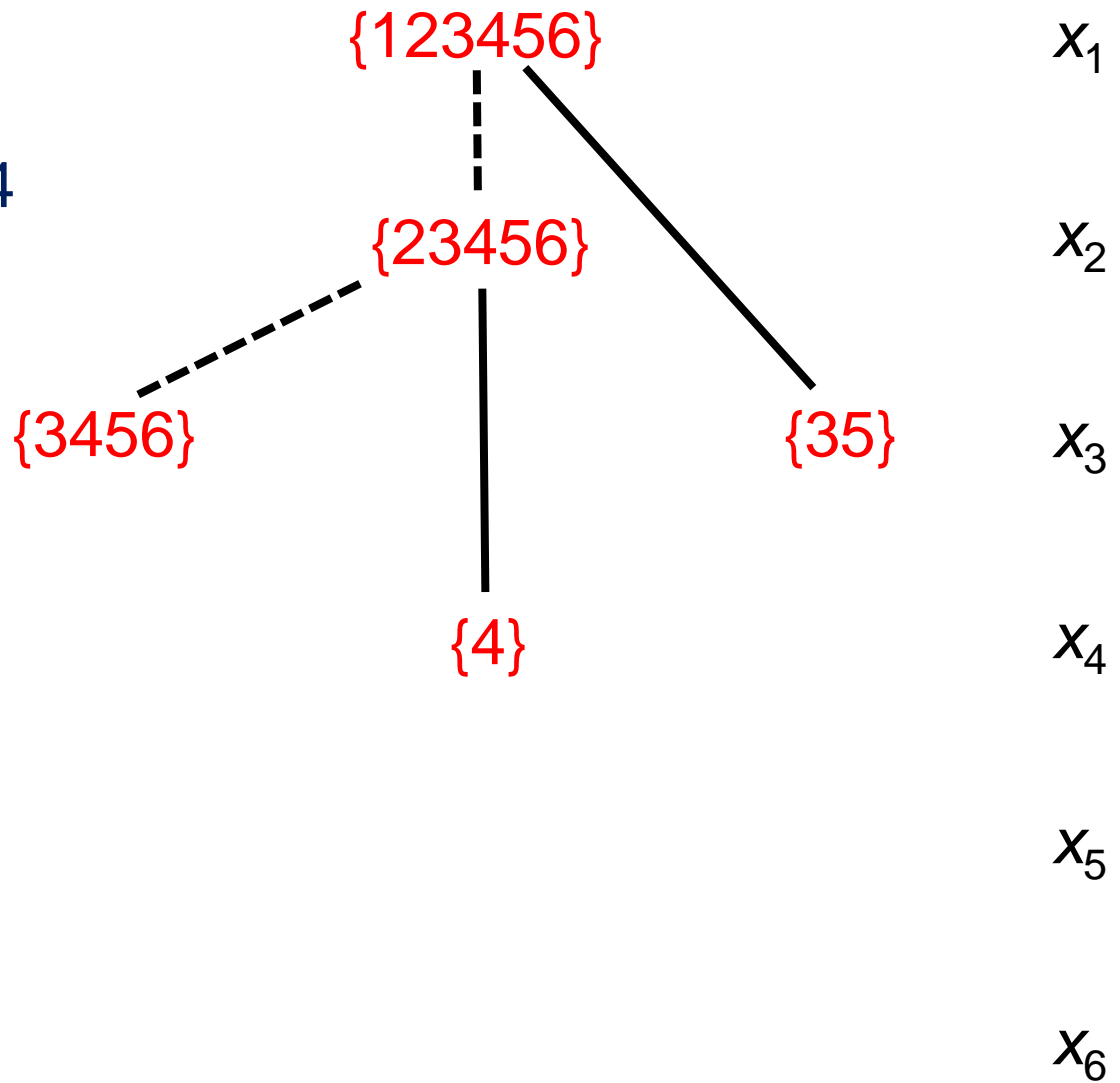
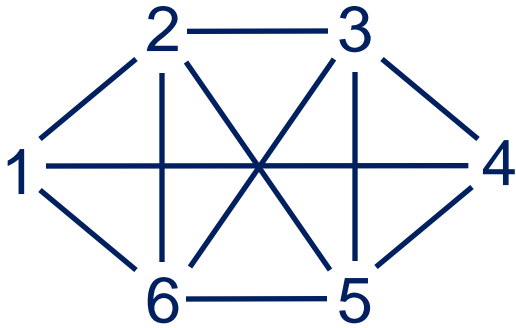
$x_5$

$x_6$

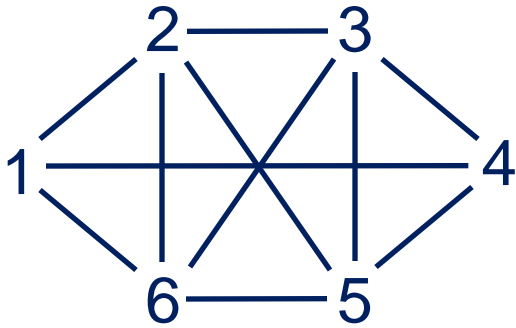
To build **relaxed**  
BDD, merge  
some additional  
nodes as we go  
along



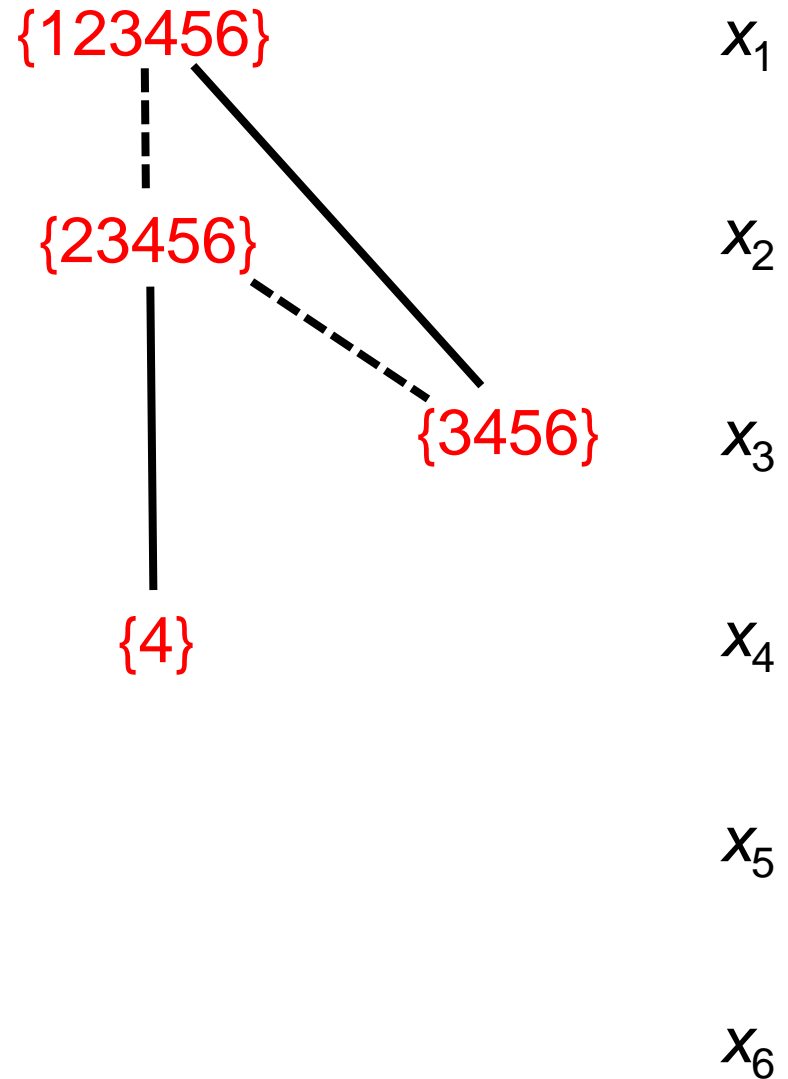
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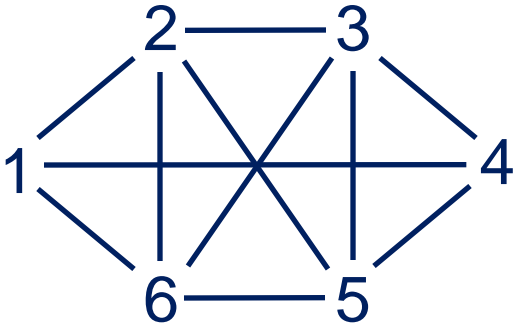
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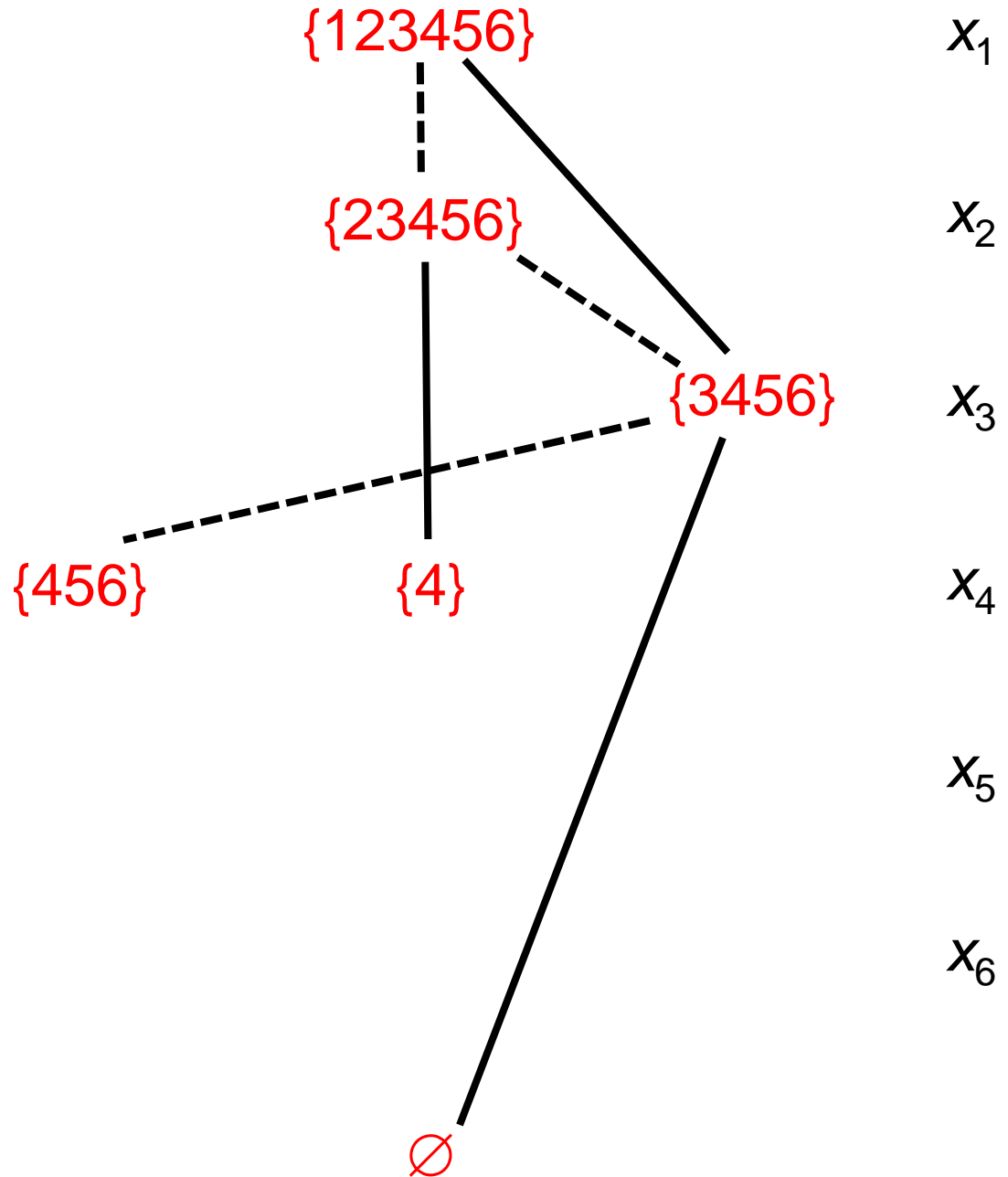
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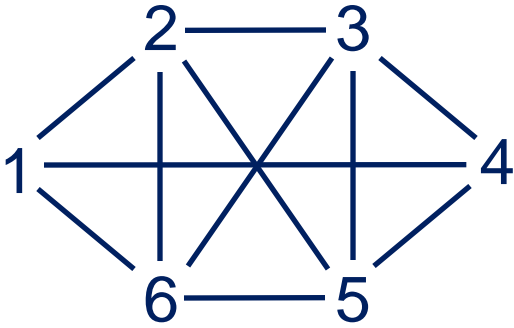




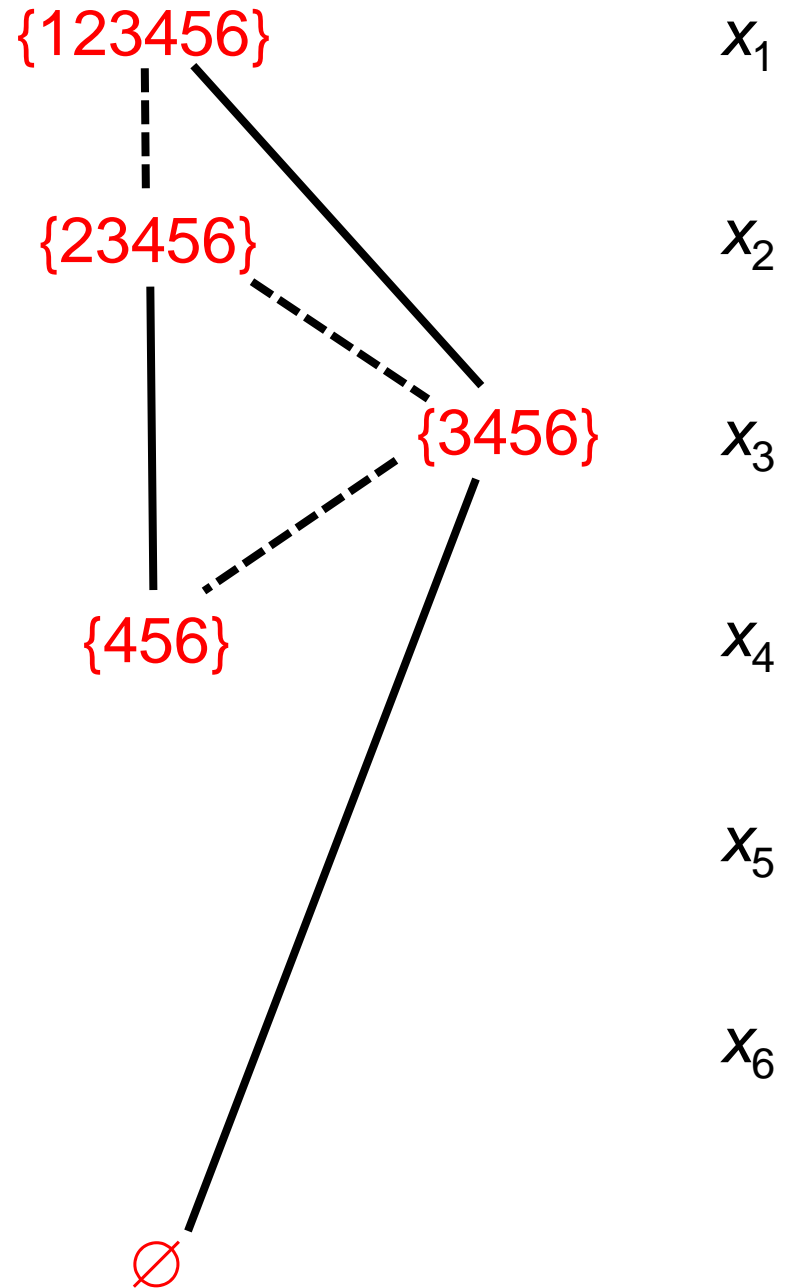


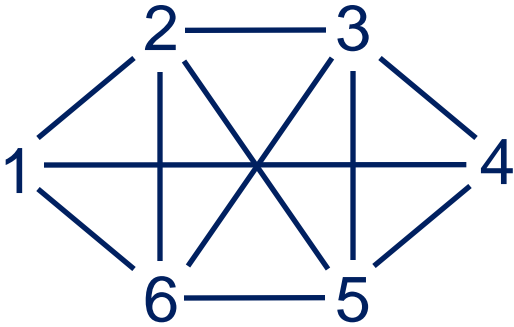
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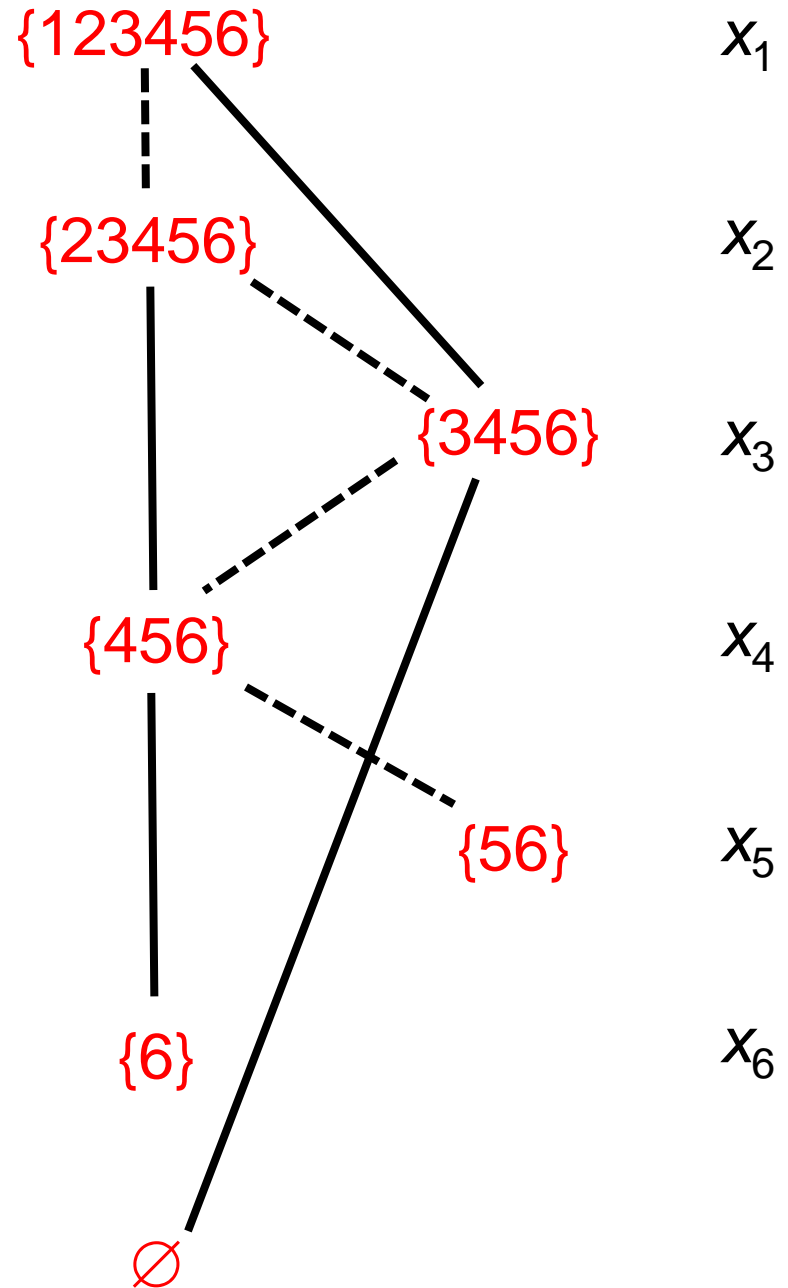


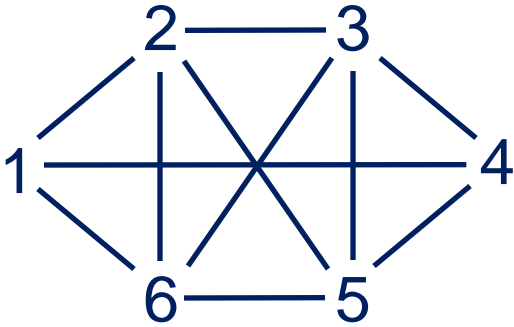
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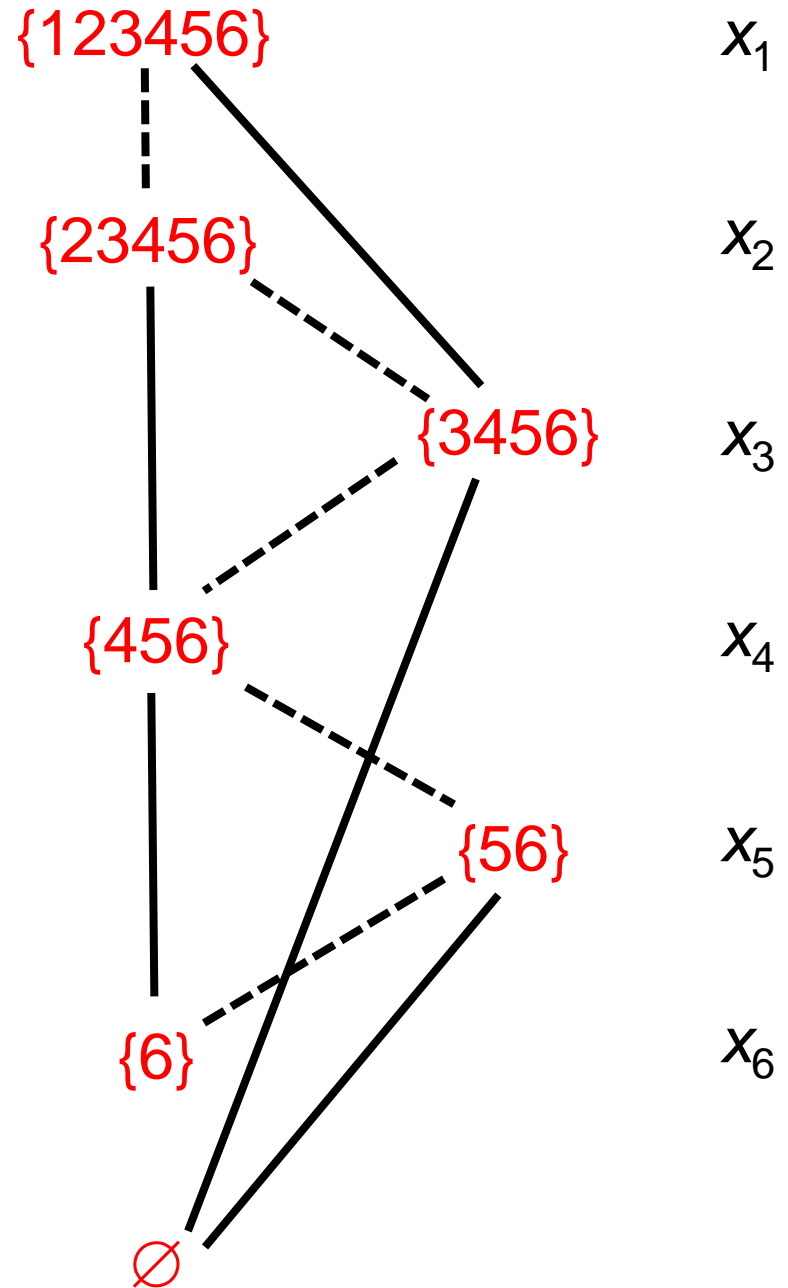


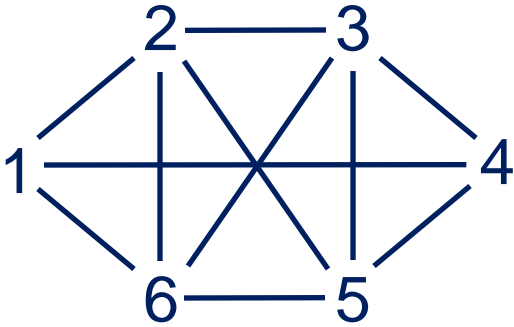
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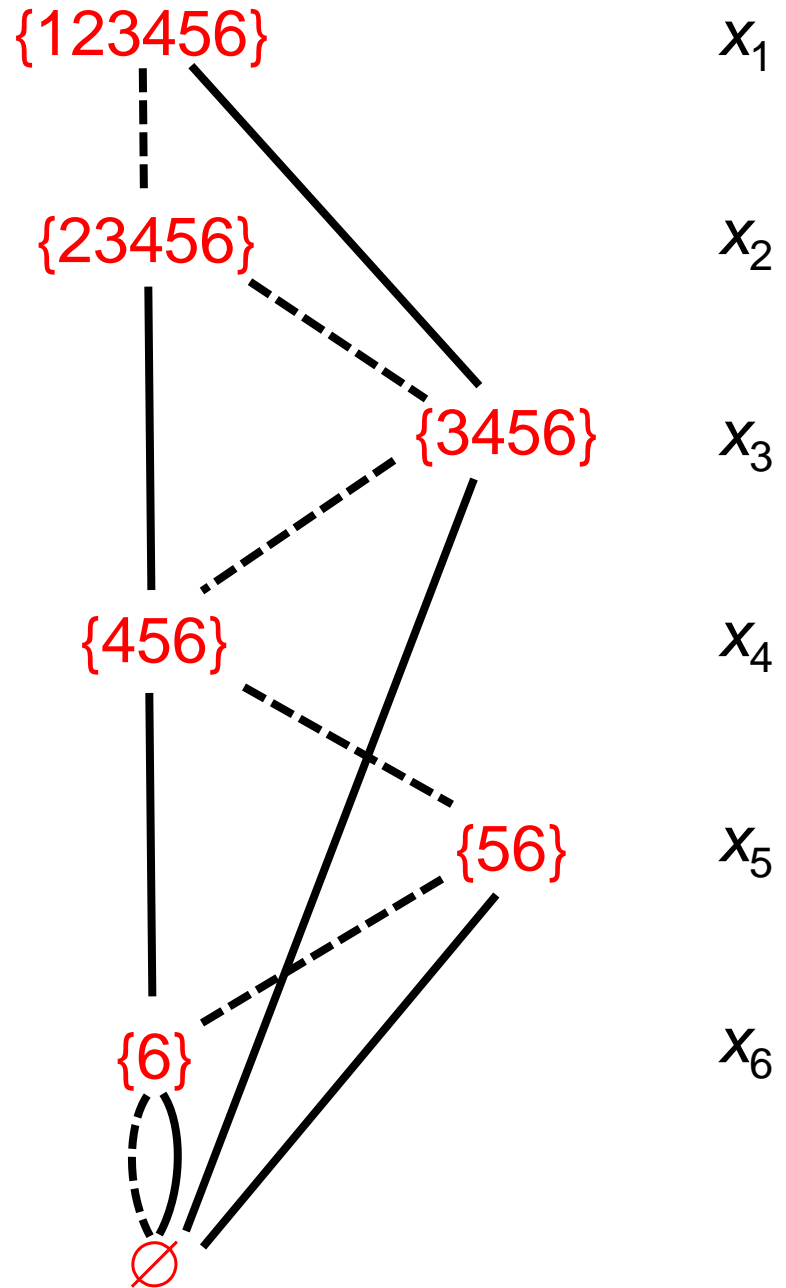
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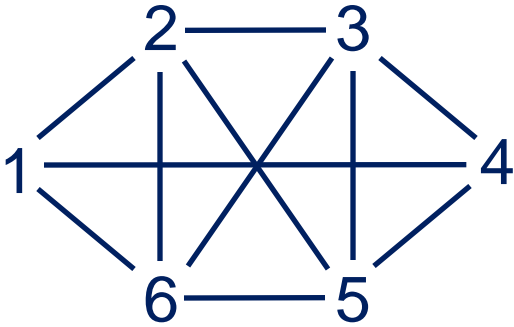




Width = 1

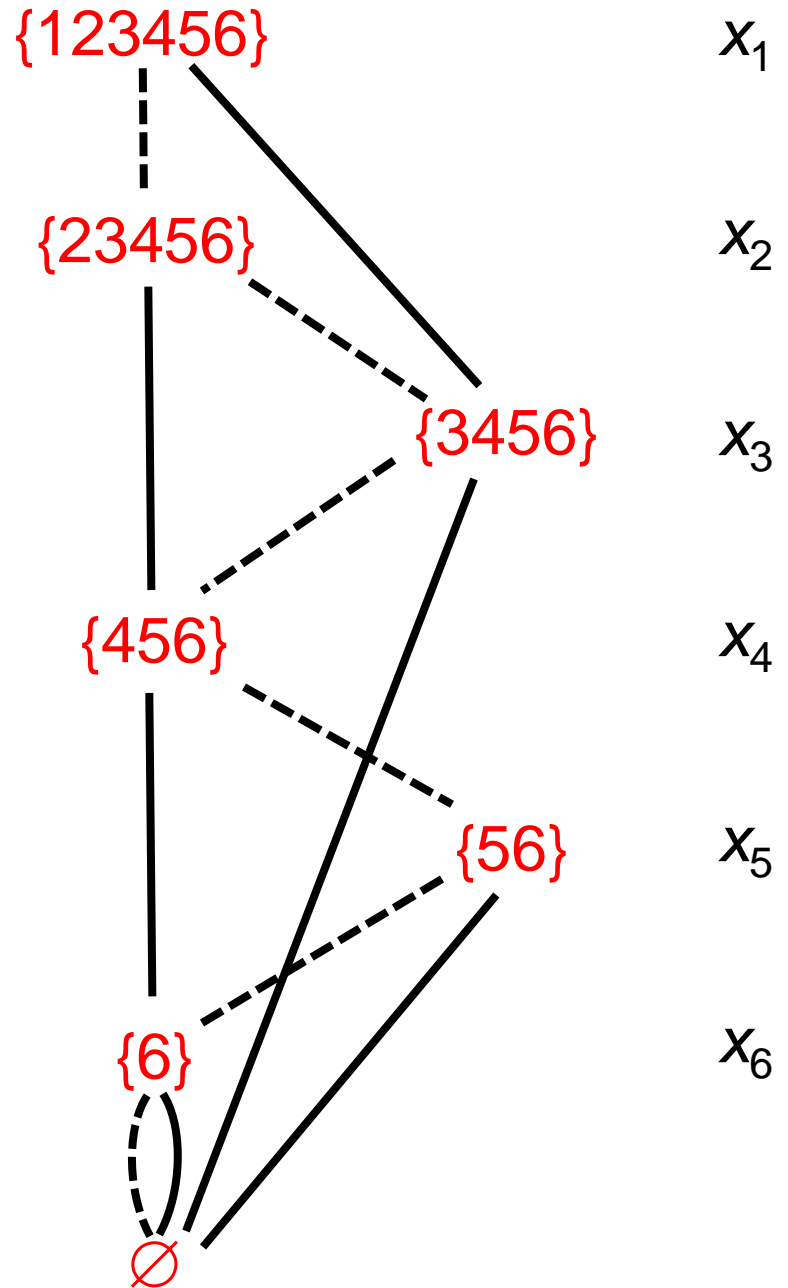
To build **relaxed**  
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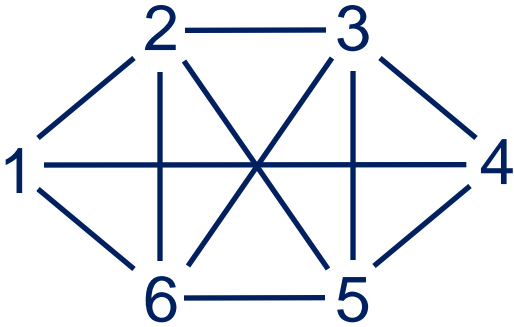




Width = 1

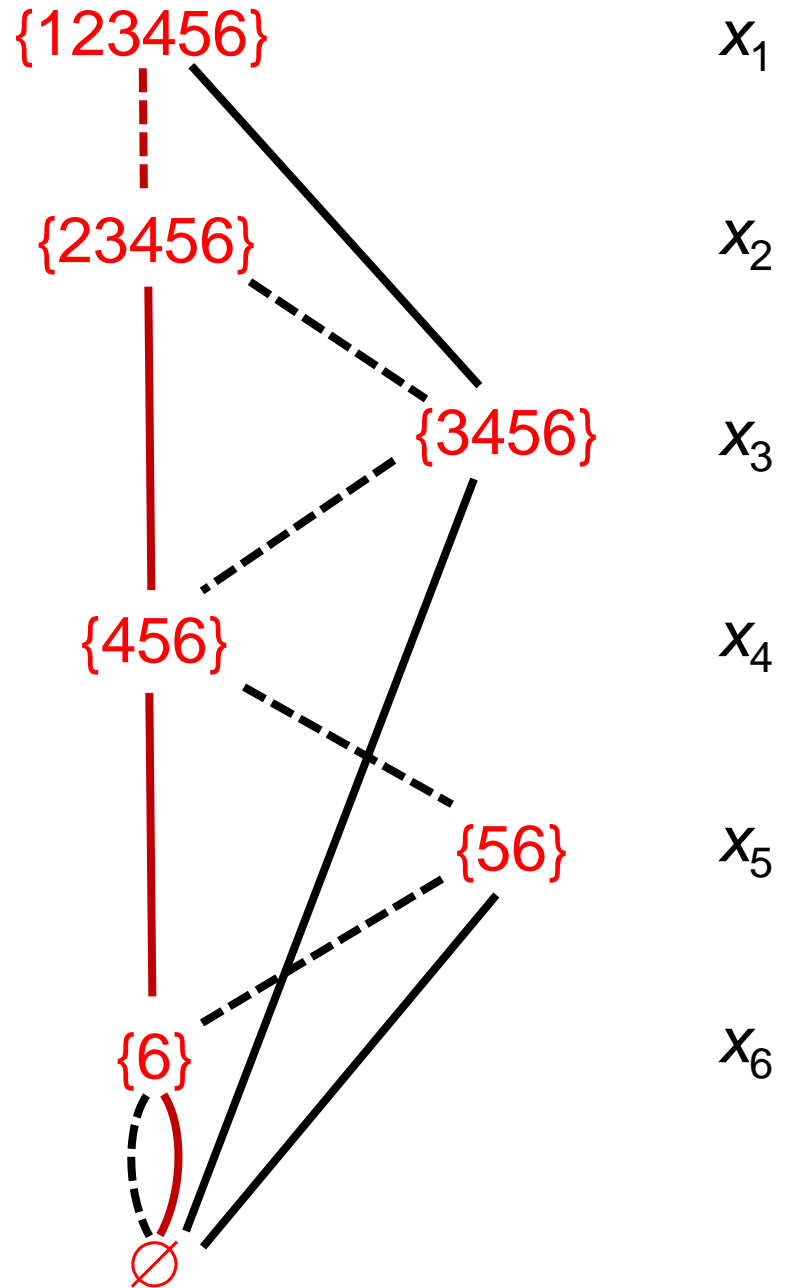
Represents 18 solutions,  
including 11  
feasible solutions





**Width = 1**

**Longest path**  
gives bound  
of 3 on optimal  
value of 2

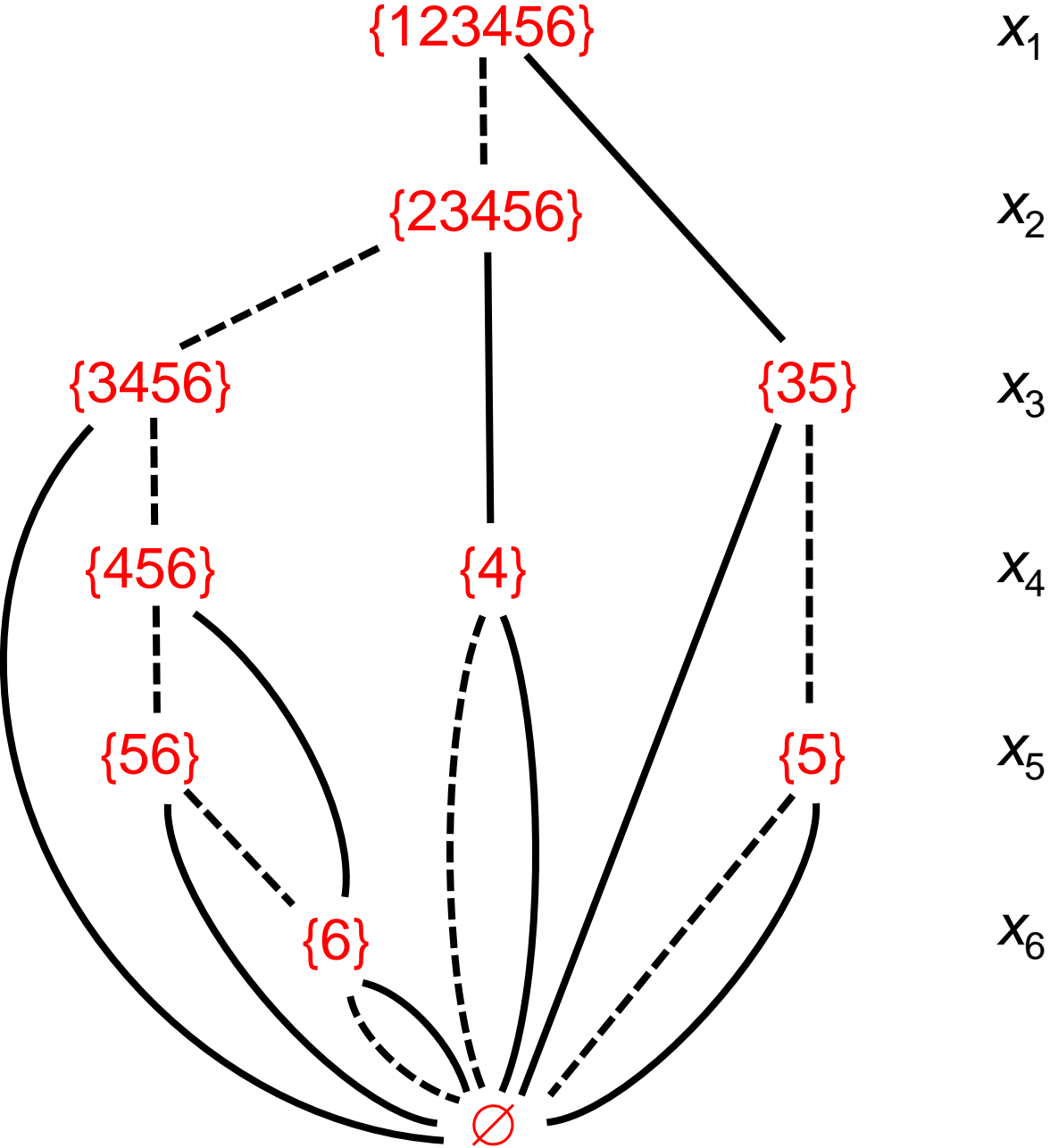


# Propagation

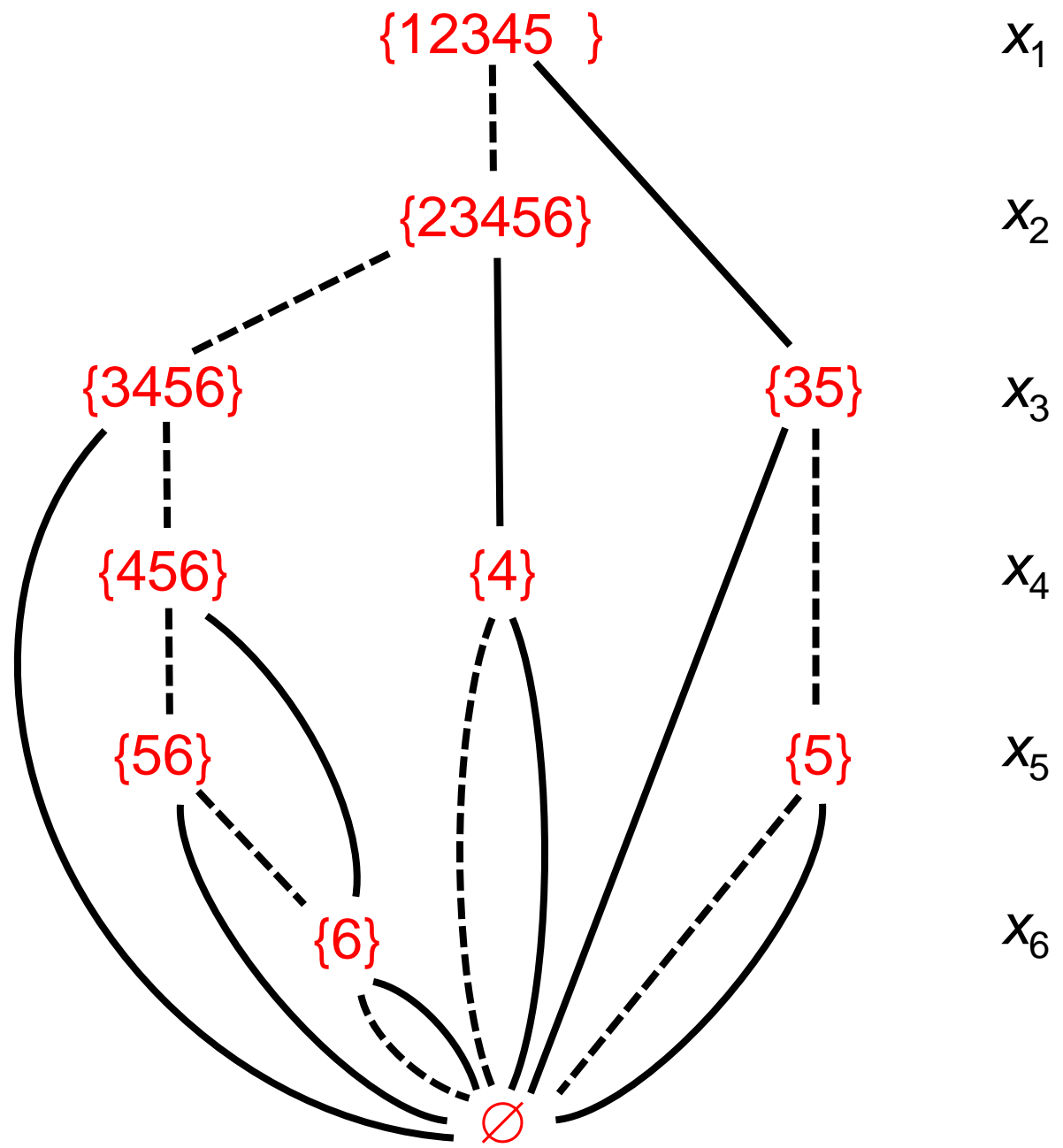
- We can propagate by removing arcs from the decision diagram.
  - Rather than removing elements from variable spans.
  - More effective than traditional domain filtering.
  - More information propagated from one constraint to the next.



Suppose this  
is the relaxed  
decision  
diagram

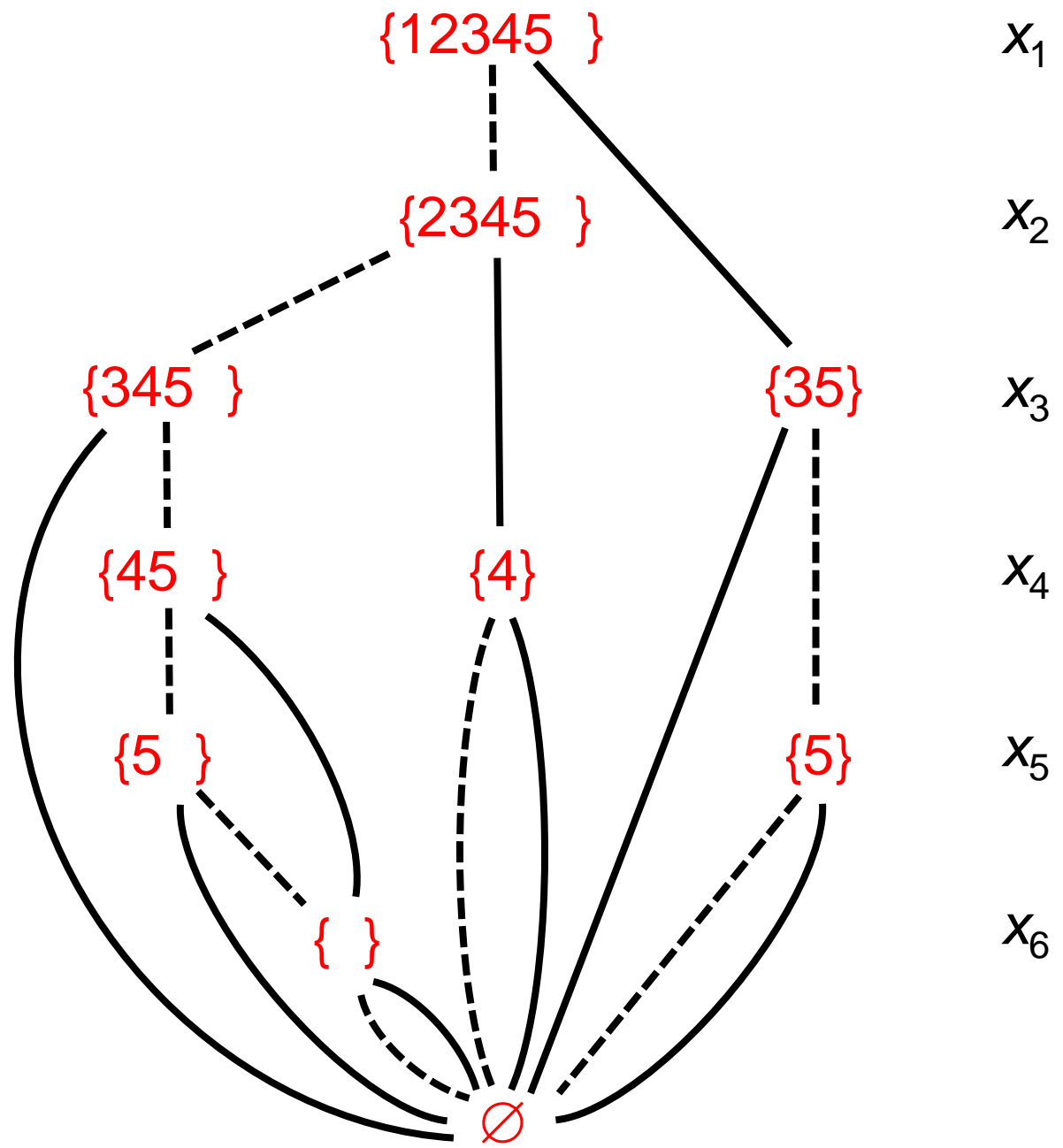


Suppose this is the relaxed decision diagram



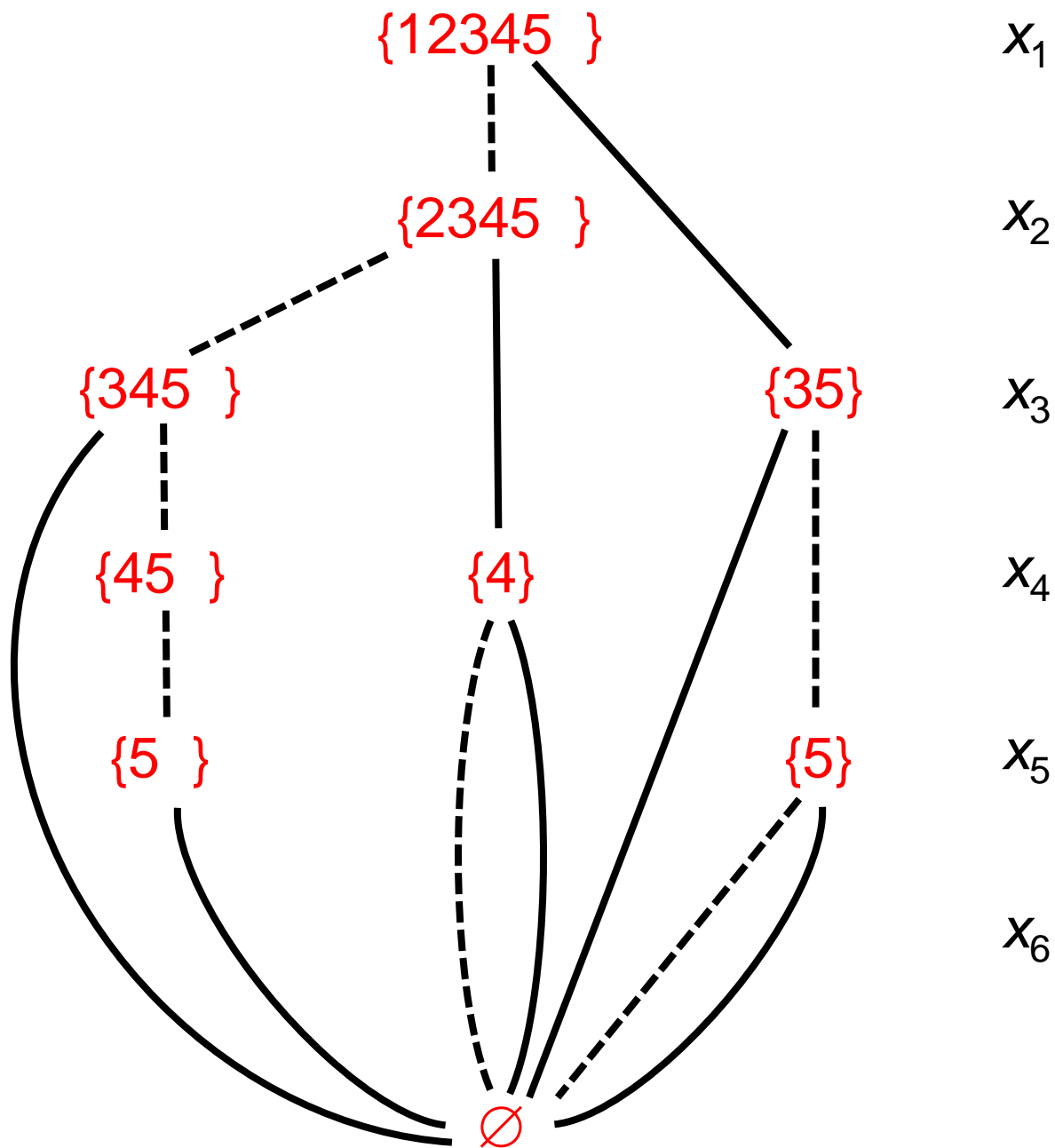
Suppose other constraints remove 6 from domain of  $x_1$

Suppose this  
is the relaxed  
decision  
diagram



This propagates  
through the  
states and  
removes some  
arcs.

Suppose this  
is the relaxed  
decision  
diagram



This propagates  
through the  
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# Variable Ordering

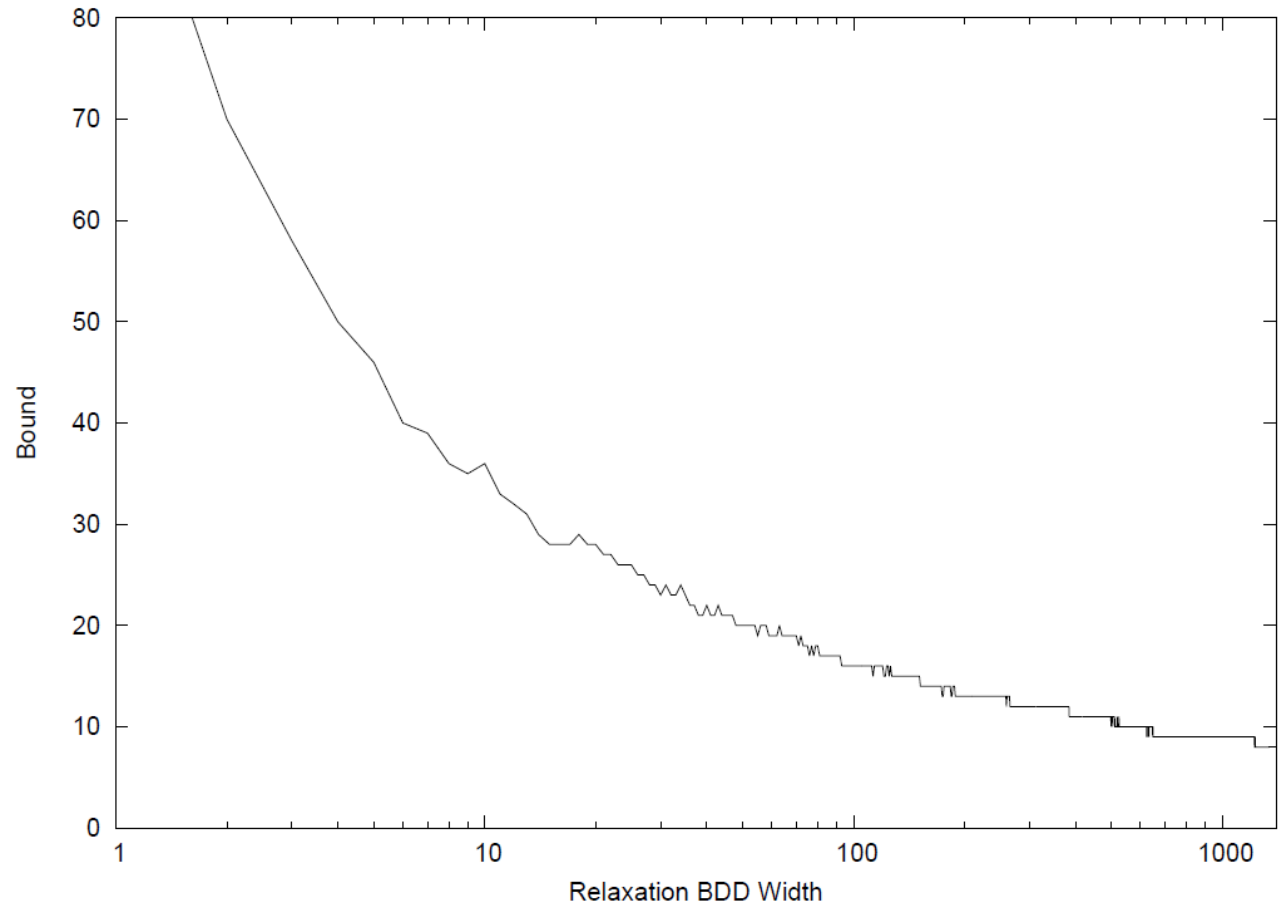
- Variable ordering is key.
  - Just as in branching methods.
  - Some orderings provide much tighter bounds.
- Width of exact decision diagram is bounded by Fibonacci numbers.
  - For an ordering induced by maximal path decomposition of the graph.
- We used a dynamic ordering heuristic.
  - Next variable is the one appearing in the smallest number of states on the current level.
    - Better than maximal path ordering.

# Merging Heuristic

- Which nodes to merge when building relaxation?
- A longest path heuristic seems by far the best.
  - Order nodes on current layer by increasing length of longest path into each node.
  - Merge nodes in this order.
  - We lose information in areas not likely to be part of the solution.
- This is better than merging nodes with more vertices in the corresponding states.

# Width of Decision Diagram

- Wider BDDs yield tighter bounds.
  - But take longer to build.



# Comparison with LP Bound

- Random and benchmark instances
- Compare with LP bound at root node in CPLEX
  - Use clique cover IP formulation
    - Requires precomputing clique cover.

$$\max \sum_i w_i y_i$$

$$\sum_{i \in C_k} y_i \leq 1, \quad \text{all cliques } C_k \text{ in clique cover}$$

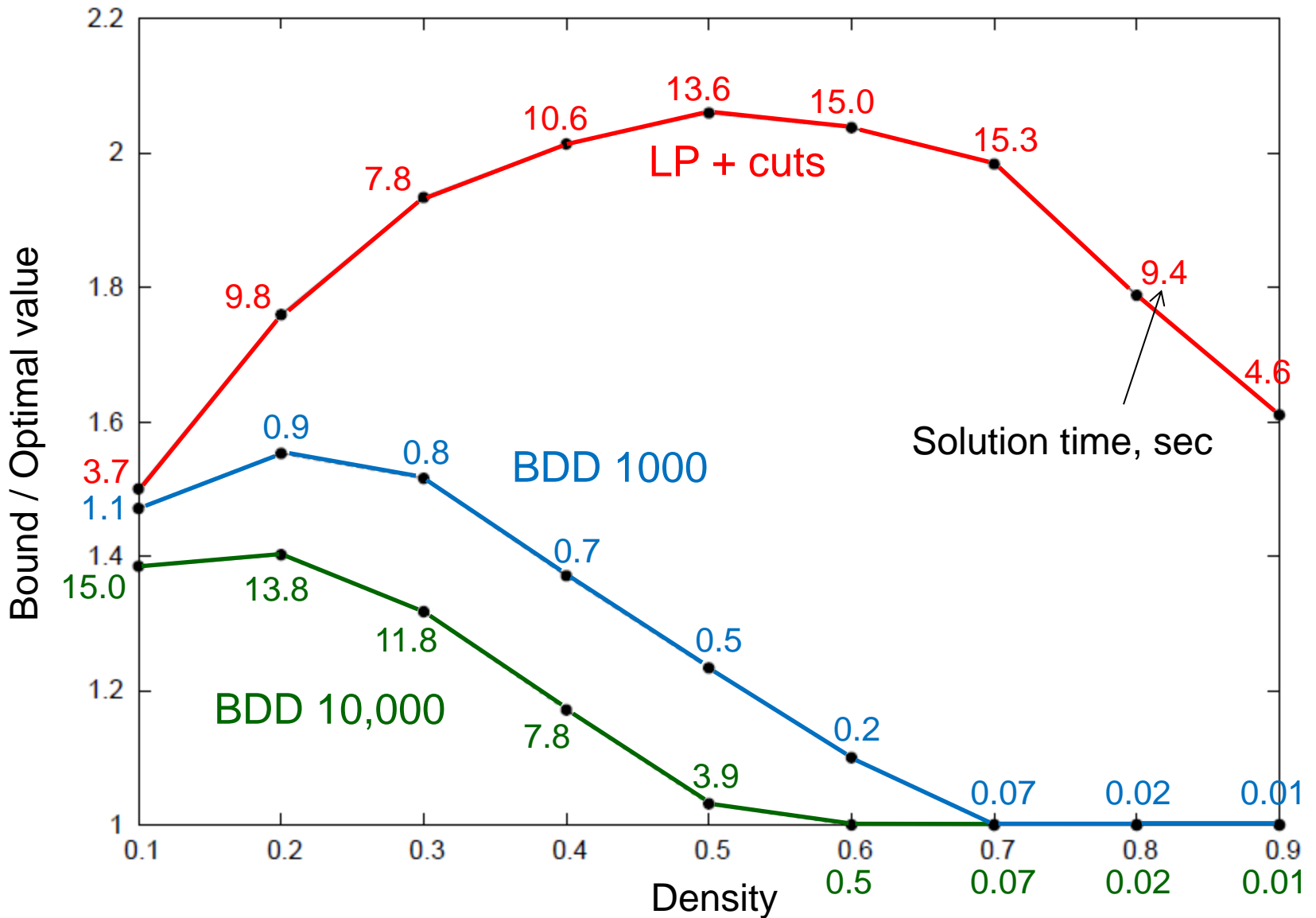
$$y_i \in \{0,1\}, \quad \text{all } i$$



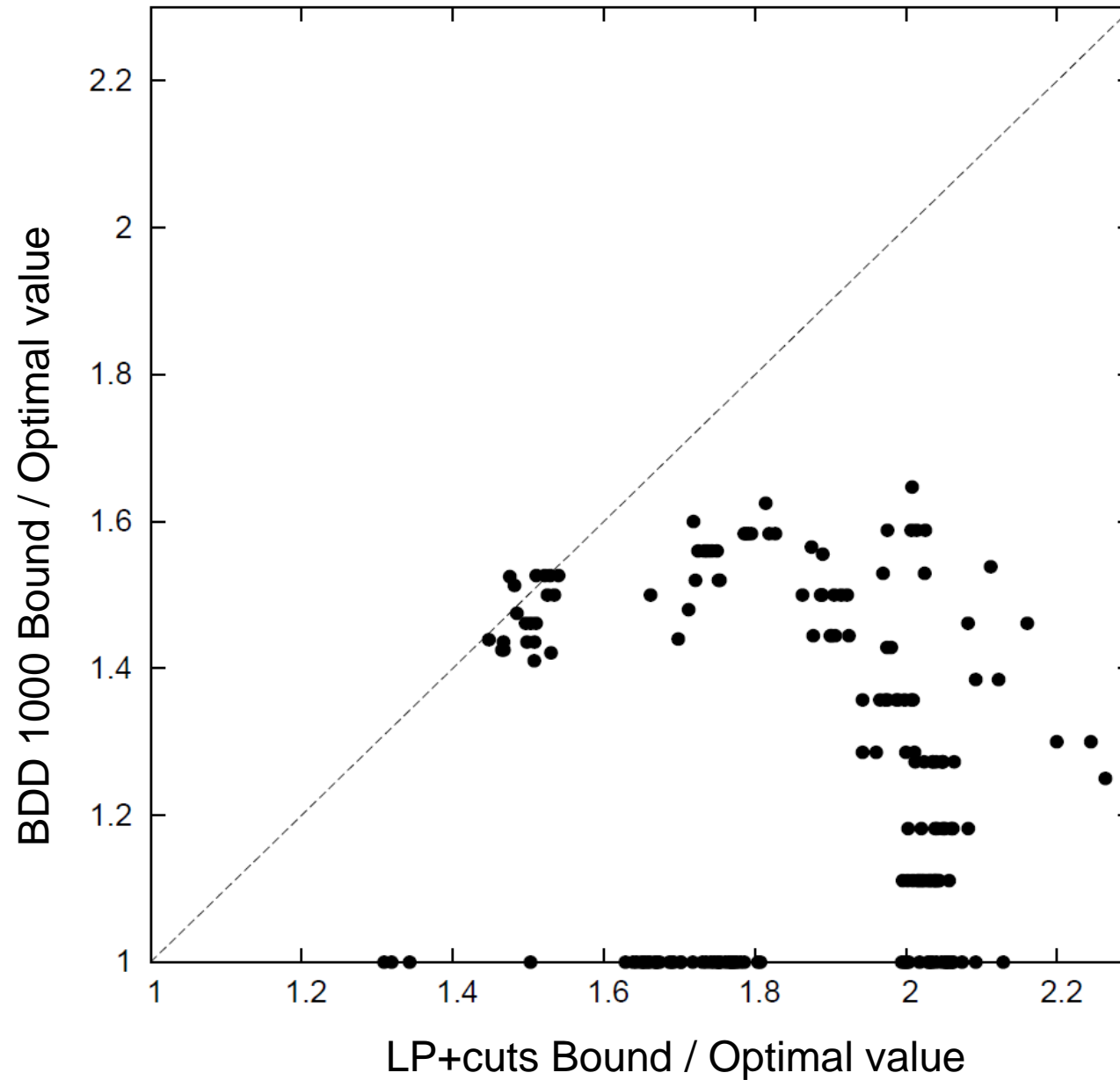
# Comparison with LP Bound

- CPLEX settings
  - Turn off presolve
    - It makes CPLEX bound worse or at most 1 better.
    - BDD bounds don't use presolve.
  - Use full cutting plane resources
  - Use interior point (barrier) LP solver
    - Faster than simplex on these instances.

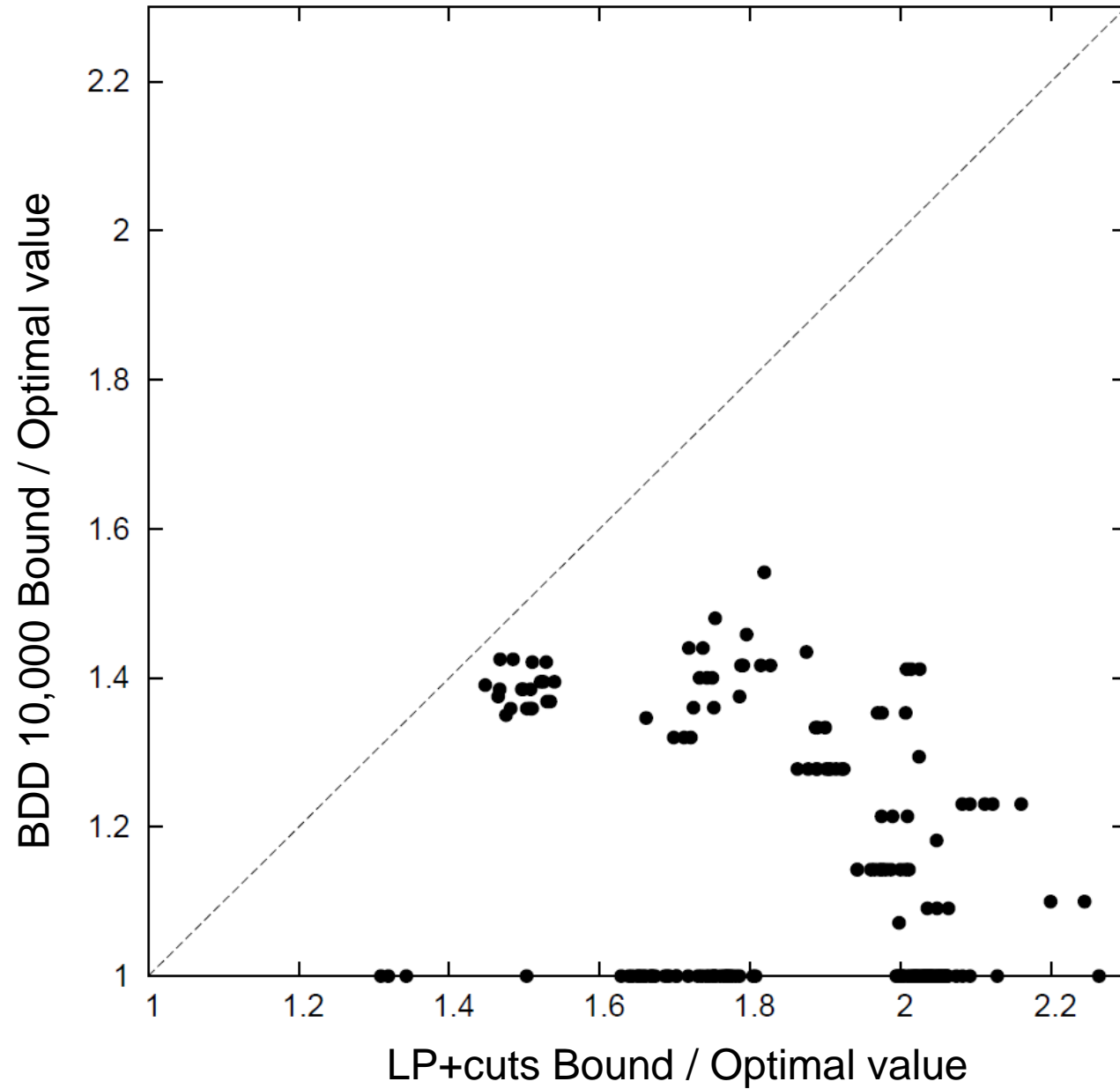
# Random instances



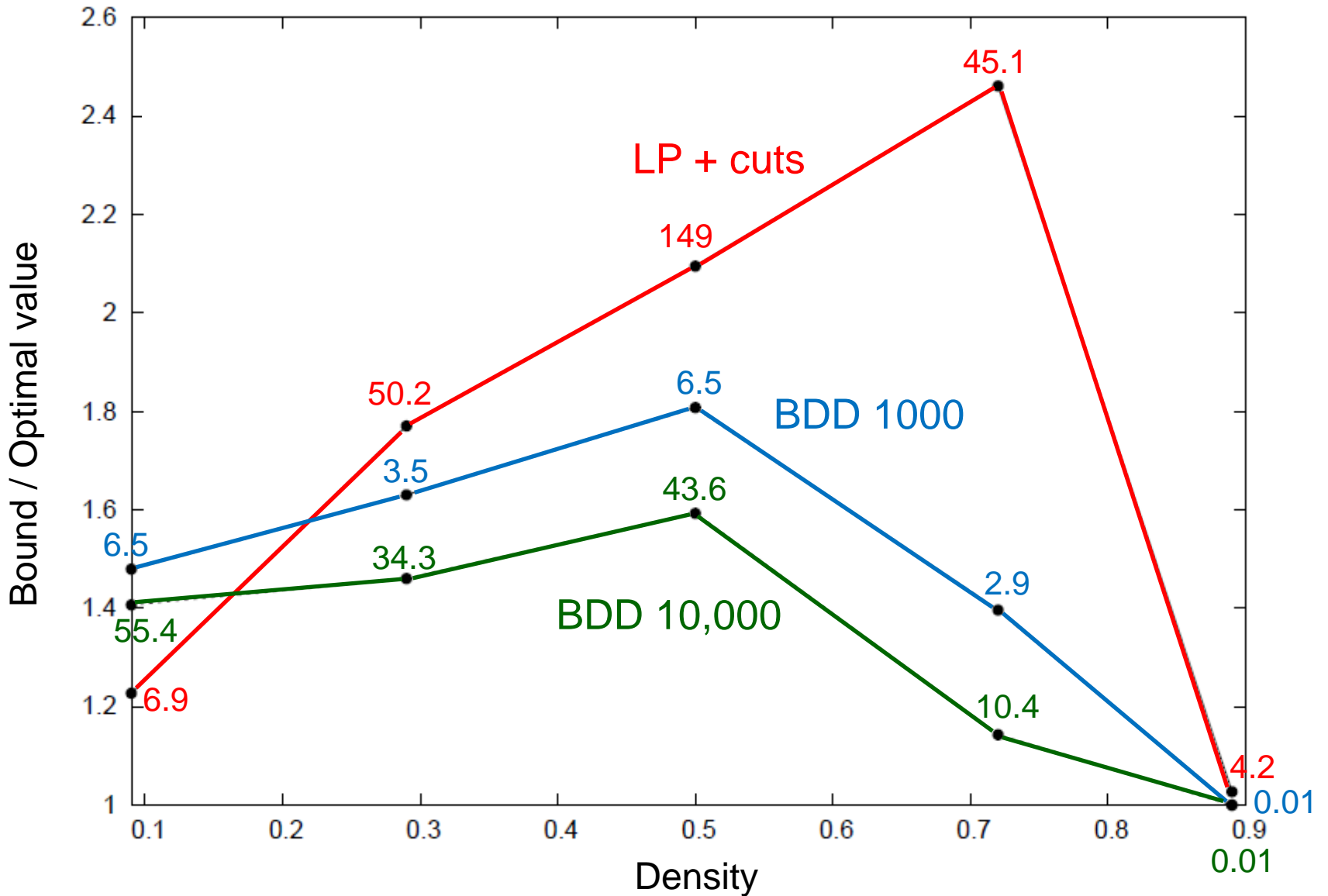
# Random instances



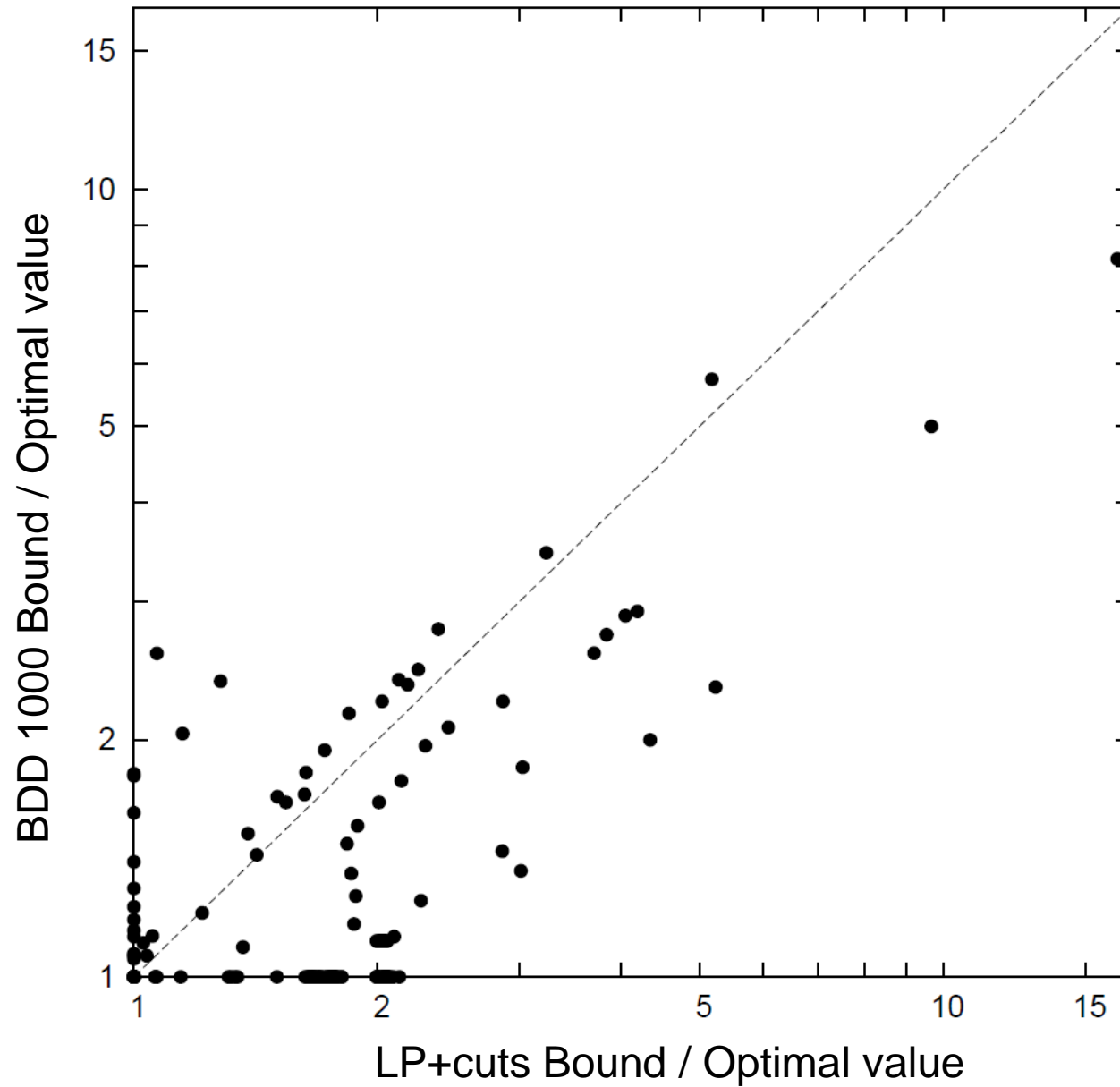
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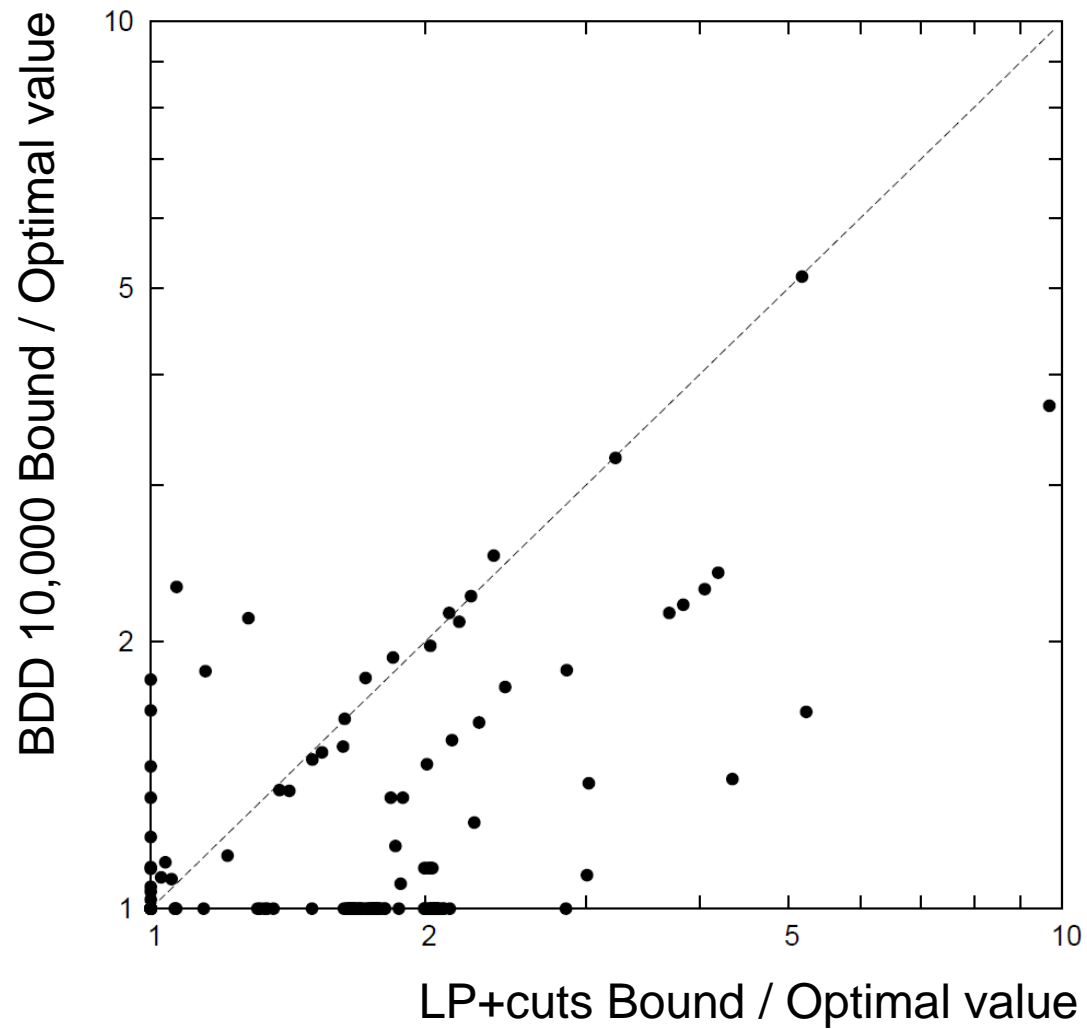
# DIMACS instances



# DIMACS instances



# DIMACS instances



# Incrementality

- Fast incremental calculation of bound.
  - Modify relaxed BDD after variable is fixed in branching tree (fast).
  - Recompute longest path (fast).
- However, may be better to rebuild relaxation from scratch.
  - Research issue.

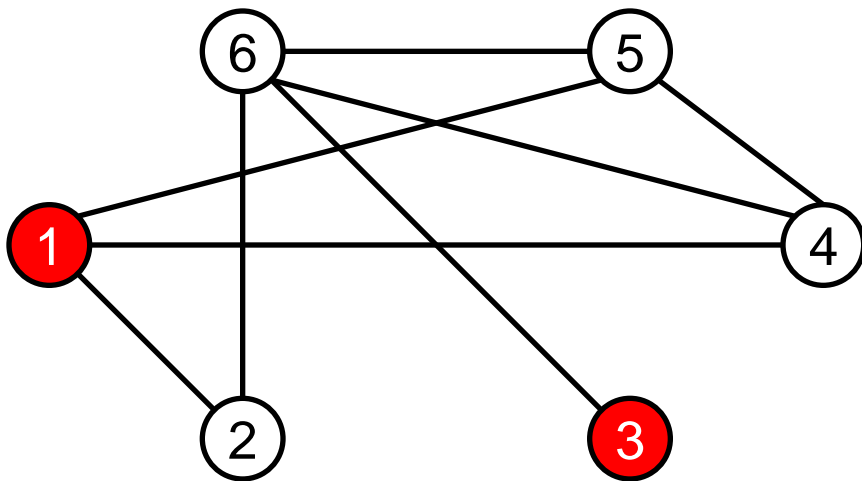


# Primal Heuristic

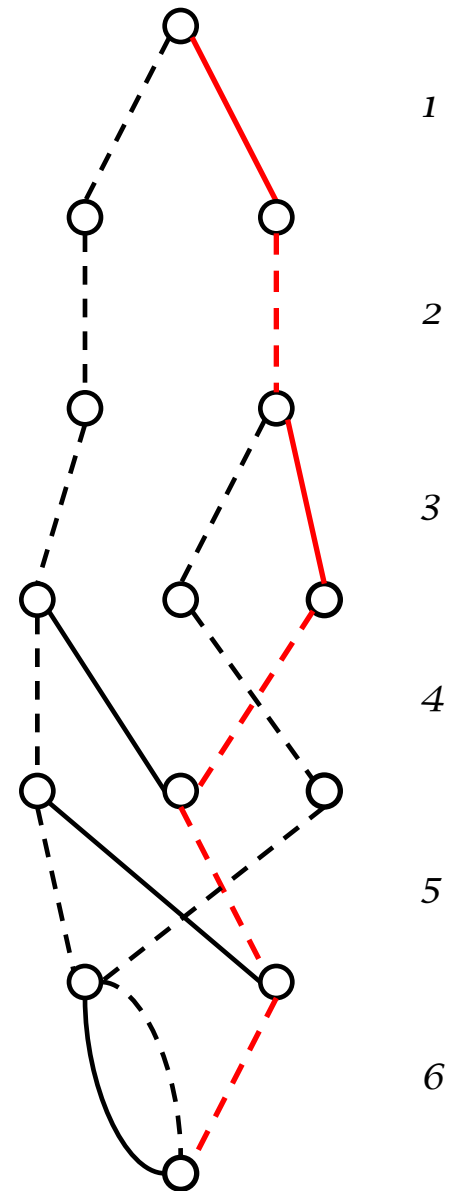
- Use **restricted** decision diagram.
  - All paths are feasible.
  - Longest path is a good feasible solution
  - Reduce width by intelligently removing nodes.

# Restricted BDD

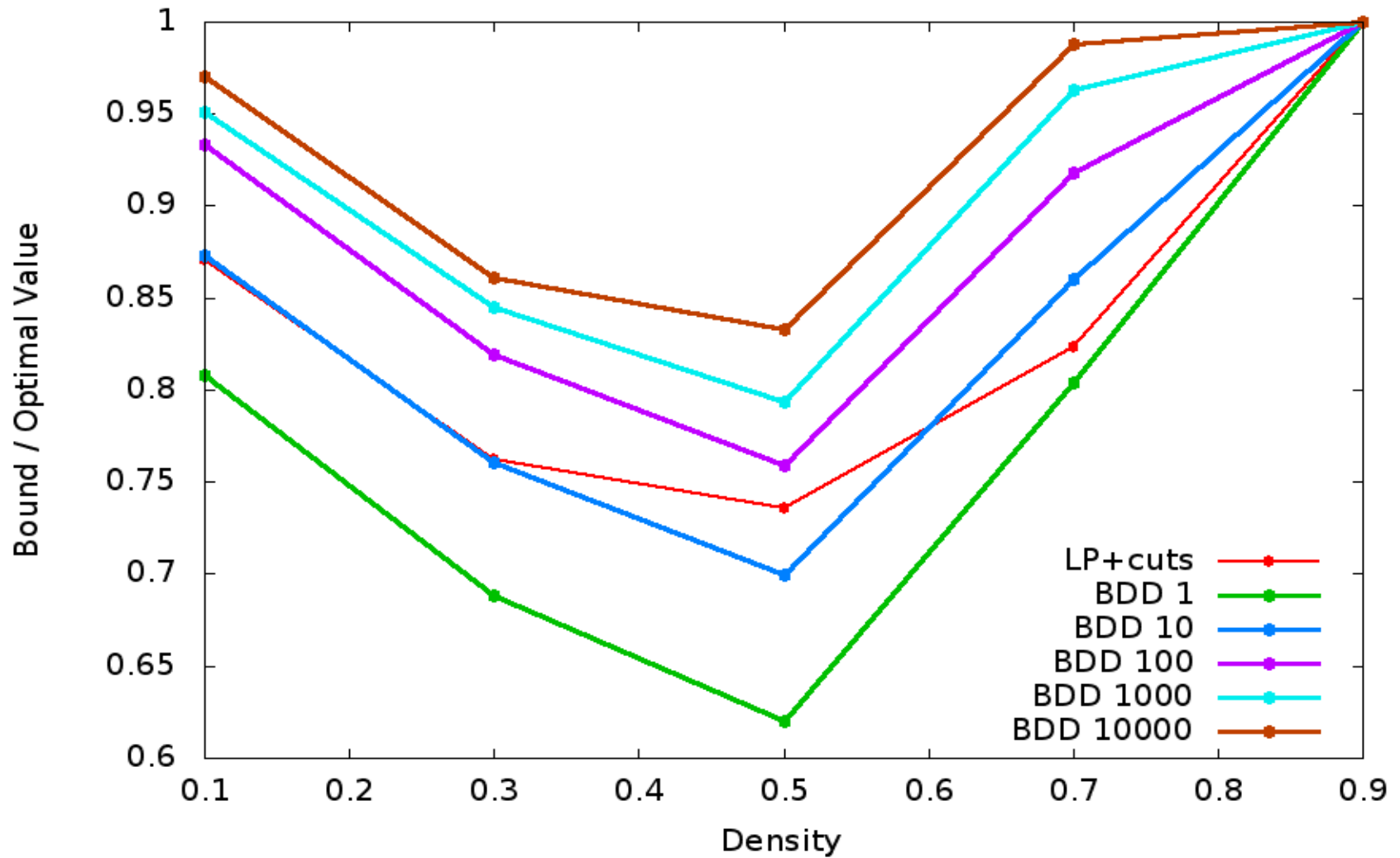
Width = 3



Longest path is  
**feasible solution**  
and gives upper bound  
of **2** on optimal value of **3**

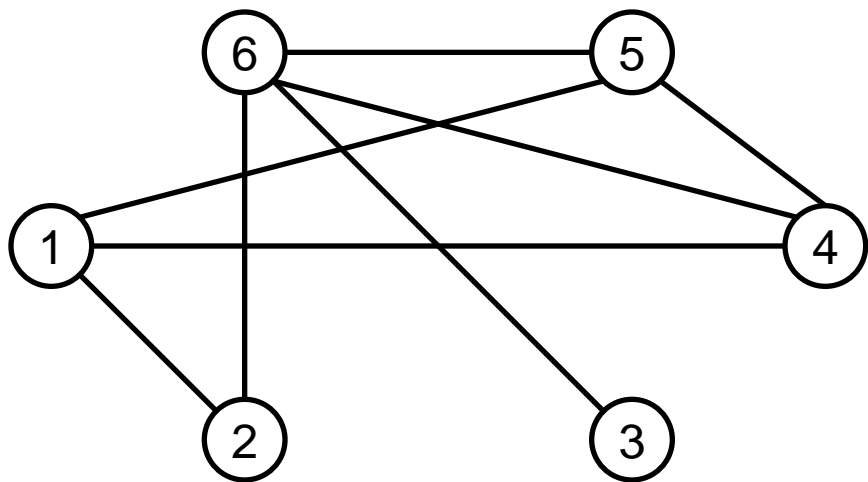


# DIMACS instances



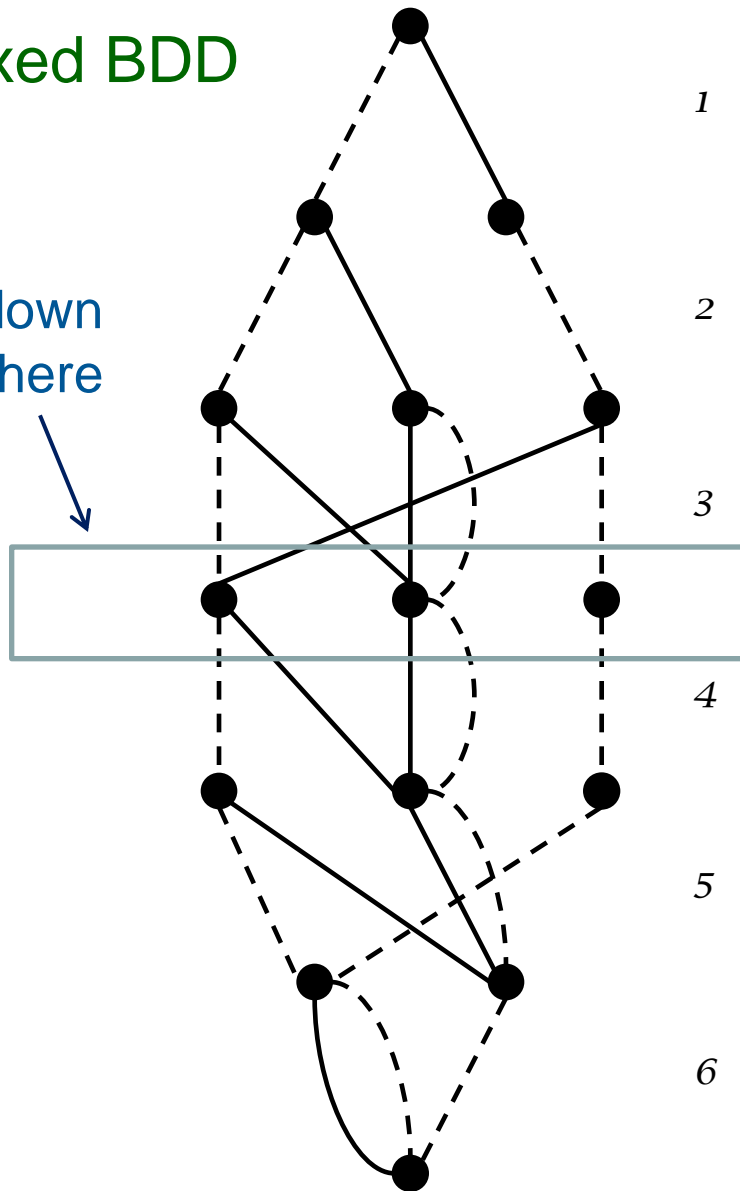
# Branching

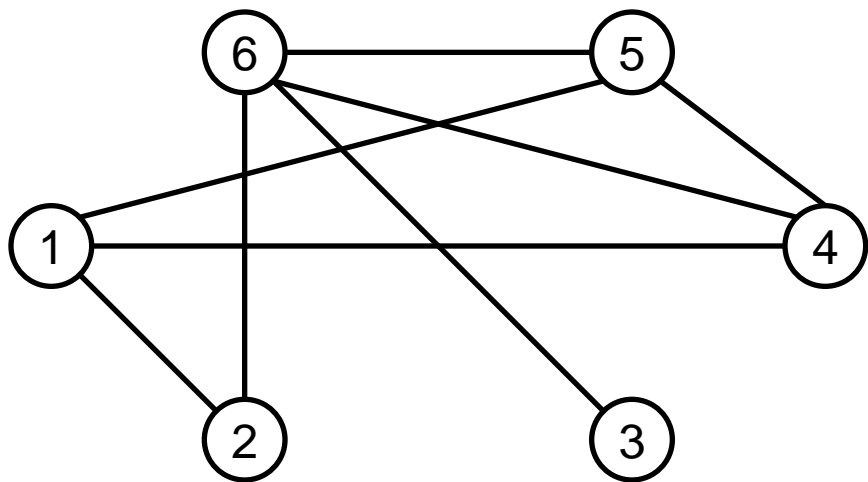
- Novel branching scheme
  - Branch in relaxed decision diagram.
  - Branch on pools of partial solutions.
  - Reduce symmetry.



## Relaxed BDD

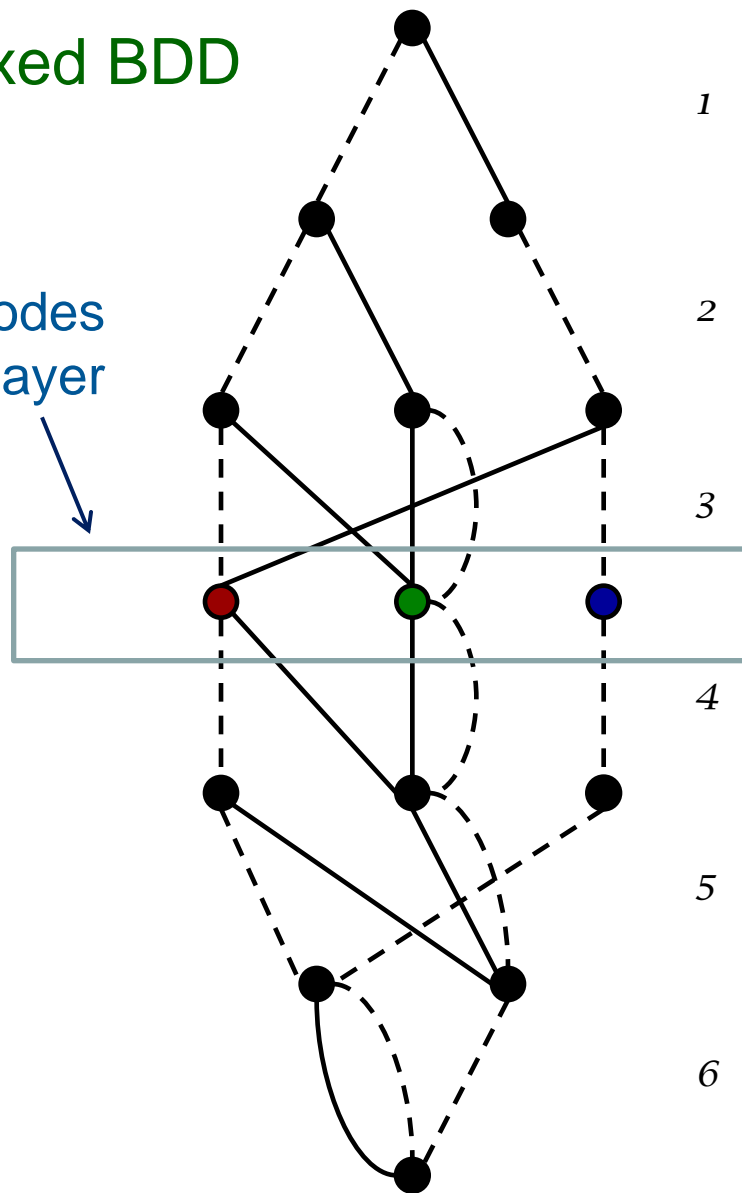
Exact down  
to here

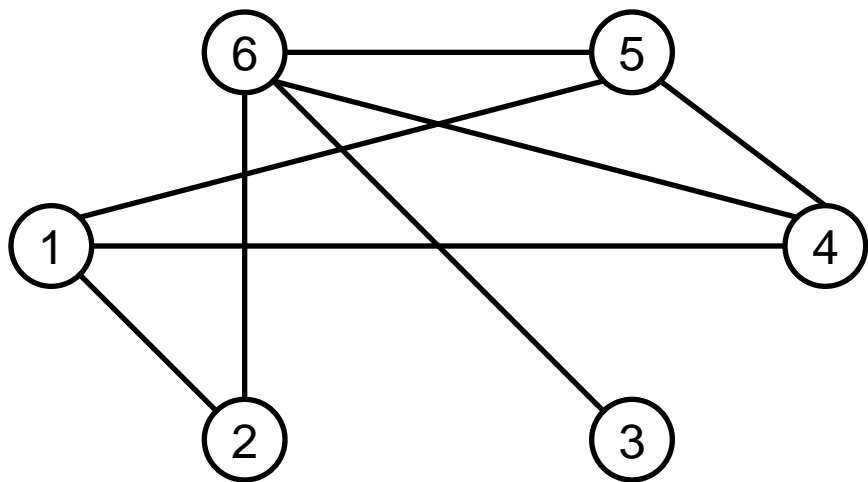




Branch on nodes  
in this layer

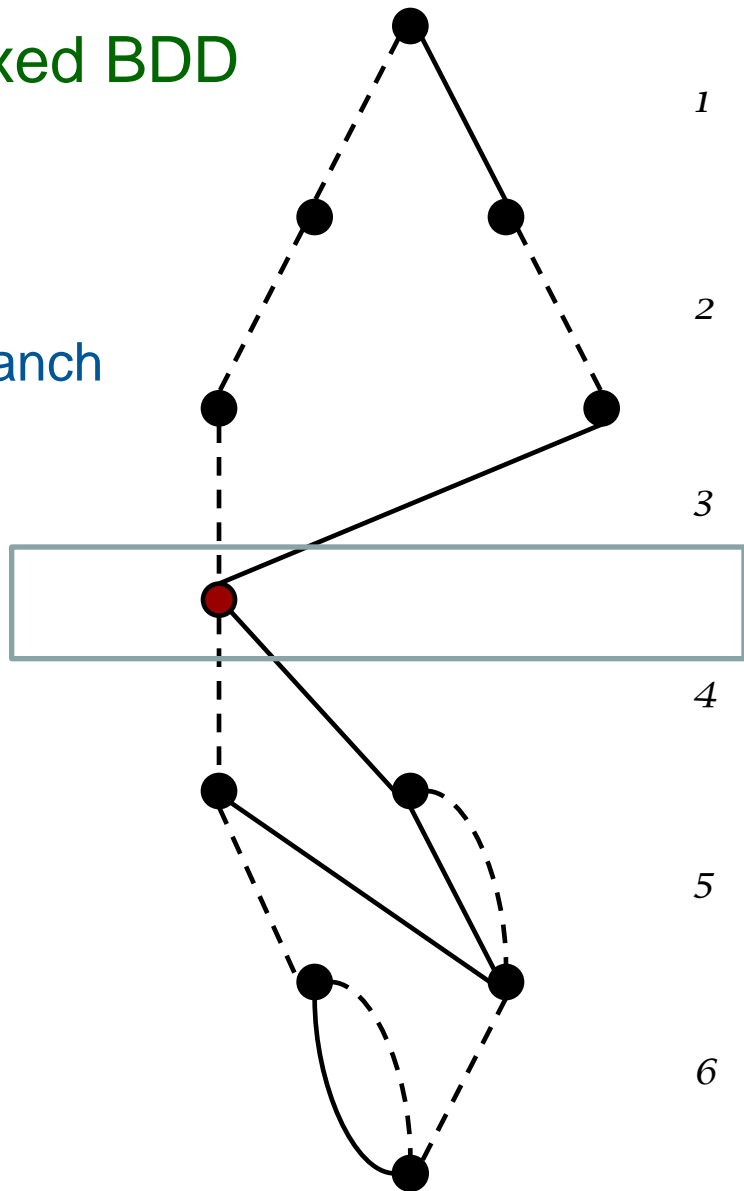
Relaxed BDD

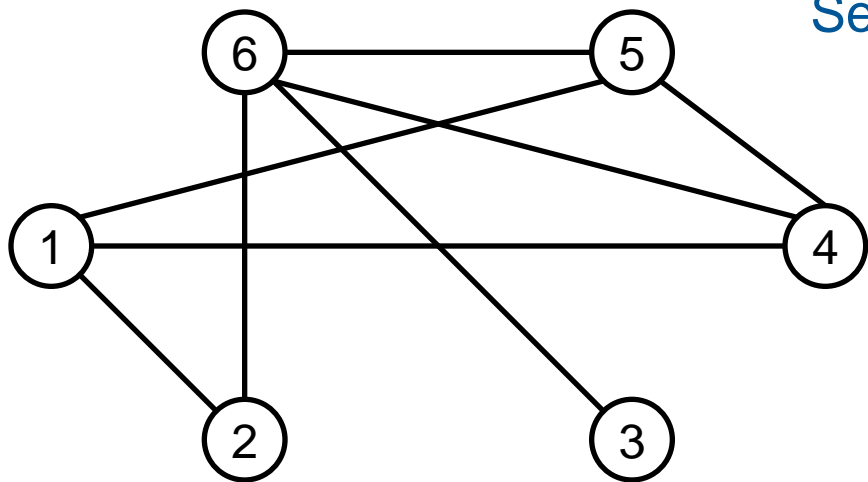




Relaxed BDD

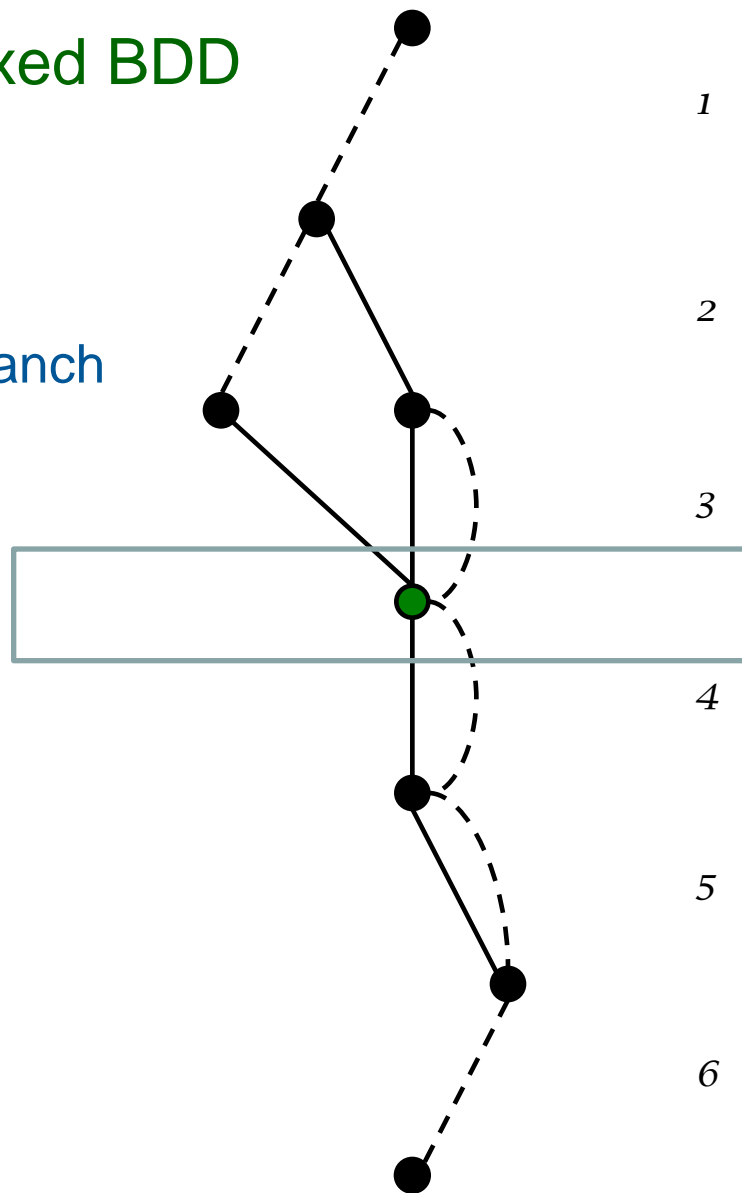
First branch





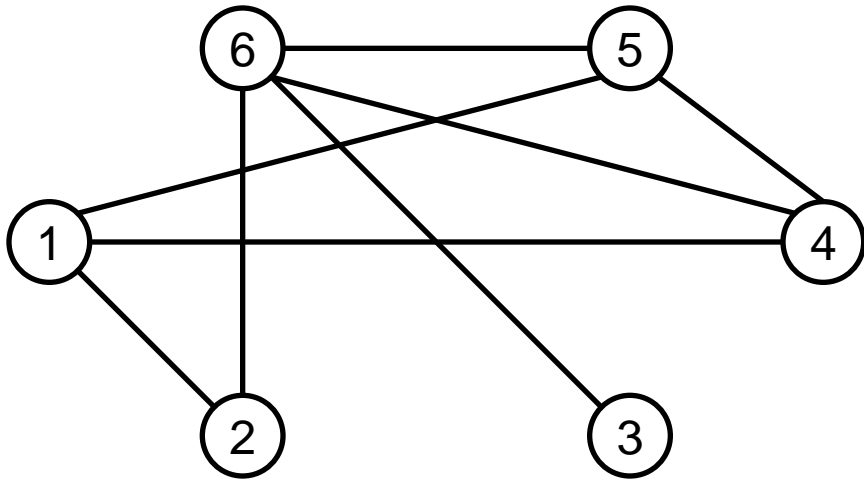
Second branch

Relaxed BDD





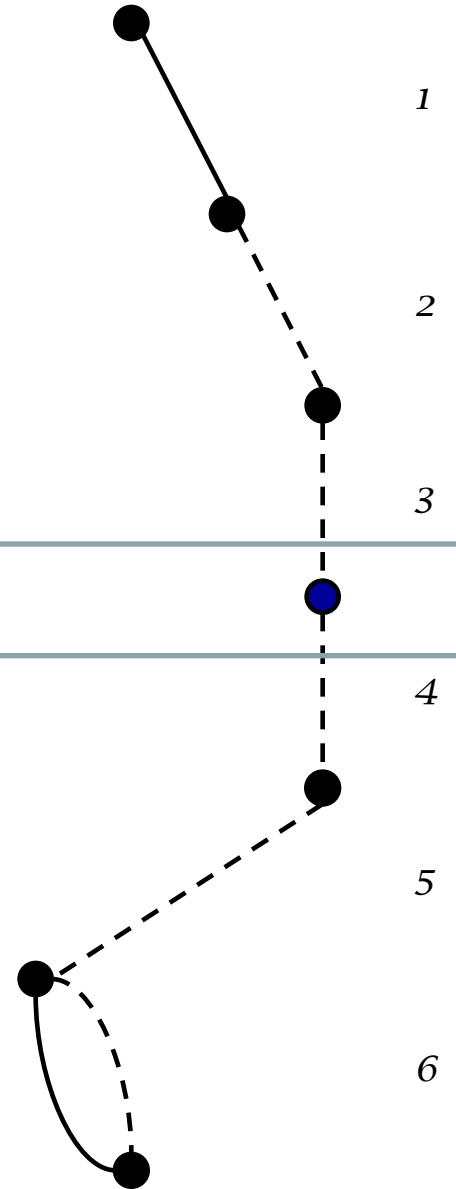
Relaxed BDD



Third branch

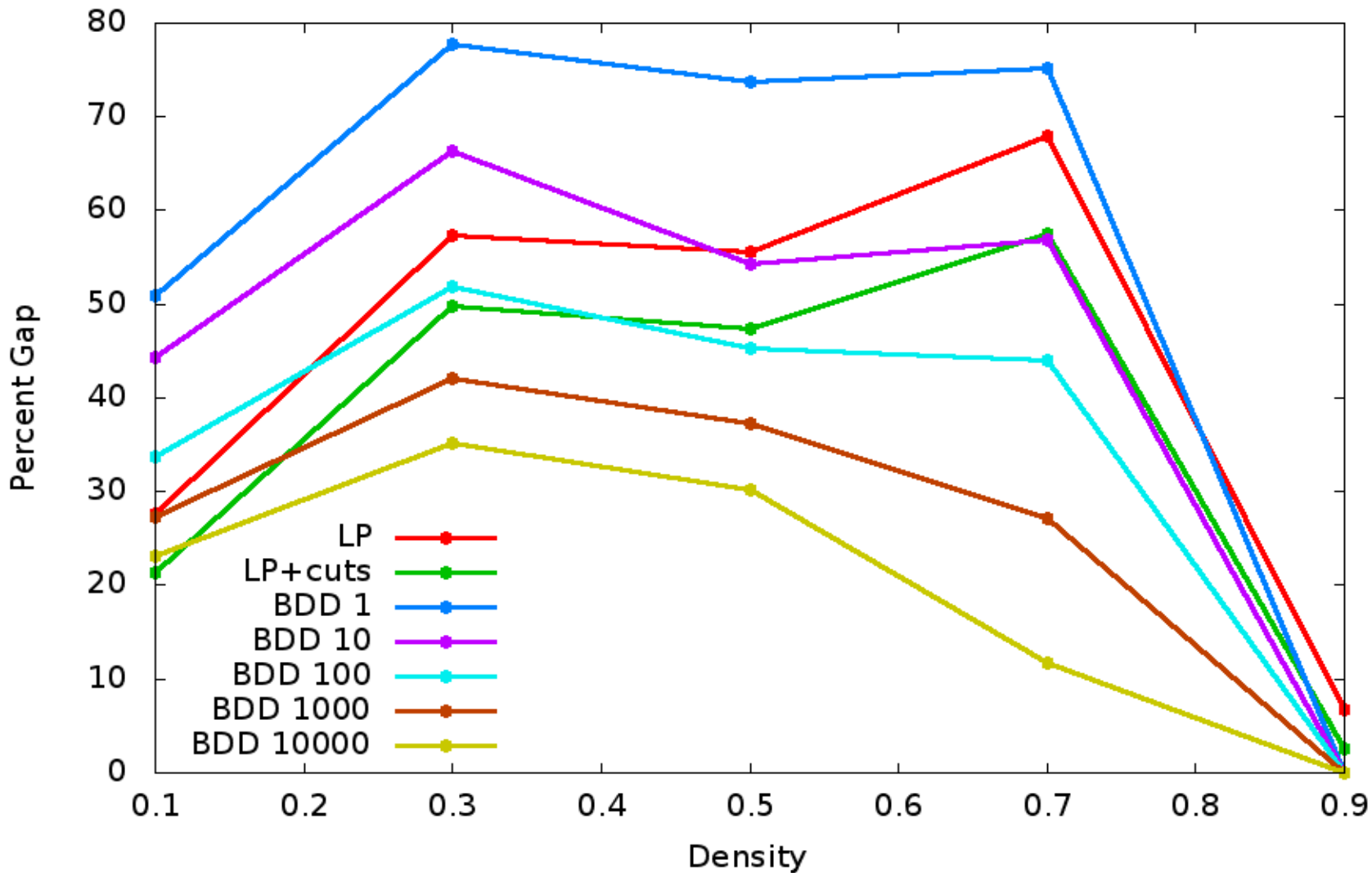


Continue recursively



# DIMACS instances

## Root Node Gap Comparison

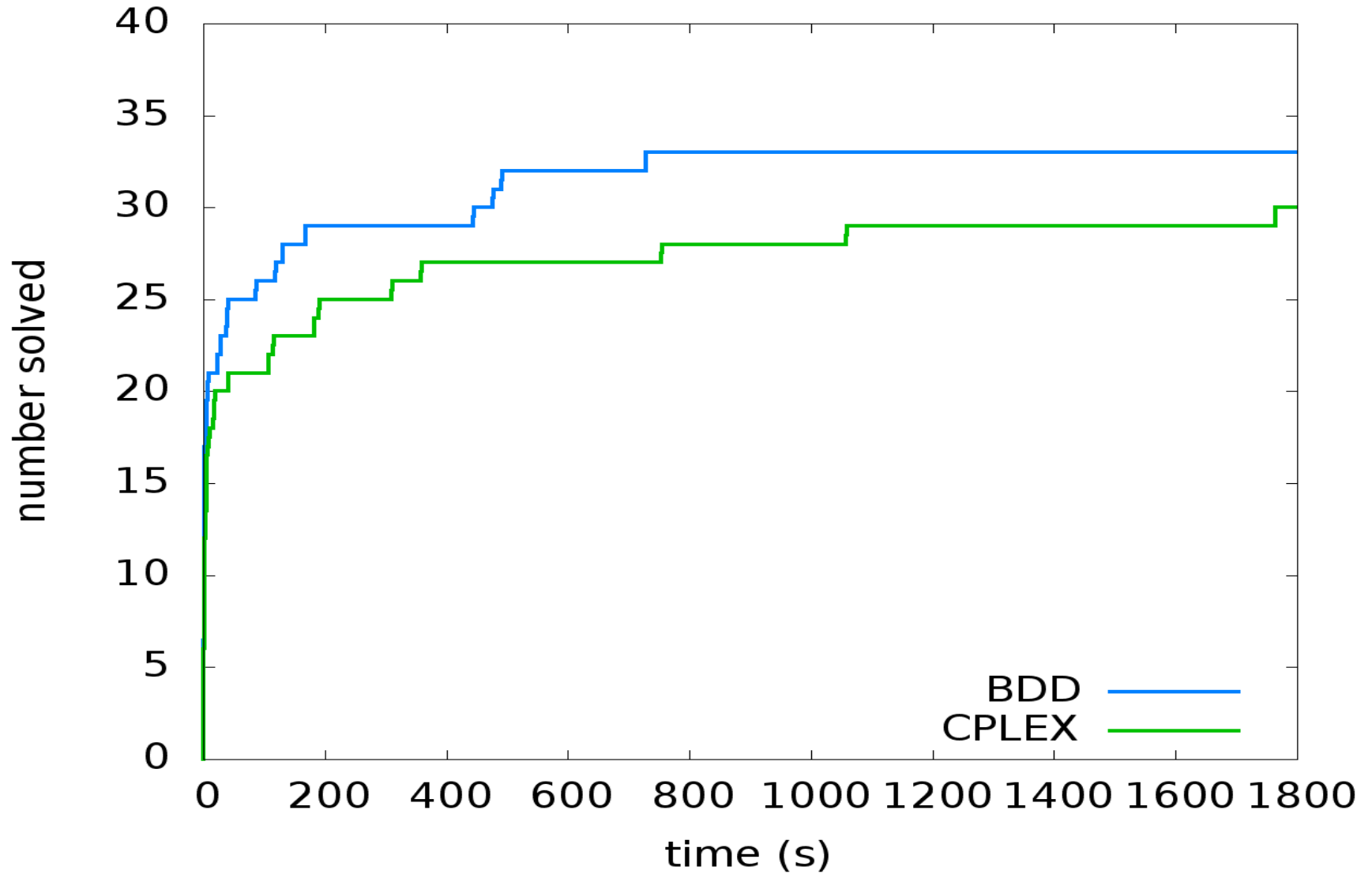


# Put it All Together

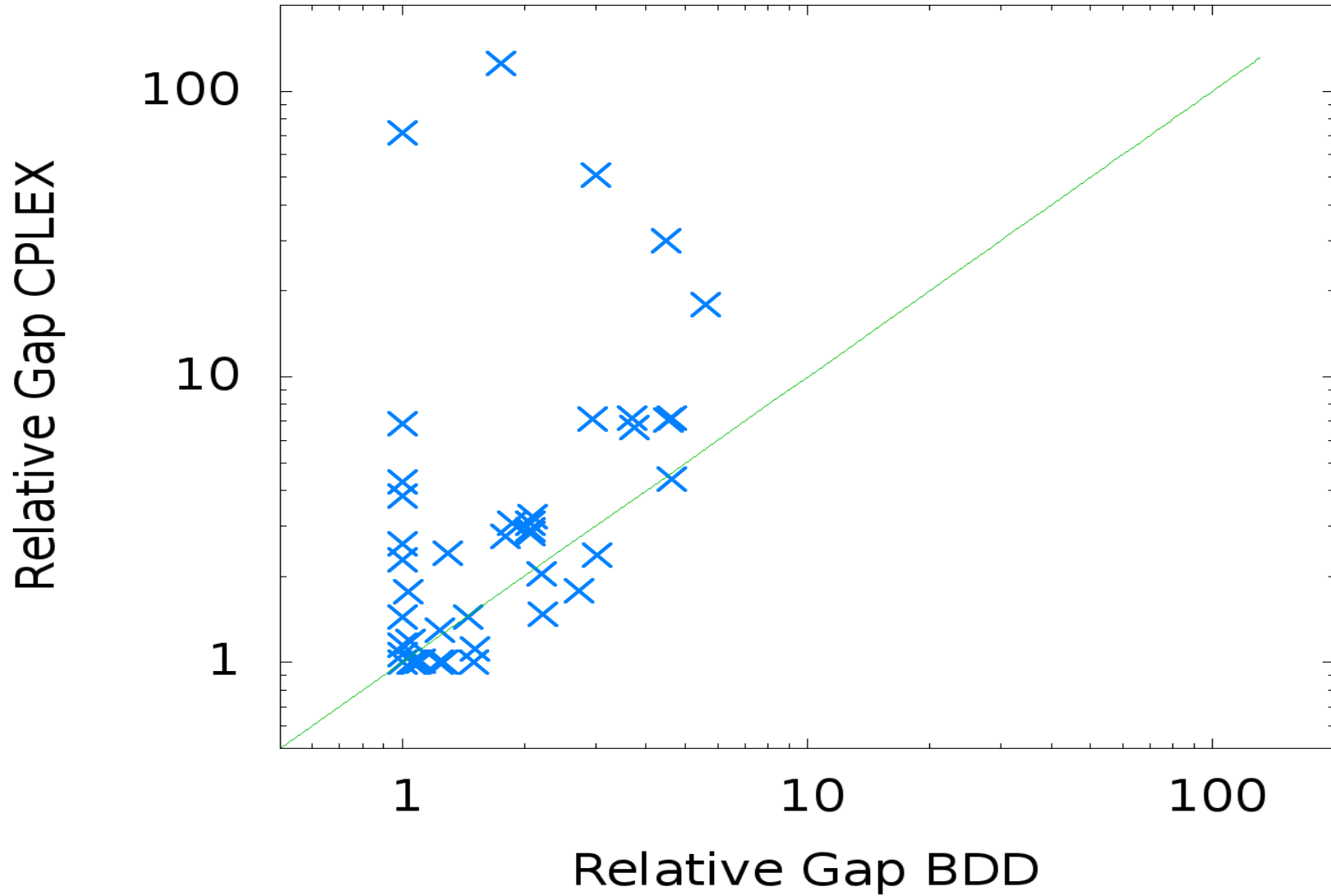
- General solver
  - Bounds from relaxed decision diagram
  - Primal heuristic from restricted decision diagram
  - Propagation in relaxed decision diagram (not used here)
  - Formulation with state variables
  - Branching in decision diagram
- Compare with CPLEX
  - Independent set problem

DIMACS instances

## Number Solved Comparison



# End Gap Comparison



# Ongoing Research

- Additional problems
  - Set covering
  - MaxSAT
  - Minimum bandwidth (linear arrangement)
- Omit complicating variables from relaxed decision diagram
  - Reduce coupling from constraints
  - Decision diagram becomes much tighter relaxation